

Braided 2 Categories and the Drinfel'd Center

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Contents

1	Let's recall some stuff first	2
2	Braided Monoidal 2-Categories	3
3	The Drinfel'd Center of a Monoidal 2-Category	9
3.1	The Base 2-Category $\mathcal{Z}(\mathcal{A})$	9
3.2	The Monoidal Structure	14
3.3	The Braiding	15
3.4	Some Facts	16
4	Braided Module Categories	17
4.1	Definitions	17
4.2	α -Inductions and the Intermediate Category $\mathbf{A}(\mathcal{B})$	22
5	Explicit Description of $\mathcal{Z}(\Sigma\mathcal{B})$	24

1 Let's recall some stuff first

Recall the periodic table of k -tuply monoidal n -categories:

	n=0	n=1	n=2	...
k=0	(pointed) set	(pointed) category	(pointed) 2-category	...
k=1	monoid	monoidal category	monoidal 2-category	...
k=2	abelian monoid	braided monoidal category	braided monoidal 2-category	...
k=3	"	symmetric monoidal category	symplectic monoidal 2-category	...
k=4	"	"	symmetric monoidal 2-category	...
k=5	"	"	"	...
⋮	⋮	⋮	⋮	⋮

We see that a *monoidal 2-category* is just a 3-category with one object.

We will however work with the strictest version possible while still retaining full generality:

Notation 1.1 — By a *monoidal 2-category* we mean a Gray monoid, i.e.

(S1) a 2-category \mathcal{A} together with

(S2) a 2-functor $\boxtimes : \mathcal{A} \boxtimes_G \mathcal{A} \rightarrow \mathcal{A}$ where \boxtimes_G is the Gray product,

such that:


(A1) \boxtimes is (strictly) associative with unit $\mathbf{1} \in \mathbb{C}$.

2 Braided Monoidal 2-Categories

A braided monoidal 2-category is just a 4-category with a single object and a single 1-morphism. Unfortunately, 4-categories are *hard*. However, we can provide a semi-strict (Gray) definition for braided monoidal 2-categories. This definition is due to [3].

Definition 2.1 — An *braided monoidal 2-category* consists of:

(S1) A monoidal 2-category $(\mathcal{A}, \boxtimes, I)$

(S2) A 2-natural equivalence $\beta_{-, -} : - \boxtimes - \xrightarrow{\sim} - \boxtimes -$ represented by 

(S3) Invertible modifications $R_{(-|-,-)}$ and $R_{(-,-|-)}$ where:



satisfying:

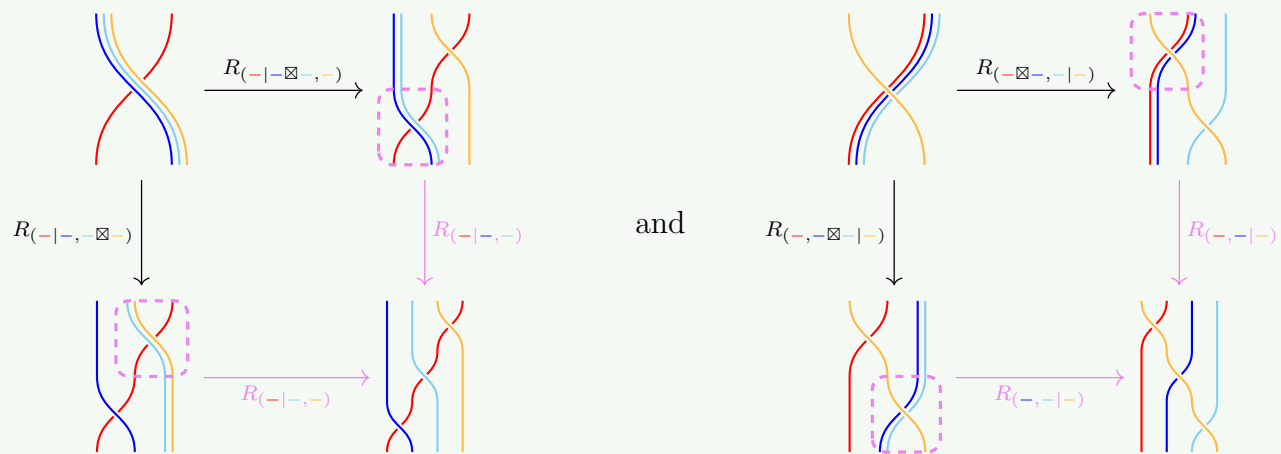
(A1) the (1,3)-crossing and (3,1)-crossing axioms,

(A2) the (2,2)-crossing axiom,

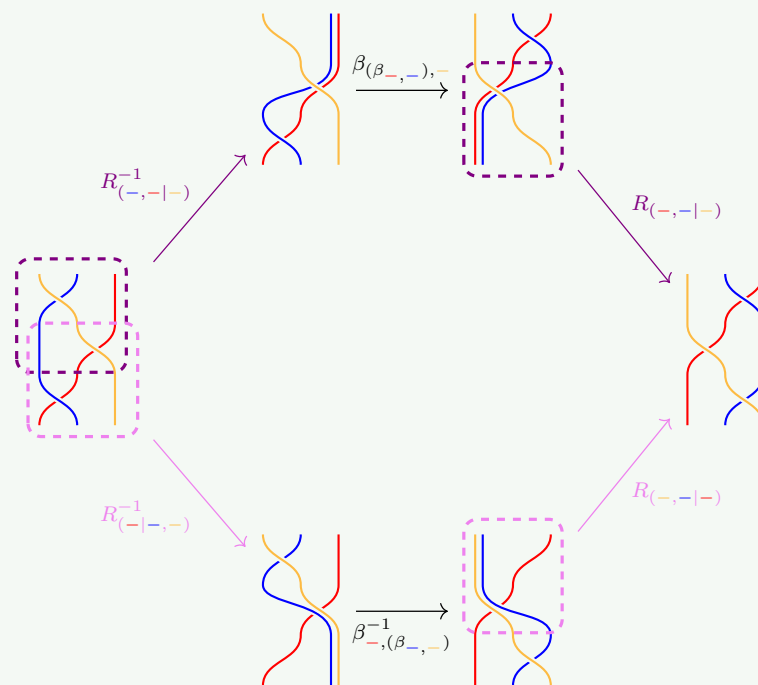
(A3) the Yang-Baxter axiom, and

(A4) unit axioms.

(A1) The (1,3)-crossing and (3,1)-crossing axioms:

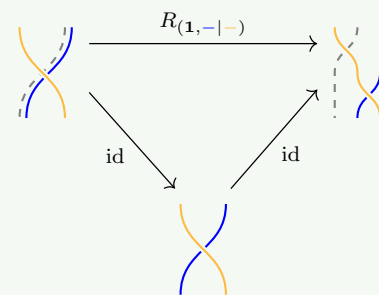
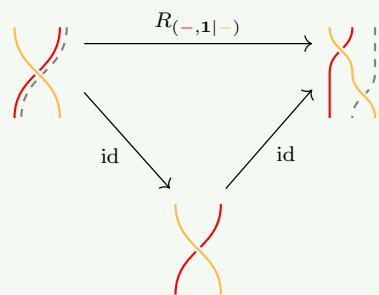
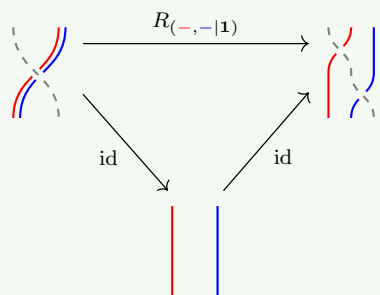
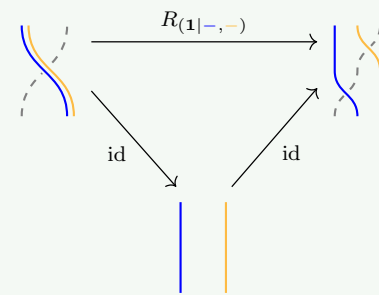
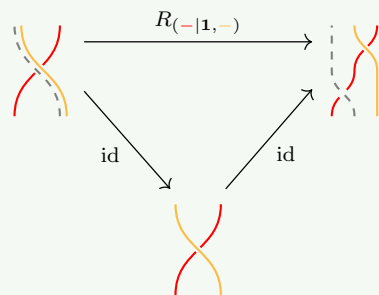
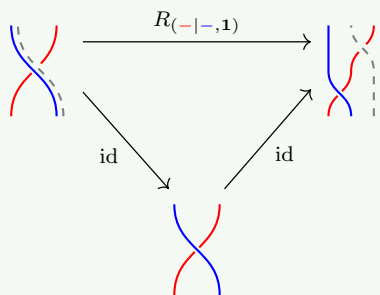


(A3) The Yang-Baxter axiom:



(A4) Unit axioms:

$$\begin{array}{c} \text{---} \\ \diagdown \\ \diagup \\ \text{---} \end{array} = \text{---} \qquad \begin{array}{c} \text{---} \\ \diagup \\ \diagdown \\ \text{---} \end{array} = \text{---}$$



Example 2.2 — Examples of braided monoidal 2-categories

- The Drinfeld center $\mathcal{Z}(\mathcal{A})$ of a monoidal 2-category \mathcal{A} ,
- The braided monoidal 2-category $\text{BrMod}\text{-}\mathcal{B}$ of braided module categories over a braided fusion 1-category \mathcal{B} .
- The braided fusion 2-categories \mathcal{S} and \mathcal{T}

Exercise 2.1. Come up with more examples of braided monoidal 2-categories.

3 The Drinfel'd Center of a Monoidal 2-Category

This definition is also from [3].

3.1 The Base 2-Category $\mathcal{Z}(\mathcal{A})$

Definition 3.1 — Given a monoidal 2-category \mathcal{A} , we define $\mathcal{Z}(\mathcal{A})$ to be the 2-category consisting of:

(0) Objects are triples $(A, \beta_{A,-}, R_{(A|-, -)})$ consisting of:

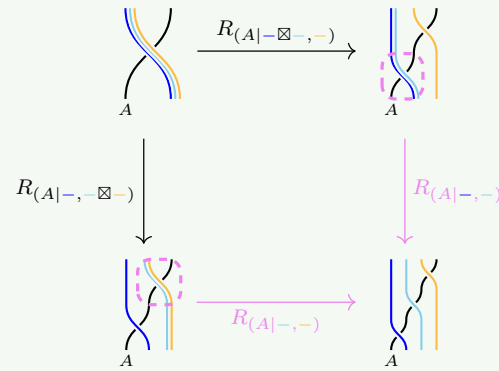
- an object $A \in \mathcal{A}$,
- a 2-natural equivalence $\beta_{A,-} : A \boxtimes - \xrightarrow{\sim} - \boxtimes A$ represented by

$$\begin{array}{c} \text{A} \\ \diagup \quad \diagdown \\ \text{A} \quad \text{X} \end{array} := \left\{ \begin{array}{c} \text{A} \\ \diagdown \quad \diagup \\ \text{A} \quad \text{X} \end{array} = \beta_{A,X} : A \boxtimes X \rightarrow X \boxtimes A \right\}_{X \in \mathcal{A}} \cup \left\{ \beta_{A,f} : \begin{array}{c} \text{X}' \quad \text{A} \\ \diagdown \quad \diagup \\ \text{A} \quad \text{X} \\ \text{f} \end{array} \Rightarrow \begin{array}{c} \text{X}' \quad \text{A} \\ \text{f} \\ \text{A} \quad \text{X} \end{array} \right\}_{f: X \rightarrow X' \text{ in } \mathcal{A}}$$

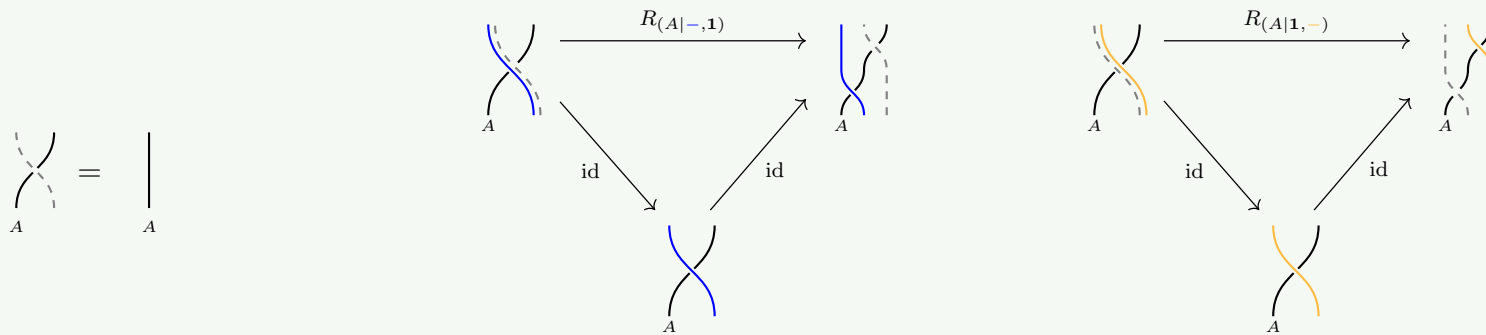
- an invertible modification $R_{(A|-, -)}$ where $\begin{array}{c} \text{A} \\ \diagup \quad \diagdown \\ \text{A} \quad \text{A} \end{array} \xrightarrow[\sim]{R_{(A|-, -)}} \begin{array}{c} \text{A} \\ \diagdown \quad \diagup \\ \text{A} \quad \text{A} \end{array}$

satisfying the (1,3)-crossing and unit axioms.

- The (1,3)-crossing axiom:



- Unit axioms:



(1) A 1-morphism $(f, \beta_{f,-}) : (A, \beta_{A,-}, R_{(A|-,-)}) \rightarrow (A', \beta_{A',-}, R_{(A'|-,-)})$ consists of:

- a 1-morphism $f : A \rightarrow A'$ in \mathcal{A} ,

- an invertible modification represented by

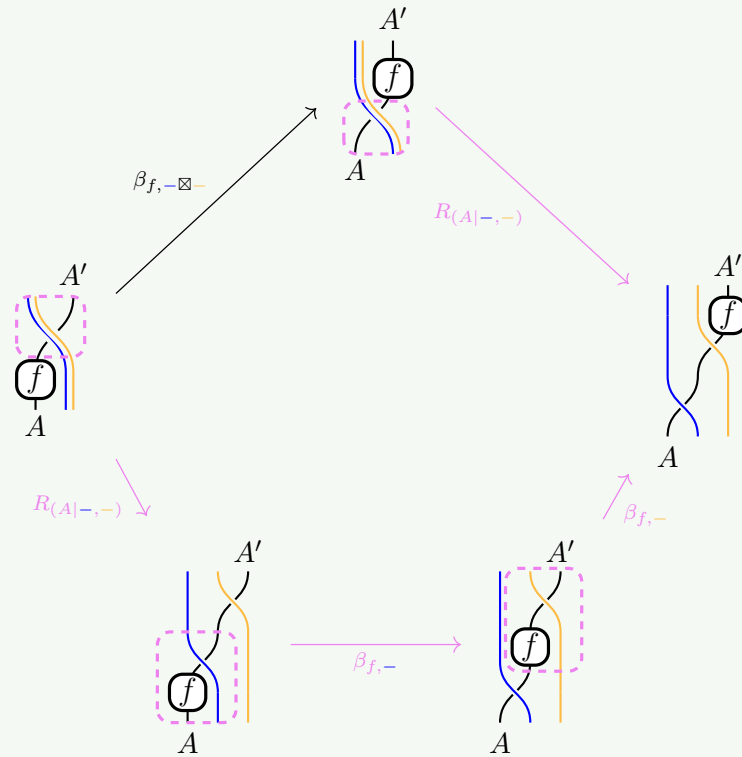
$$\begin{array}{c} A' \\ \diagup \\ \textcircled{f} \\ \diagdown \\ A \end{array} \xrightarrow[\beta_{f,-}]{\sim} \begin{array}{c} A' \\ \textcircled{f} \\ A \end{array} := \left\{ \beta_{f,X} : \begin{array}{c} X \quad A' \\ \diagup \quad \diagdown \\ \textcircled{f} \\ \diagdown \quad \diagup \\ A \quad X \end{array} \Rightarrow \begin{array}{c} X \quad A' \\ \textcircled{f} \\ A \quad X \end{array} \right\}_{X \in \mathcal{A}}$$

such that $\beta_{f,-}$ satisfies a unit axiom and $R_{(A|-,-)}$ becomes natural in $f : A \rightarrow A'$.

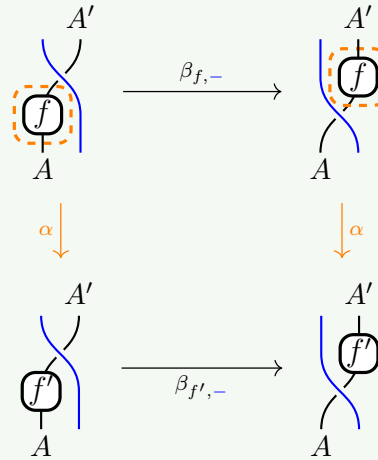
- Unit axiom:

$$\begin{array}{ccc} \begin{array}{c} 1 \quad A' \\ \diagup \quad \diagdown \\ \textcircled{f} \\ \diagdown \quad \diagup \\ A \quad 1 \end{array} & \xrightarrow{\beta_{f,1}} & \begin{array}{c} 1 \quad A' \\ \diagup \quad \diagdown \\ \textcircled{f} \\ \diagdown \quad \diagup \\ A \quad 1 \end{array} \\ \text{id} \searrow & & \nearrow \text{id} \\ & A' & \\ & \textcircled{f} & \\ & A & \end{array}$$

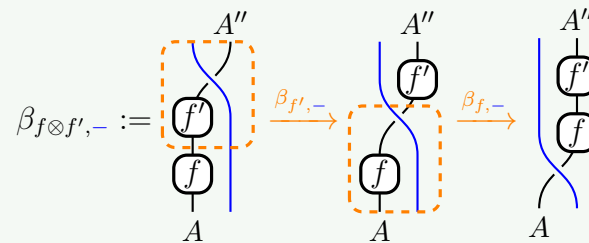
- Naturality of $R_{(A|-, -)}$:



(2) A 2-morphism $\alpha : (f, \beta_{f,-}) \Rightarrow (f', \beta_{f',-})$ is a 2-morphism $\alpha : f \Rightarrow f'$ in \mathcal{A} such that $\beta_{f,-}$ becomes 2-natural in $\alpha : f \Rightarrow f'$, i.e.



For $(A, \beta_{A,-}, R_{(A|-,-)}) \xrightarrow{(f, \beta_{f,-})} (A', \beta_{A',-}, R_{(A'|-,-)}) \xrightarrow{(f', \beta_{f',-})} (A'', \beta_{A'',-}, R_{(A''|-,-)})$ in $\mathcal{Z}(\mathcal{A})$, their 1-composite $(f, \beta_{f,-}) \otimes (f', \beta_{f',-})$ is defined to be $(f \otimes f', \beta_{f \otimes f',-})$ where:



The compositions \otimes and \circ of 2-morphisms in $\mathcal{Z}(\mathcal{A})$ are the same as in \mathcal{A} .

3.2 The Monoidal Structure

Definition 3.2 — For objects $(A, \beta_{A,-}, R_{(A|-, -)})$ and $(B, \beta_{B,-}, R_{(B|-, -)})$, we define $(A, \beta_{A,-}, R_{(A|-, -)}) \boxtimes (B, \beta_{B,-}, R_{(B|-, -)})$ to be $(A \boxtimes B, \beta_{A \boxtimes B, -}, R_{(A \boxtimes B|-, -)})$ where:

$$\beta_{A \boxtimes B, -} = \begin{array}{c} \text{Diagram 1: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 2: Two strands, left black, right orange, crossing. Labels } A \text{ and } B \text{ at bottom.} \end{array} := \begin{array}{c} \text{Diagram 3: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 4: Two strands, left black, right orange, crossing. Labels } A \text{ and } B \text{ at bottom.} \end{array}$$

$$R_{(A \boxtimes B|-, -)} : \begin{array}{c} \text{Diagram 5: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 6: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A red dashed box around the crossing, a blue dashed box around the strands.} \\ \text{Diagram 7: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 8: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 9: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \end{array} \xrightarrow{R_{(A|-, -)}} \begin{array}{c} \text{Diagram 10: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 11: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \end{array} \cong \begin{array}{c} \text{Diagram 12: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \\ \text{Diagram 13: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \end{array} = \begin{array}{c} \text{Diagram 14: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom.} \end{array}$$

For an object $(A, \beta_{A,-}, R_{(A|-, -)})$ and a 1-morphism $(g : B \rightarrow B', \beta_{g,-})$, we define:

$$(A, \beta_{A,-}, R_{(A|-, -)}) \boxtimes (g, \beta_{g,-}) := (A \boxtimes g, \beta_{A \boxtimes g, -}) \quad \text{and} \quad (g, \beta_{g,-}) \boxtimes (A, \beta_{A,-}, R_{(A|-, -)}) := (g \boxtimes A, \beta_{g \boxtimes A, -}) \quad \text{where:}$$

$$\beta_{A \boxtimes g, -} := \begin{array}{c} \text{Diagram 15: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A red dashed box around the crossing, a black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 16: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 17: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \end{array} \xrightarrow{R_{g,-}} \begin{array}{c} \text{Diagram 18: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 19: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \end{array} \cong \begin{array}{c} \text{Diagram 20: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 21: Two strands, left orange, right black, crossing. Labels } A \text{ and } B \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \end{array}$$

and

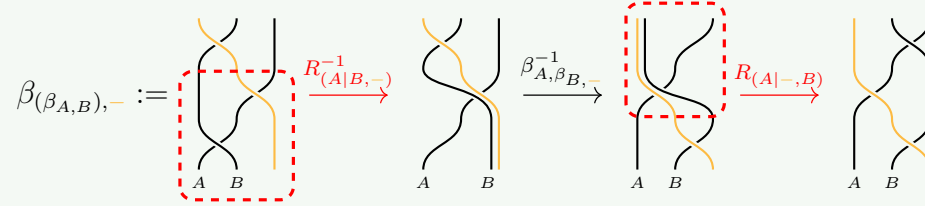
$$\beta_{g \boxtimes A, -} := \begin{array}{c} \text{Diagram 22: Two strands, left orange, right black, crossing. Labels } B \text{ and } A \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 23: Two strands, left orange, right black, crossing. Labels } B \text{ and } A \text{ at bottom. A red dashed box around the crossing, a black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 24: Two strands, left orange, right black, crossing. Labels } B \text{ and } A \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \end{array} \xrightarrow{R_{g,-}} \begin{array}{c} \text{Diagram 25: Two strands, left orange, right black, crossing. Labels } B \text{ and } A \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \\ \text{Diagram 26: Two strands, left orange, right black, crossing. Labels } B \text{ and } A \text{ at bottom. A black circle labeled } g \text{ around the crossing.} \end{array}$$

The other \boxtimes -products are defined as in \mathcal{A} .

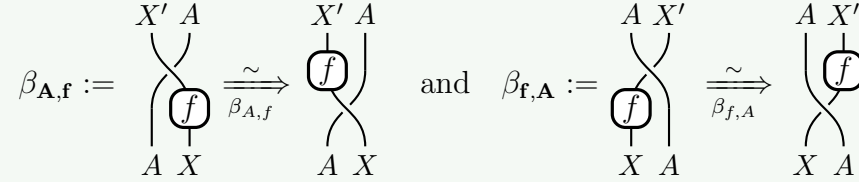
3.3 The Braiding

Definition 3.3 — We define the braiding $(\beta_{-,-}, R_{(-|-,-)}, R_{(-,-|-)})$ on $\mathcal{Z}(\mathcal{A})$ as follows:

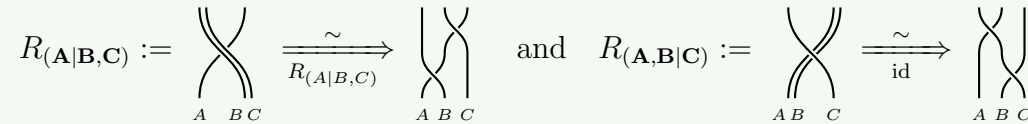
(β_0) For objects $\mathbf{A} = (A, \beta_{A,-}, R_{(A|-,-)})$ and $\mathbf{B} = (B, \beta_{B,-}, R_{(B|-,-)})$, we define $\beta_{\mathbf{A},\mathbf{B}} := (\beta_{A,B}, \beta_{(\beta_{A,B}),-})$ where



(β_1) For an object $\mathbf{A} = (A, \beta_{A,-}, R_{(A|-,-)})$ and a 1-morphism $\mathbf{f} = (f : X \rightarrow X', \beta_{f,-})$, we define



(R) For objects $\mathbf{A} = (A, \beta_{A,-}, R_{(A|-,-)})$, $\mathbf{B} = (B, \beta_{B,-}, R_{(B|-,-)})$, and $\mathbf{C} = (C, \beta_{C,-}, R_{(C|-,-)})$, we define



3.4 Some Facts

Theorem 3.4. Given any monoidal 2-category $(\mathcal{A}, \boxtimes, \mathbf{1})$, the Drinfeld center $\mathcal{Z}(\mathcal{A})$ is a braided monoidal 2-category.

Proof. An incomplete proof of this theorem appears in [Baez + Neuchl], which is completed and corrected by [Crans]. \square

Theorem 3.5. Given any braided monoidal 2-category $(\mathcal{A}, \boxtimes, \mathbf{1}, \beta_{-, -}, R_{(-|-,-)}, R_{(-,-|-)})$, there exists an embedding $\zeta : \mathcal{A} \hookrightarrow \mathcal{Z}(\mathcal{A})$ given by:

$$\begin{aligned} (A \in \mathcal{A}) &\mapsto (A, \beta_{A, -}, R_{(A|-,-)}) \\ (f : A \rightarrow A') &\mapsto (f, \beta_{f, -}) \\ (\alpha : f \Rightarrow f') &\mapsto \alpha. \end{aligned}$$

Note that this implies that every braided monoidal 2-category is equivalent to one for which $R_{(-,-|-)}$ is trivial.

4 Braided Module Categories

This section is based on chapter 3 of [4]. From here on out, \mathcal{B} will always be braided fusion 1-category.

4.1 Definitions

Definition 4.1 — A braided (right) module category of \mathcal{B} is:

(S1) a finite semisimple (right) \mathcal{B} -module category $(\mathcal{M}, \triangleleft : \mathcal{M} \boxtimes \mathcal{B} \rightarrow \mathcal{M}, \dots)$

(S2) a natural isomorphism $\sigma_{-, -} : \triangleleft \Rightarrow \triangleleft$ represented by: $\begin{array}{c} \text{blue} \\ \text{red} \end{array} \left. \begin{array}{c} \text{blue} \\ \text{red} \end{array} \right\} = \left\{ \begin{array}{c} \text{blue} \\ \text{red} \end{array} \right\} := \sigma_{m,x} : m \triangleleft x \rightarrow m \triangleleft x \right\}_{m \in \mathcal{M}, x \in \mathcal{B}}$

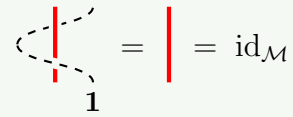
satisfying:

(A1) a unit axiom,

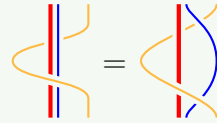
(A2) compatibility with \triangleleft and braiding, and

(A3) compatibility with the \otimes -product on \mathcal{B} .

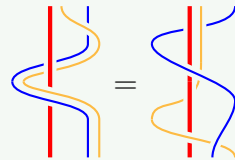
(A1) Unit axiom:


$$\text{loop}_1 = \text{line} = \text{id}_{\mathcal{M}}$$

(A2) Compatibility with \triangleleft and braiding:


$$\text{braiding} = \text{braiding}$$

(A3) Compatibility with the \otimes -product on \mathcal{B} :


$$\text{braiding} = \text{braiding}$$

Remark 4.1. The term *braided* in the previous definition is justified as follows:

- Recall that the Artin braid group of type B is the group B_n generated by $\sigma_0, \dots, \sigma_{n-1}$ subject to the relations:

$$\begin{aligned} \sigma_1 \sigma_0 \sigma_1 \sigma_0 &= \sigma_0 \sigma_1 \sigma_0 \sigma_1, \\ \sigma_i \sigma_j &= \sigma_j \sigma_i && \text{whenever } |i - j| > 1, \\ \sigma_i \sigma_{i+1} \sigma_i &= \sigma_{i+1} \sigma_i \sigma_{i+1} && \text{for } i = 1, \dots, n - 1. \end{aligned}$$

- Given $X_1, \dots, X_{n-1} \in \mathcal{B}$ and $M \in \mathcal{M}$, there are isomorphisms

$$M \triangleleft X_1 \triangleleft \cdots \triangleleft X_{n-1} \rightarrow M \triangleleft X_{\sigma(1)} \triangleleft \cdots \triangleleft X_{\sigma(n-1)}, \quad \text{for } \sigma \in B_n,$$

compatible with the composition of braids.

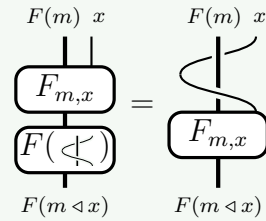
Definition 4.2 — A braided module functor $(\mathcal{M}, \triangleleft, \sigma_{-, -}) \xrightarrow{(F, F_{-, -})} (\mathcal{M}', \triangleleft', \sigma'_{-, -})$ is:

(S1) a linear functor $F : \mathcal{M} \rightarrow \mathcal{M}'$, and

(S2) a natural isomorphism $F_{-, -} = \{F_{m,x} : F(m \triangleleft x) \rightarrow F(m) \triangleleft' x\}_{m \in \mathcal{M}, x \in \mathcal{B}}$

such that:

(A1) $F_{m,x} \circ F(\sigma_{m,x}) = \sigma'_{f(m),x} \circ F_{m,x}$.



Remark 4.2. Note that being a *braided* module functor is a property of a module functor, *not* extra structure.

Definition 4.3 — A transformation $\alpha : (F, F_{-, -}) \Rightarrow (F', F'_{-, -})$ of braided module functors is simply a natural transformation $\alpha : F \Rightarrow F'$ of the underlying \mathcal{B} -module functors.

Definition 4.4 — We define $\text{BrMod-}\mathcal{B}$ to be the 2-category of braided modules over \mathcal{B} , braided module functors, and natural transformations.

Example 4.5 — Any braided monoidal functor $F : \mathcal{B} \rightarrow \mathcal{C}$ of braided fusion 1-categories equips \mathcal{C} with the structure of a braided \mathcal{B} -module category with $c \triangleleft b := c \otimes_{\mathcal{C}} F(b)$ and module braiding:

$$\begin{array}{c} \text{Diagram 1} \\ \text{Diagram 2} \\ F(-) \end{array} := \left\{ \begin{array}{c} \text{Diagram 3} \\ m \quad F(x) \end{array} = \beta_{F(x),m}^{\mathcal{C}} \circ \beta_{m,F(x)}^{\mathcal{C}} : m \triangleleft x \rightarrow m \triangleleft x \right\}_{m \in \mathcal{C}, x \in \mathcal{B}}$$

- In particular, when \mathcal{C} is a braided tensor category containing \mathcal{B} , we can equip \mathcal{C} with this braided \mathcal{B} -module category structure. In this case, the category of braided module endofunctors is braided equivalent to:

$$\mathcal{Z}_{(2)}(\mathcal{B} \subset \mathcal{C}) := \{c \in \mathcal{C} \mid \beta_{c,b} \circ \beta_{b,c} = \text{id}_{b \otimes c} \text{ for all } b \in \mathcal{B}\}$$

- A special case of this is when $\mathcal{C} = \mathcal{B}$, where we see \mathcal{B} as the rank one free braided \mathcal{B} -module category. Then, the category of braided module endofunctors of \mathcal{B} is braided equivalent to the Müger center $\mathcal{Z}_{(2)}(\mathcal{B}) := \mathcal{Z}_{(2)}(\mathcal{B} \subset \mathcal{B})$.

4.2 α -Inductions and the Intermediate Category $\mathbf{A}(\mathcal{B})$

Definition 4.6 — The α -inductions [2] for a right \mathcal{B} -module category \mathcal{M} are tensor functors:

$$\alpha_{\mathcal{M}}^{\pm} : \mathcal{B} \rightarrow \text{End}_{\mathcal{B}}(\mathcal{M}), \quad \alpha_{\mathcal{M}}^{\pm}(X) := - \triangleleft X = \begin{array}{c} | \\ \hline X \end{array}, \text{ for every } X \in \mathcal{B}.$$

where $\text{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{M})$ is the category of right exact \mathcal{B} -module endofunctors of \mathcal{M} . The \mathcal{B} -module functor structures on $\alpha_{\mathcal{M}}^{\pm}(X) : \mathcal{M} \rightarrow \mathcal{M}$ do differ and are given by:

$$\begin{array}{ccc} \alpha_{\mathcal{M}}^+(X)(M \triangleleft Y) & = & \begin{array}{c} M \quad Y \quad X \\ | \quad \diagdown \quad / \\ \hline M \quad X \quad Y \end{array} \\ \alpha_{\mathcal{M}}^+(X)_{M,Y} \uparrow & := & \\ \alpha_{\mathcal{M}}^+(X)(M) \triangleleft Y & = & \end{array} \quad \text{and} \quad \begin{array}{ccc} \alpha_{\mathcal{M}}^-(X)(M \triangleleft Y) & = & \begin{array}{c} M \quad Y \quad X \\ | \quad \diagdown \quad / \\ \hline M \quad X \quad Y \end{array} \\ \alpha_{\mathcal{M}}^-(X)_{M,Y} \uparrow & := & \\ \alpha_{\mathcal{M}}^-(X)(M) \triangleleft Y & = & \end{array}$$

The monoidal functor structures of $\alpha_{\mathcal{M}}^{\pm}$ also differ and are given by:

$$\begin{array}{ccc} \alpha_{\mathcal{M}}^+(X \otimes Y) & = & \begin{array}{c} | \quad X \quad Y \\ \diagdown \quad / \\ \hline Y \quad X \end{array} \\ (\alpha_{\mathcal{M}}^+)_{X,Y} \uparrow\uparrow & := & \\ \alpha_{\mathcal{M}}^+(X) \circ \alpha_{\mathcal{M}}^+(Y) & = & \end{array} \quad \text{and} \quad \begin{array}{ccc} \alpha_{\mathcal{M}}^-(X \otimes Y) & = & \begin{array}{c} | \quad X \quad Y \\ \diagdown \quad / \\ \hline Y \quad X \end{array} \\ (\alpha_{\mathcal{M}}^-)_{X,Y} \uparrow\uparrow & := & \\ \alpha_{\mathcal{M}}^-(X) \circ \alpha_{\mathcal{M}}^-(Y) & = & \end{array}$$

Remark 4.3. Notice that, by definition, every \mathcal{B} -module functor $F : \mathcal{M} \rightarrow \mathcal{N}$ gives rise to natural transformations $F_{-,X}^\pm$ of \mathcal{B} -module functors:

$$\begin{array}{ccc}
 \mathcal{M} & \xrightarrow{F} & \mathcal{N} \\
 \alpha_{\mathcal{M}}^\pm(X) \downarrow & \nearrow F_{-,X}^\pm & \downarrow \alpha_{\mathcal{N}}^\pm(X) \\
 \mathcal{M} & \xrightarrow{F} & \mathcal{N}
 \end{array}
 \quad \text{for all } X \in \mathcal{B}.$$

Definition 4.7 — Let $\mathbf{A}(\mathcal{B})$ be the 2-category consisting of:

- (0) objects are pairs (\mathcal{M}, η) where \mathcal{M} is a \mathcal{B} -module category and $\eta : \alpha_{\mathcal{M}}^+ \xrightarrow{\cong} \alpha_{\mathcal{M}}^-$ is an isomorphism of tensor functors,
- (1) a 1-morphism $F : (\mathcal{M}, \eta) \rightarrow (\mathcal{N}, \eta')$ is a \mathcal{B} -module functor $F : \mathcal{M} \rightarrow \mathcal{N}$ such that

$$\begin{array}{ccc}
 \mathcal{M} & \xrightarrow{F} & \mathcal{N} \\
 \alpha_{\mathcal{M}}^+(X) \downarrow & \nearrow F_{-,X}^+ & \downarrow \alpha_{\mathcal{N}}^+(X) \\
 \mathcal{M} & \xrightarrow{F} & \mathcal{N}
 \end{array}
 \xrightarrow{\eta'_X}
 \begin{array}{ccc}
 \mathcal{M} & \xrightarrow{F} & \mathcal{N} \\
 \alpha_{\mathcal{M}}^-(X) \downarrow & \nearrow F_{-,X}^- & \downarrow \alpha_{\mathcal{N}}^-(X) \\
 \mathcal{M} & \xrightarrow{F} & \mathcal{N}
 \end{array}$$

- (2) A 2-morphism $\alpha : F \Rightarrow F'$ is just a \mathcal{B} -module natural transformation.

5 Explicit Description of $\mathcal{Z}(\Sigma\mathcal{B})$

This section is based on sections 4.1 and 4.2 of [4]. First recall that $\Sigma\mathcal{B} \cong \mathbf{Mod}\text{-}\mathcal{B}$, the category of finite semisimple (right) \mathcal{B} -module categories, pointed by the rank one free module \mathcal{B} ([5], 1.3.13). The goal of this section is to prove the following:

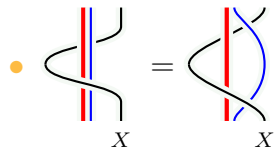
Theorem 5.1. $\mathcal{Z}(\Sigma\mathcal{B}) \cong \mathbf{BrMod}\text{-}\mathcal{B}$

We will proceed by showing $\mathcal{Z}(\mathbf{Mod}\text{-}\mathcal{B}) \cong \mathbf{A}(\mathcal{B}) \cong \mathbf{BrMod}\text{-}\mathcal{B}$.

Lemma 5.2 — There is a canonical 2-equivalence $\mathbf{A}(\mathcal{B}) \cong \mathbf{BrMod}\text{-}\mathcal{B}$.

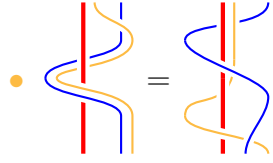
Proof of Lemma. Let \mathcal{M} be a \mathcal{B} -module category. Note that a module braiding $\sigma_{-, -}$ on \mathcal{M} is actually the same thing as a natural isomorphism $\eta : \alpha_{\mathcal{M}}^+ \xrightarrow{\sim} \alpha_{\mathcal{M}}^-$:

$$\begin{array}{ccc} M \triangleleft X & \xrightarrow{\sigma_{M,X}} & M \triangleleft X \\ \parallel & & \parallel \\ \alpha_{\mathcal{M}}^+(X)(M) & \xrightarrow{\eta_X(M)} & \alpha_{\mathcal{M}}^-(X)(M) \end{array} \quad \text{for } X \in \mathcal{B} \text{ and } M \in \mathcal{M}.$$



is equivalent to $\eta_X : \alpha_{\mathcal{M}}^+(X) \xrightarrow{\sim} \alpha_{\mathcal{M}}^-(X)$ being an isomorphism of left \mathcal{B} -module functors, i.e.:

$$\begin{array}{ccc} \alpha_{\mathcal{M}}^+(X)(M) \triangleleft Y & \xrightarrow{(\alpha_{\mathcal{M}}^+(X))_{M,Y}} & \alpha_{\mathcal{M}}^+(X)(M \triangleleft Y) \\ \eta_X(M) \triangleleft \text{id}_Y \downarrow & & \downarrow \eta_X(M \triangleleft Y) \\ \alpha_{\mathcal{M}}^-(X)(M) \triangleleft Y & \xrightarrow{(\alpha_{\mathcal{M}}^-(X))_{M,Y}} & \alpha_{\mathcal{M}}^-(X)(M \triangleleft Y) \end{array}$$



is equivalent to the monoidality of the natural isomorphism η , i.e.:

$$\begin{array}{ccc}
 \alpha_{\mathcal{N}}^+(A_1) \otimes \alpha_{\mathcal{N}}^+(A_2) & \xrightarrow{\eta_{A_1 \otimes A_2}} & \alpha_{\mathcal{N}}^-(A_1) \otimes \alpha_{\mathcal{N}}^-(A_2) \\
 (\alpha_{\mathcal{N}}^+)_{A_1, A_2} \downarrow & & \downarrow (\alpha_{\mathcal{N}}^-)_{A_1, A_2} \\
 \alpha_{\mathcal{N}}^+(A_1 \otimes A_2) & \xrightarrow{\eta_{A_1 \otimes A_2}} & \alpha_{\mathcal{N}}^-(A_1 \otimes A_2)
 \end{array}$$

Furthermore, for a \mathcal{B} -module functor $F : \mathcal{M} \rightarrow \mathcal{N}$:

• F being braided, that is $\begin{array}{c} F(m) \ x \\ \downarrow \\ \boxed{F_{m,x}} \\ \downarrow \\ \boxed{F(\swarrow)} \\ \downarrow \\ F(m \triangleleft x) \end{array} = \begin{array}{c} F(m) \ x \\ \downarrow \\ \boxed{F_{m,x}} \\ \downarrow \\ \boxed{F(\swarrow)} \\ \downarrow \\ F(m \triangleleft x) \end{array}$, is equivalent to:

$$\begin{array}{ccc}
 \mathcal{M} & \xrightarrow{F} & \mathcal{N} \\
 \downarrow \alpha_{\mathcal{M}}^+(X) & \nearrow F_{-,X}^+ & \downarrow \alpha_{\mathcal{N}}^+(X) \\
 \mathcal{M} & \xrightarrow{F} & \mathcal{N}
 \end{array}
 \xrightarrow{\eta'_X}
 \begin{array}{ccc}
 \mathcal{M} & \xrightarrow{F} & \mathcal{N} \\
 \downarrow \alpha_{\mathcal{M}}^-(X) & \nearrow F_{-,X}^- & \downarrow \alpha_{\mathcal{N}}^-(X) \\
 \mathcal{M} & \xrightarrow{F} & \mathcal{N}
 \end{array}$$

Lastly, the 2-morphisms in each category just all \mathcal{B} -module natural transformations. ■

Sketch of proof of Theorem 5.1. In light of lemma 5.2, we need only show $\mathcal{Z}(\mathbf{Mod}\text{-}\mathcal{B}) \cong \mathbf{A}(\mathcal{B})$.

(\Leftarrow) We construct a 2-functor $\mathbf{A}(\mathcal{B}) \rightarrow \mathcal{Z}(\mathbf{Mod}\text{-}\mathcal{B})$ as follows:

- Let $(\mathcal{N}, \eta : \alpha_{\mathcal{N}}^+ \xrightarrow{\sim} \alpha_{\mathcal{N}}^-)$ be an object of $\mathbf{A}(\mathcal{B})$, so $\mathcal{N} \in \mathbf{Mod}\text{-}\mathcal{B}$.
- Recall that for any $\mathcal{M} \in \mathbf{Mod}\text{-}\mathcal{B}$, there exists an algebra A in \mathcal{B} such that $\mathcal{M} \cong A - \mathbf{Mod}_{\mathcal{B}}$, the category of (left) \mathcal{A} -modules in \mathcal{B} ([6], Cor. 7.10.5). Then,

$$\mathcal{N} \boxtimes_{\mathcal{B}} \mathcal{M} \cong \mathcal{N} \boxtimes_{\mathcal{B}} (A - \mathbf{Mod}_{\mathcal{B}}) \cong \mathbf{Mod}_{\mathcal{N}} - \alpha_{\mathcal{N}}^+(A),$$

where $\alpha_{\mathcal{N}}^+(A)$ is an algebra in $\mathbf{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{N})$ and its modules in \mathcal{N} are objects in $N \in \mathcal{N}$ together with an action $\alpha_{\mathcal{N}}^+(A)(N) \rightarrow N$ satisfying the usual axioms.

Similarly, $\mathcal{M} \boxtimes_{\mathcal{B}} \mathcal{N} \cong \mathbf{Mod}_{\mathcal{N}} - \alpha_{\mathcal{N}}^-(A)$.

Hence, the isomorphism $\eta_A : \alpha_{\mathcal{N}}^+(A) \xrightarrow{\sim} \alpha_{\mathcal{N}}^-(A)$ of algebras in $\mathbf{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{N})$ yields a 2-natural \mathcal{B} -module equivalence

$$\beta_{\mathcal{M}, \mathcal{N}} : \mathcal{N} \boxtimes_{\mathcal{B}} \mathcal{M} \rightarrow \mathcal{M} \boxtimes_{\mathcal{B}} \mathcal{N}.$$

- Let $\mathcal{L} \cong A_1 - \mathbf{Mod}_{\mathcal{B}}$ and $\mathcal{M} \cong A_2 - \mathbf{Mod}_{\mathcal{B}}$. The invertible modification $R_{(\mathcal{N}|\mathcal{L}, \mathcal{M})}$ arises from the following commutative diagram of algebra isomorphisms:

$$\begin{array}{ccc} \alpha_{\mathcal{N}}^+(A_1) \otimes \alpha_{\mathcal{N}}^+(A_2) & \xrightarrow{\eta_{A_1} \otimes \eta_{A_2}} & \alpha_{\mathcal{N}}^-(A_1) \otimes \alpha_{\mathcal{N}}^-(A_2) \\ (\alpha_{\mathcal{N}}^+)_{A_1, A_2} \downarrow & & \downarrow (\alpha_{\mathcal{N}}^-)_{A_1, A_2} \\ \alpha_{\mathcal{N}}^+(A_1 \otimes A_2) & \xrightarrow{\eta_{A_1 \otimes A_2}} & \alpha_{\mathcal{N}}^-(A_1 \otimes A_2) \end{array}$$

Indeed, since $\alpha_{\mathcal{N}}^{\pm}$ is a central functor, $\alpha_{\mathcal{N}}^{\pm}(A_1) \otimes \alpha_{\mathcal{N}}^{\pm}(A_2)$ are algebras in $\mathbf{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{N})$ and $\eta_{A_1} \otimes \eta_{A_2}$ is an algebra isomorphism.

The fact that $R_{(\mathcal{N}|\text{-}, \text{-})}$ satisfies the (1, 3)-crossing and unit axioms follows from the monoidality of η .

This gives rise to a 2-functor $(\mathcal{N}, \eta) \mapsto (\mathcal{N}, \beta_{\mathcal{N}}, \text{-}, R_{(\mathcal{N}|\text{-}, \text{-})})$.

(\Rightarrow) We construct a 2-functor $\mathcal{Z}(\text{Mod-}\mathcal{B}) \rightarrow \mathbf{A}(\mathcal{B})$ as follows:

- Note that for any $X \in \text{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{B}) \cong \mathcal{B}$ and $\mathcal{N} \in \text{Mod-}\mathcal{B}$

$$\left| \begin{array}{c} \ominus \\ \mathcal{N} \end{array} \right| : \mathcal{B} \rightarrow \text{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{N}) \quad \text{given by} \quad X \mapsto \rho_{\mathcal{N}} \circ (\text{id}_{\mathcal{N}} \boxtimes_{\mathcal{B}} X) \circ \rho_{\mathcal{N}}^{-1}$$

$$\left| \begin{array}{c} \ominus \\ \mathcal{N} \end{array} \right| : \mathcal{B} \rightarrow \text{End}_{\mathcal{B}}^{\text{r.e.}}(\mathcal{N}) \quad \text{given by} \quad X \mapsto \lambda_{\mathcal{N}} \circ (X \boxtimes_{\mathcal{B}} \text{id}_{\mathcal{N}}) \circ \lambda_{\mathcal{N}}^{-1}$$

are isomorphic to $\alpha_{\mathcal{N}}^+$ and $\alpha_{\mathcal{N}}^-$ respectively.

- For an object $(\mathcal{N}, \beta_{\mathcal{N},-}, R_{(\mathcal{N}|-,-)}) \in \mathcal{Z}(\text{Mod-}\mathcal{B})$, consider:

$$\eta_X : \alpha_{\mathcal{N}}^+(X) \cong \left| \begin{array}{c} \textcircled{X} \\ \mathcal{N} \end{array} \right| = \begin{array}{c} \bullet \\ \diagdown \\ \textcircled{X} \\ \diagup \\ \mathcal{N} \end{array} \xrightarrow{\beta_{\mathcal{N},X}} \begin{array}{c} \textcircled{X} \\ \diagdown \\ \mathcal{N} \\ \diagup \\ \bullet \end{array} = \left| \begin{array}{c} \textcircled{X} \\ \mathcal{N} \end{array} \right| \cong \alpha_{\mathcal{N}}^-(X)$$

Since $\beta_{\mathcal{N},-}$ is a 2-natural transformation, $\beta_{\mathcal{N},X \circ Y} = \beta_{\mathcal{N},X} \circ \beta_{\mathcal{N},Y}$, implying that η is an isomorphism of tensor functors.

This gives rise to a 2-functor $(\mathcal{N}, \beta_{\mathcal{N},-}, R_{(\mathcal{N}|-,-)}) \mapsto (\mathcal{N}, \eta)$ which is a quasi-inverse to $\mathbf{A}(\mathcal{B}) \rightarrow \mathcal{Z}(\text{Mod-}\mathcal{B})$. □

Corollary 5.3 — $\text{BrMod-}\mathcal{B}$ may be equipped with the structure of a braided monoidal 2-category due to the equivalence $\mathcal{Z}(\Sigma\mathcal{B}) \cong \text{BrMod-}\mathcal{B}$ of monoidal 2-categories. This structure can be described explicitly (see [3], Remark 4.13).

Remark 5.1. The forgetful functor $\mathcal{Z}(\Sigma\mathcal{B}) \cong \text{BrMod-}\mathcal{B} \xrightarrow{\text{forget}} \text{Mod-}\mathcal{B} \cong \Sigma\mathcal{B}$ is fully faithful on 2-morphisms since every module natural transformations between two braided module functors is allowed.

Remark 5.2. Any \mathcal{B} -module summand of a braided \mathcal{B} -module category can be equipped with the structure of a braided \mathcal{B} -module category.

Hence, a braided \mathcal{B} -module category is indecomposable if and only if it is indecomposable as a \mathcal{B} -module category.

Warning 5.4 — For a general non-connected fusion 2-category \mathcal{A} , the canonical map $\mathcal{Z}(\mathcal{A}) \rightarrow \mathcal{A}$ is faithful on 2-morphisms but *not necessarily* full.

Exercise 5.1. Find such a fusion 2-category \mathcal{A} such that the map $\mathcal{Z}(\mathcal{A}) \rightarrow \mathcal{A}$ is *not* full.

Corollary 5.5 — $\Omega\mathcal{Z}(\Sigma\mathcal{B}) = \mathcal{Z}_{(2)}(\mathcal{B})$.

Proof of Corollary. The unit object in $\mathcal{Z}(\Sigma\mathcal{B})$ corresponds to the “rank-1 free module” $\mathcal{B} \in \text{BrMod-}\mathcal{B}$. Then recall that in example 4.5, we saw that the endomorphism category of \mathcal{B} is $\mathcal{Z}_{(2)}(\mathcal{B})$. ■

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