

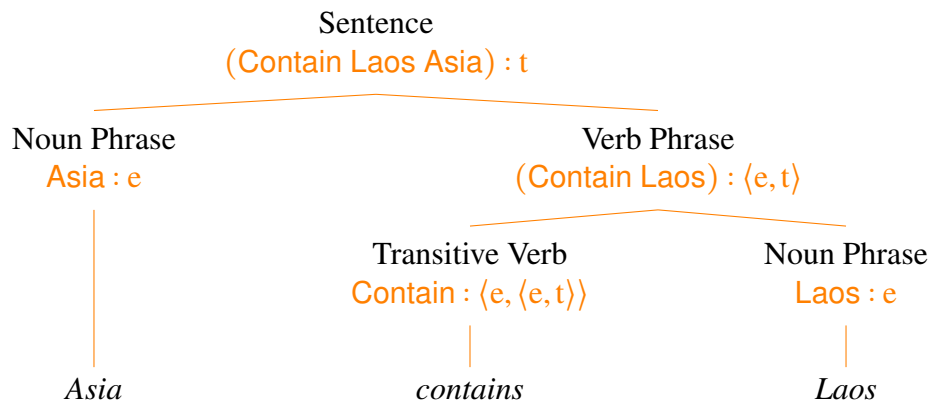
LING4400: Lecture Notes 10

Composition and Schematized Functions

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Earlier we drew translation trees that were isomorphic to syntax, which translate directly to logic:



Today we preserve this isomorphism in conjunction and negation of incomplete propositions.

10.1 Schematized conjunction and disjunction [Partee & Rooth, 1983]

We often encounter conjunctions of incomplete propositions:

- (1) a. *Etna erupts or is dormant.*
- b. (entail and entailed by 1a:) *Etna erupts or Etna is dormant.*

You might think we could just copy the subject into the conjuncts (called **conjunction reduction**).

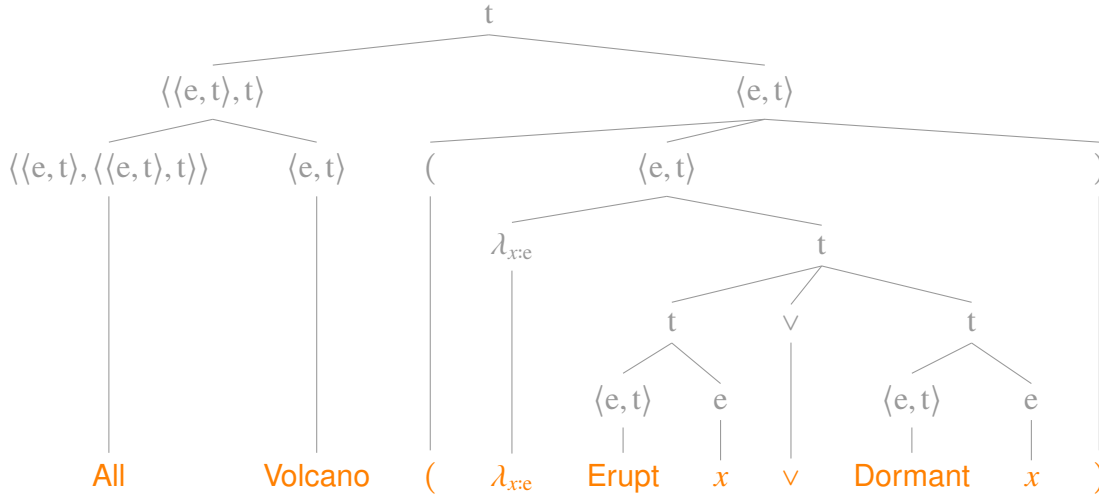
But that doesn't work with quantified subjects:

- (2) a. *All volcanoes erupt or are dormant.*
- b. (**not** entailed by 2a:) *All volcanoes erupt or all volcanoes are dormant.*

Sentence 2a is true if some volcanoes erupt and the rest are dormant; sentence 2b is not.

To model 2a, we need to get the disjunction inside the nuclear scope of *All*.

This can be done in a derivation tree:



but not in a translation tree, since no words in that sentence translate as lambda.

To do this we generalize across these types using **schemas**, defined with meta-variables γ and δ .

Specifically, we can allow functions to take any type γ_n with any number n of arguments $\delta_1, \dots, \delta_n$:

$$\begin{aligned}\gamma_0 &= t \\ \gamma_n &= \langle \delta_n, \gamma_{n-1} \rangle\end{aligned}$$

For example, using the second rule to substitute γ_2 and γ_1 and then the first rule to substitute γ_0 :

$$\gamma_2 = \langle \delta_2, \gamma_1 \rangle = \langle \delta_2, \langle \delta_1, \gamma_0 \rangle \rangle = \langle e, \langle e, t \rangle \rangle \quad (\text{if } \delta_1 = e \text{ and } \delta_2 = e)$$

(These schematized types are also called **polymorphic**, because they have several forms.)

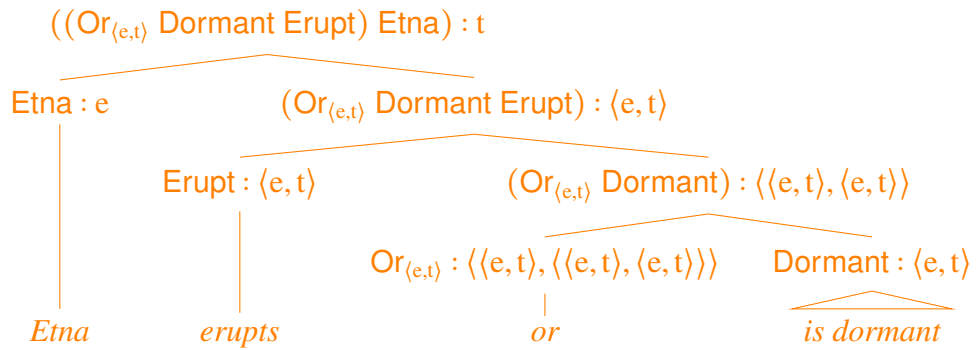
We can now define **schematized conjunction and disjunction** of type $\langle \gamma_n, \langle \gamma_n, \gamma_n \rangle \rangle$:

$$\begin{aligned}\llbracket \text{And}_{\gamma_n} \rrbracket^M &= \llbracket \lambda_{f:\gamma_n} \lambda_{g:\gamma_n} \lambda_{x_n:\delta_n} \dots \lambda_{x_1:\delta_1} f x_n \dots x_1 \wedge g x_n \dots x_1 \rrbracket^M \\ \llbracket \text{Or}_{\gamma_n} \rrbracket^M &= \llbracket \lambda_{f:\gamma_n} \lambda_{g:\gamma_n} \lambda_{x_n:\delta_n} \dots \lambda_{x_1:\delta_1} f x_n \dots x_1 \vee g x_n \dots x_1 \rrbracket^M\end{aligned}$$

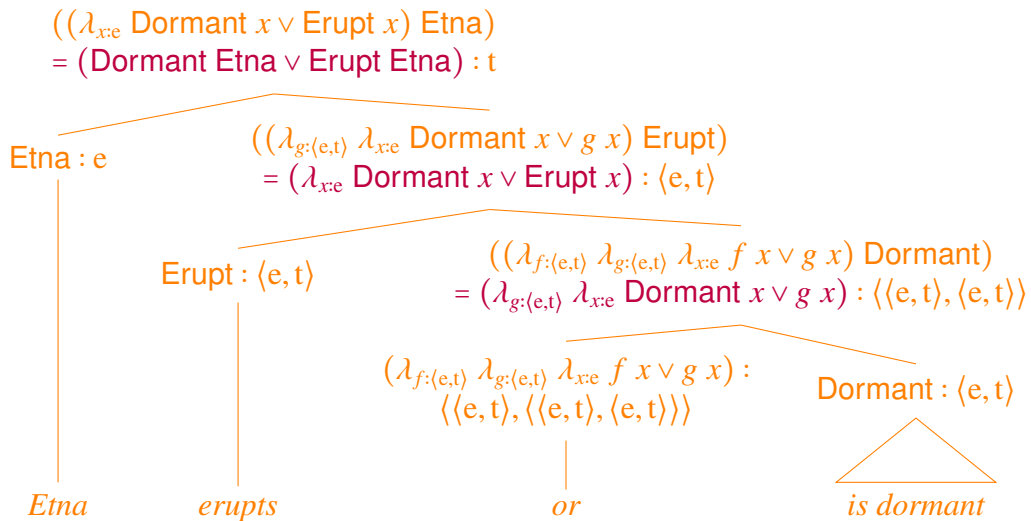
For example, if $n = 1$ and $\gamma_1 = \langle e, t \rangle$, we have a disjunction over an intransitive verb or verb phrase:

$$\llbracket \text{Or}_{\langle e, t \rangle} \rrbracket^M = \llbracket \lambda_{f:\langle e, t \rangle} \lambda_{g:\langle e, t \rangle} \lambda_{x_1:e} f x_1 \vee g x_1 \rrbracket^M$$

Here is an example translation, again isomorphic to syntax (read the translation off the top):

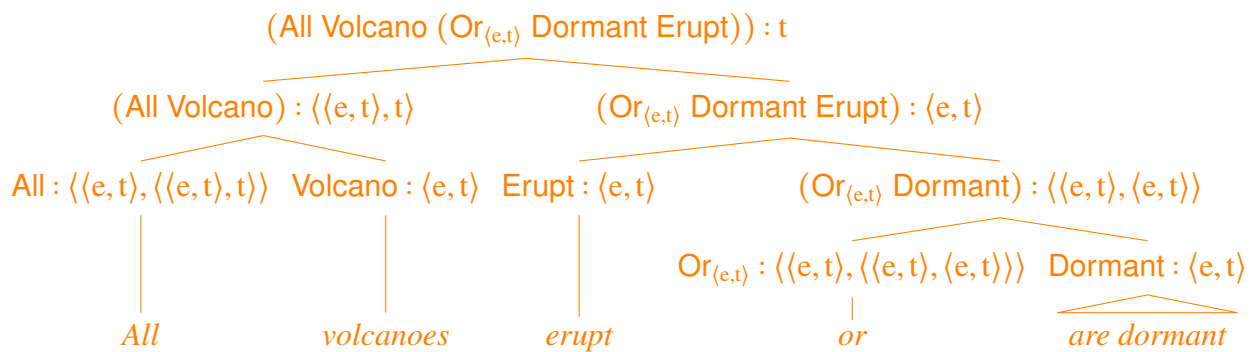


Example translation with variables (requires **beta reduction**):

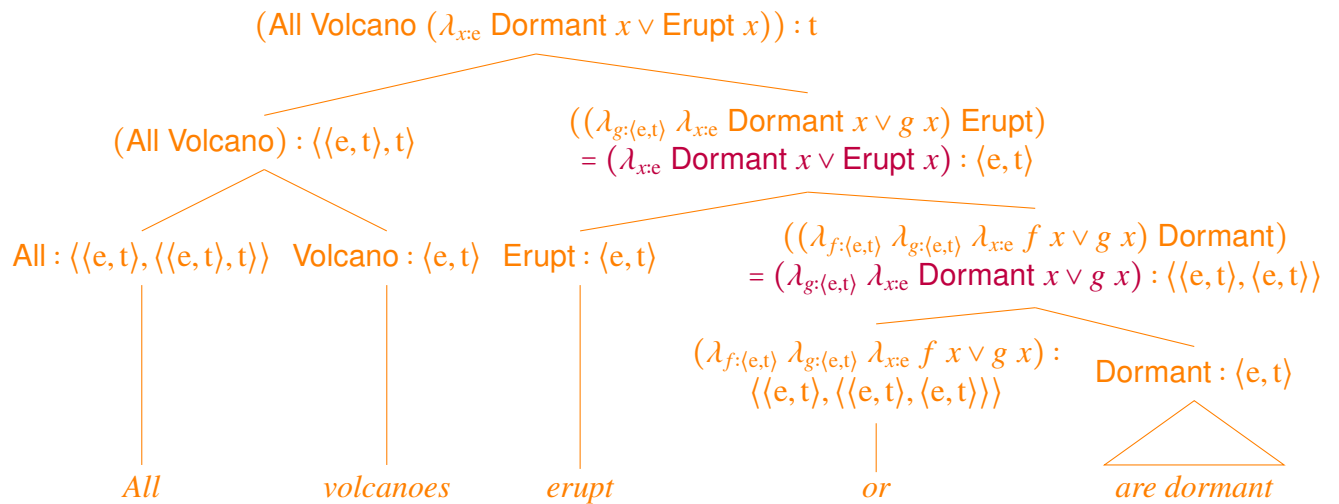


The translation $\text{Dormant Etna} \vee \text{Erupt Etna}$ is the same as for *Etna erupts or Etna is dormant*.

Here's the analysis for the quantified noun phrase:



And here's the example with variables (requires **beta reduction**):

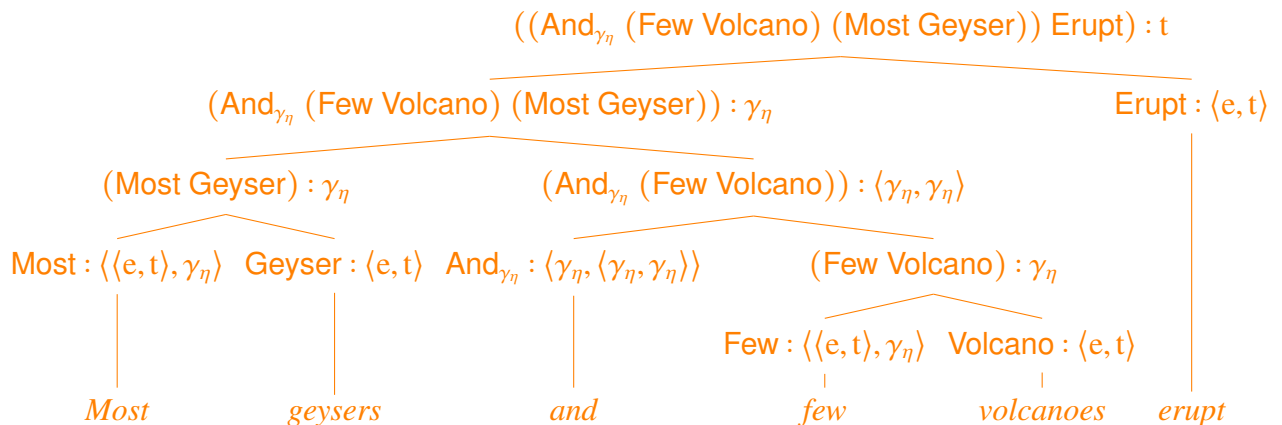


Note this is not **All Volcano Erupt ∨ All Volcano Dormant** – the disjunction is for each volcano.

Schematization also works for conjunctions of quantified noun phrases:

- (3) a. *Most geysers and few volcanoes erupt.*
- b. (entail and entailed by 3a:) *Most geysers erupt and few volcanoes erupt.*

Derivation using schematized conjunction (where $\gamma_\eta = \langle \langle e, t \rangle, t \rangle$):



Practice 10.1: schematized function

Define a schematized **And** function for conjoining transitive verbs like *peel* and *eat* of type $\langle e, \langle e, t \rangle \rangle$.

10.2 Schematized negation

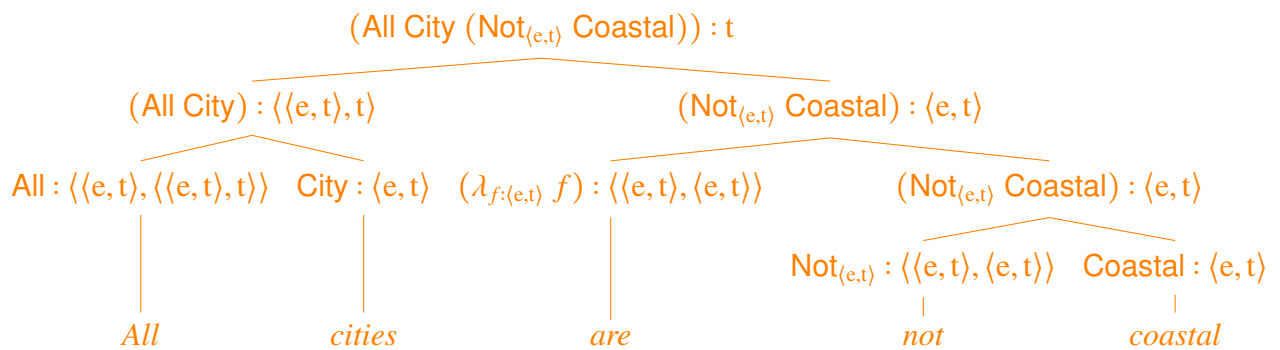
We also need schemas of type $\langle \gamma_n, \gamma_n \rangle$ in order to negate phrases which are not type **t**:

$$\llbracket \text{Not}_{\gamma_n} \rrbracket^M = \llbracket \lambda_{f:\gamma_n} \lambda_{x_n:\delta_n} \dots \lambda_{x_1:\delta_1} \neg (f x_n \dots x_1) \rrbracket^M$$

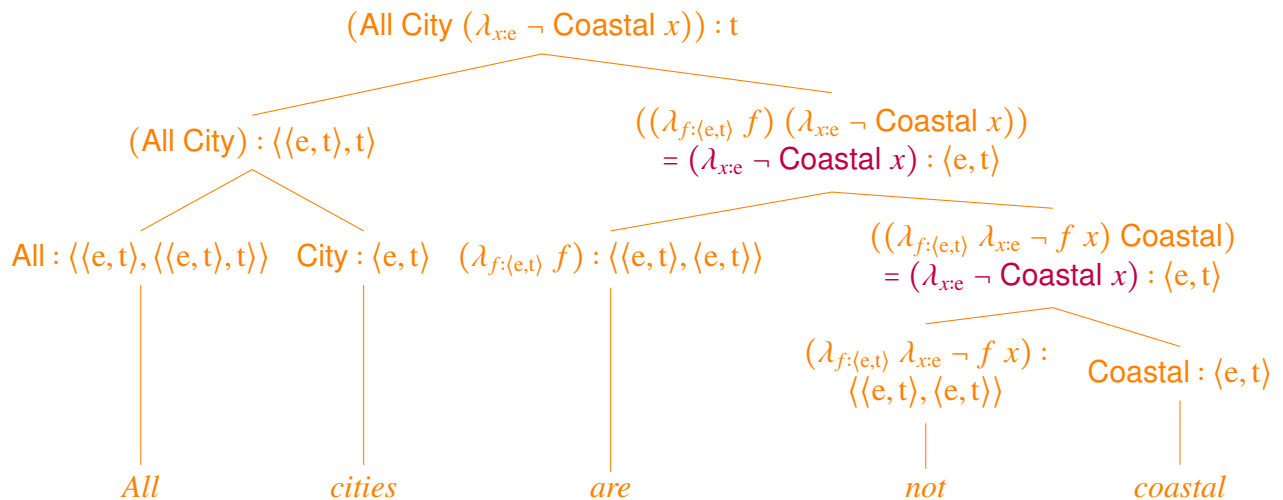
Here's an example schema definition:

$$\llbracket \text{Not}_{\langle e,t \rangle} \rrbracket^M = \llbracket \lambda_{f:\langle e,t \rangle} \lambda_{x_1:e} \neg (f x_1) \rrbracket^M$$

And here's the full translation, again isomorphic to syntax (read the translation off the top):



Example translation with variables (requires **beta reduction**):



Practice 10.2: schematized function

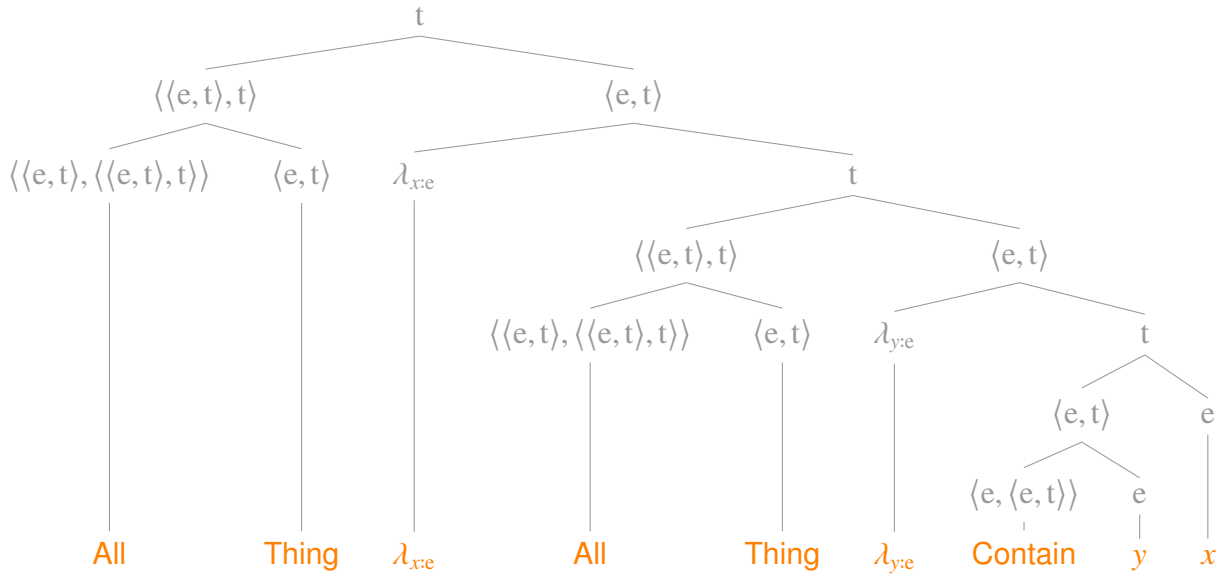
Define a schematized Not_{γ_n} function that can combine with **All**.

Practice 10.3: tree drawing

Draw a translation tree for *Not all countries are coastal* using the above function.

10.3 Schematized quantifiers

We have a similar problem with quantifiers as syntactic objects – we can do this in derivations:



But we can't mark up a syntax tree this way and read off the translation – lambdas aren't words!

To translate by making derivations isomorphic to phrase structure trees, we can't use abstraction.

Subject/object quantifiers, which take intransitive/transitive arguments, must have different types.

A schematized quantifier of type $\langle\langle e, t \rangle, \langle\langle e, \gamma_n \rangle, \gamma_n \rangle\rangle$, can then be defined for any type γ_n :

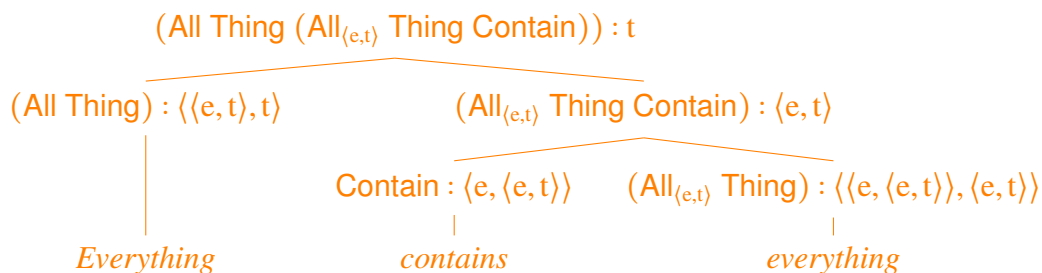
$$\llbracket \text{All}_{\gamma_n} \rrbracket^M = \llbracket \lambda_{r:\langle e, t \rangle} \lambda_{s:\langle e, \gamma_n \rangle} \lambda_{x_n:\delta_n} \dots \lambda_{x_1:\delta_1} \text{All } r (\lambda_{x_{n+1}:e} s x_{n+1} \dots x_1) \rrbracket^M$$

It takes restriction r , nuclear scope s and 'extra' arguments $x_{1..n}$ and passes the extras along to s .

For example, if $n = 1$ and $\gamma_1 = \langle e, t \rangle$ we have a quantifier over a second argument (direct object):

$$\llbracket \text{All}_{\langle e, t \rangle} \rrbracket^M = \llbracket \lambda_{r:\langle e, t \rangle} \lambda_{s:\langle e, \langle e, t \rangle \rangle} \lambda_{x_1:e} \text{All } r (\lambda_{x_2:e} s x_2 x_1) \rrbracket^M$$

Here it is in a translation, which is now isomorphic to syntax (read the translation off the top):



Practice 10.4: tree drawing

Draw a translation tree for *Everyone sent everyone everything*, using type $\langle e, \langle e, \langle e, t \rangle \rangle \rangle$ for *sent*.

Practice 10.5: translate English to logic

Translate the following into logic by drawing a tree with a logical expression at each branch:

Few people see a volcano.

References

[Partee & Rooth, 1983] Partee, B. & Rooth, M. (1983). Generalized conjunction and type ambiguity. In R. Bauerle, C. Schwarze, & A. von Stechow (Eds.), *Meaning, Use and Interpretation of Language* (pp. 361–383). Berlin: Walter de Gruyter.