

LING5702: Lecture Notes 19

Quantifier Scope

Sentences can often have multiple different readings, depending on **scope**:

- *Two people are in each booth.*

These sentences can be understood with an **in-situ** interpretation:

Two $(\lambda_x \text{ Person } x)$
 $(\lambda_x \text{ All } (\lambda_y \text{ Booth } y)$
 $(\lambda_y \text{ Some } (\lambda_e \text{ In } y \ x \ e)$
 $(\lambda_e \text{ True})))$

or a **raised** interpretation:

All $(\lambda_y \text{ Booth } y)$
 $(\lambda_y \text{ Two } (\lambda_x \text{ Person } x)$
 $(\lambda_x \text{ Some } (\lambda_e \text{ In } y \ x \ e)$
 $(\lambda_e \text{ True})))$

The raised interpretation is often non-local (can cross clause boundaries):

- *At least two CEOs said a representative was going to each country.*

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19.1 Evidence for explicit scoping [Dotlačil & Brasoveanu, 2015]

Scope seems to be explicitly calculated after each sentence (i.e. doesn't remain underspecified):

- **stimuli:** sentences presented in eye-tracking:

- (a) *A caregiver comforted a child every night. The caregivers wanted the children to...*
- (b) *A caregiver comforted a child every night. The caregivers wanted the child to...*
- (c) *A caregiver comforted a child every night. The caregiver wanted the children to...*
- (d) *A caregiver comforted a child every night. The caregiver wanted the child to...*

These analyses are eliminated at *caregiver*, but neither is the preferred in-situ analysis:

All $(\lambda_t \text{ Night } t)$	Some $(\lambda_c \text{ Child } c)$
$(\lambda_t \text{ Some } (\lambda_k \text{ Caregiver } k)$	$(\lambda_c \text{ All } (\lambda_t \text{ Night } t)$
$(\lambda_k \text{ Some } (\lambda_c \text{ Child } c)$	$(\lambda_t \text{ Some } (\lambda_k \text{ Caregiver } k)$
$(\lambda_c \text{ Comfort } t \ k \ c)))$	$(\lambda_k \text{ Comfort } t \ k \ c)))$

The preferred in-situ (first) analysis is eliminated at *children*:

Some (λ_k Caregiver k) (λ_k Some (λ_c Child c) (λ_c All (λ_t Night t) (λ_t Comfort $t k c$)))	Some (λ_k Caregiver k) (λ_k All (λ_t Night t) (λ_t Some (λ_c Child c) (λ_c Comfort $t k c$)))
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- **measure:** eye-tracking fixation durations at *children* (and spillover word).
- **results:** singular-plural (c) is slowest at *children*, suggests dynamic reanalysis there.

19.2 Computational-level scope resolution [Cooper, 1983, Keller, 1988]

We can define scope raising at the computational level, using logic.

First we generalize across types using **schemas**, defined with meta-variables γ and δ .

Specifically, we allow functions to take any type γ_n with any number n of arguments $\delta_1, \dots, \delta_n$:

$$\begin{aligned}\gamma_0 &= t \\ \gamma_n &= \langle \delta_n, \gamma_{n-1} \rangle\end{aligned}$$

Then we augment our semantics with a **store** or **context** Γ, Δ, Θ , delimited by ‘ $\vdash$ ’:

$$\underbrace{\underbrace{\underbrace{\Delta}_{\text{store}} \vdash \varphi : \underbrace{\langle \gamma_n, \gamma_{n-1} \rangle}_{\text{type}}}_{\text{expression}}}_{\text{sequent}}$$

These are called **sequents**. The store Γ, Δ, Θ of each sequent is a list of other sequents and variables.

Quantifier functions φ can now be stored, leaving variables x of type δ_n in their place:

$$\underbrace{\Delta \vdash \varphi : \langle \gamma_n, \gamma_{n-1} \rangle}_{\text{sequent}} \Rightarrow \underbrace{(\Delta \vdash \varphi : \langle \gamma_n, \gamma_{n-1} \rangle)}_{\text{stored sequent}}, \underbrace{x : \delta_n}_{\text{variable}} \vdash x : \delta_n \quad (\text{Quantifier Storage})$$

Stored functions φ are then retrieved and applied to bind the new variable x at a wider scope ψ :

$$\Gamma, \underbrace{(\Delta \vdash \varphi : \langle \gamma_n, \gamma_{n-1} \rangle)}_{\text{stored sequent}}, \underbrace{x : \delta_n}_{\text{variable}}, \Theta \vdash \psi : \gamma_{n-1} \Rightarrow \Gamma, \Delta, \Theta \vdash (\varphi (\lambda_{x:\delta_n} \psi)) : \gamma_{n-1} \quad (\text{Quantifier Retrieval})$$

Note the retrieved sequents and functions need not be retrieved in the same order they are stored!

Other rules are then **augmented with stores** that are just concatenated without being changed.

$$\Gamma \vdash \varphi : \langle \alpha, \beta \rangle \quad \Delta \vdash \chi : \alpha \Rightarrow \Gamma, \Delta \vdash (\varphi \chi) : \beta \quad (\text{Forward Function Application})$$

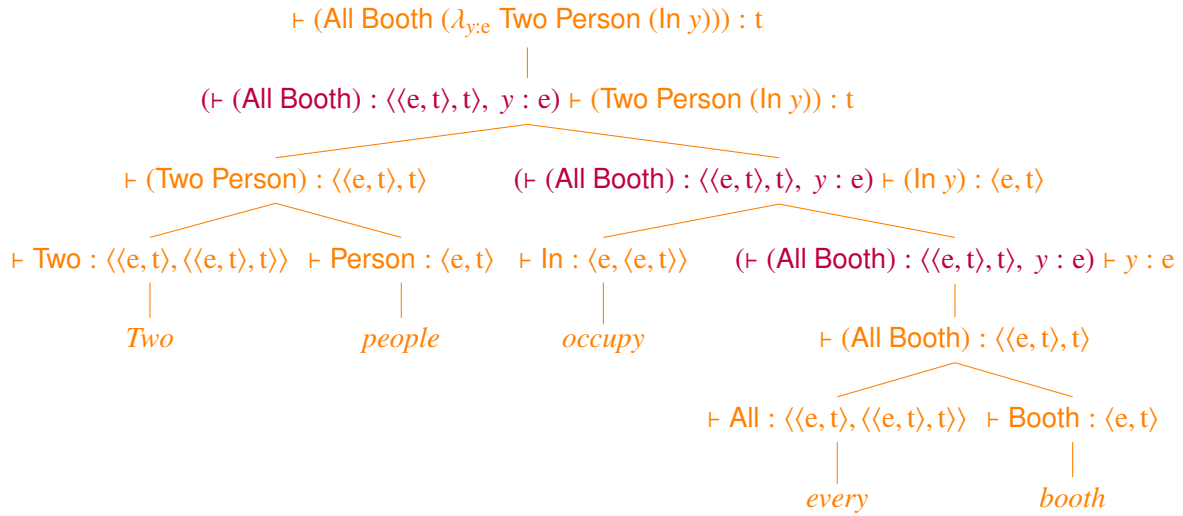
$$\Gamma \vdash \chi : \alpha \quad \Delta \vdash \varphi : \langle \alpha, \beta \rangle \Rightarrow \Gamma, \Delta \vdash (\varphi \chi) : \beta \quad (\text{Backward Function Application})$$

and similarly for Modification and Closure rules.

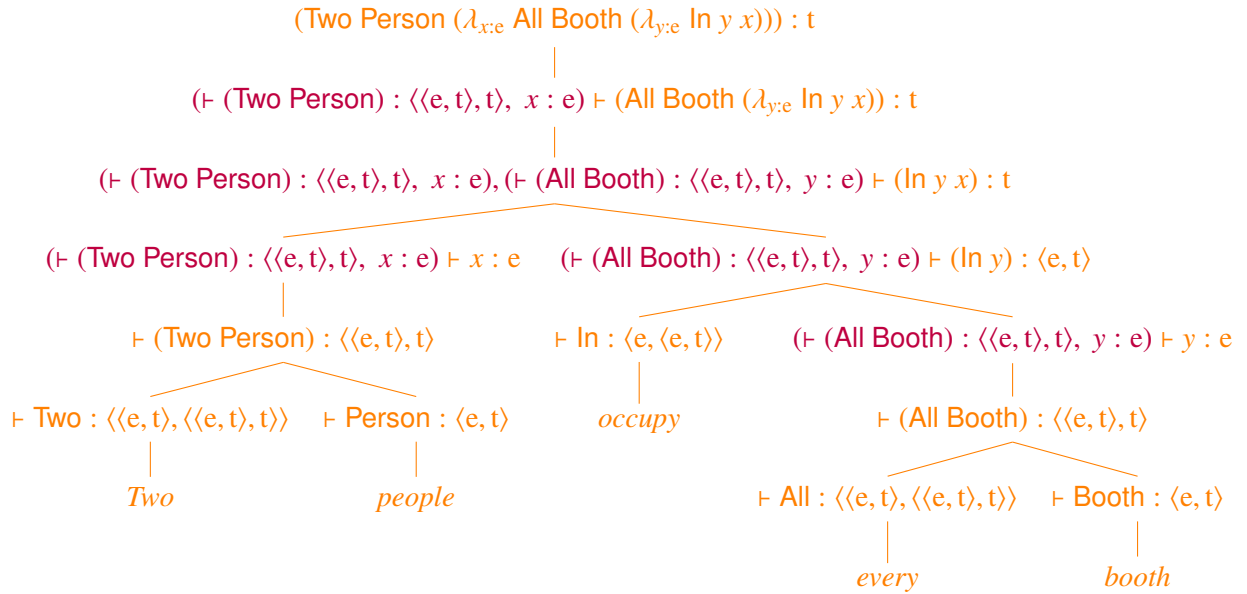
Now stored sequents can propagate up the translation and be retrieved in any order!

This approach to quantifier scope raising is called **(nested) Cooper storage**.

Now we can translate our preferred reading of *Two people occupy every booth*:



We can also translate our in-situ reading without using schematized quantifiers:



We can likewise model negation without schemas, since (recall) it can be modeled as quantification.

4. $\mathbf{E}_n \in \mathbb{R}^{V \times V}$: a matrix of associations from functions to arguments numbered by n ;

We'll also assume **inheritance** associations ('rin') from the lecture notes on sentence processing:

$$\mathbf{E}_{\text{rin}} = \mathbf{E}_1 \text{diag}(\mathbf{q}) \mathbf{E}_2^\top$$

We'll need **closure** matrices directly associating states connected by any number of associations:

$$\begin{aligned} \mathbf{E}_p &= \mathbf{I} + \sum_{n=1}^N \prod_{i=1}^n \sum_{\ell \in \{1,2,3,\dots\}} \mathbf{E}_\ell \text{diag}(\mathbf{1}-\mathbf{q}) + \text{diag}(\mathbf{1}-\mathbf{q}) \mathbf{E}_\ell^\top \\ \mathbf{E}_I &= \mathbf{I} + \sum_{n=1}^N \prod_{i=1}^n \sum_{\ell \in \{\text{cin}, \text{ein}, \text{rin}\}} \mathbf{E}_\ell + \mathbf{E}_\ell^\top \end{aligned}$$

First, initialize iteration-dependent variables:

1. $\mathbf{Q}_0 = \mathbf{0}^{V \times V}$: an initially empty matrix of immediate outscopings;
2. $\mathbf{P}_0 = \mathbf{E}_p + \mathbf{I} - \text{diag}(\mathbf{E}_p)$: a matrix of fully-connected partitions, starting with no inheritances;
3. $\mathbf{u}_0 = \sum_{v \text{ s.t. } v = \text{argmax}_v \text{diag}(\mathbf{v}) \mathbf{P}_0 \delta_v} \delta_v$: a vector of used referential states, starting with the readiest.

Then, for each iteration $i \in \{1, 2, 3, \dots\}$ such that some states remain un-used ($\mathbf{u}_{i-1} \neq \mathbf{1}$):

1. $u_i = \text{argmax}_v \underbrace{\text{diag}(\mathbf{v}) \text{diag}(\mathbf{1}-\mathbf{E}_I (\mathbf{1}-\mathbf{u}_{i-1}))}_{\text{not connected to unused}} (\mathbf{1}-\mathbf{Q}_{i-1}^\top \mathbf{1})$: get readiest used un-scoped state;
2. $\mathbf{P}_i = \mathbf{a} \mathbf{a}^\top + \underbrace{\mathbf{P}_{i-1} \text{diag}(\mathbf{1}-\mathbf{a})}_{\text{copy non-merged partitions}}$ where $\mathbf{a} = \mathbf{P}_{i-1} \mathbf{E}_I \delta_{u_i}$: merge partitions connected via u_i ;
3. $v_i = \text{argmax}_v \text{diag}(\mathbf{v}) \text{diag}(\mathbf{1}-\mathbf{u}_{i-1}) \mathbf{P}_i \delta_{u_i}$: find readiest unused state in new partition;
4. $\mathbf{Q}_i = \mathbf{Q}_{i-1} + \delta_{v_i} \delta_{u_i}^\top \mathbf{E}_I$: associate referential states in scope matrix;
5. $\mathbf{u}_i = \mathbf{u}_{i-1} + \delta_{v_i}$: add v_i as used.

Participant and scope associations define lambda calculus expressions as described earlier.

References

- [Cooper, 1983] Cooper, R. (1983). *Quantification and syntactic theory*. Dordrecht, Holland: D. Reidel.
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- [Keller, 1988] Keller, W. R. (1988). Nested co-oper storage: The proper treatment of quantifiers in ordinary noun phrases. In E. U. Reyle & E. C. Rohrer (Eds.), *Natural Language Parsing and Linguistic Theories* (pp. 432–447). D. Reidel.
- [Schuler & Wheeler, 2014] Schuler, W. & Wheeler, A. (2014). Cognitive compositional semantics using continuation dependencies. In *Third Joint Conference on Lexical and Computational Semantics (*SEM'14)*.