

# EDGEWORTH EXPANSIONS FOR NETWORK MOMENTS

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Network method of moments [20] is an important tool for nonparametric network inference. However, there has been little investigation on accurate descriptions of the sampling distributions of network moment statistics. In this paper, we present the first higher-order accurate approximation to the sampling CDF of a studentized network moment by Edgeworth expansion. In sharp contrast to classical literature on *noiseless* U-statistics, we show that the Edgeworth expansion of a network moment statistic as a *noisy* U-statistic can achieve higher-order accuracy without non-lattice or smoothness assumptions but just requiring weak regularity conditions. Behind this result is our surprising discovery that the two typically-hated factors in network analysis, namely, sparsity and edge-wise observational errors, jointly play a blessing role, contributing a crucial *self-smoothing* effect in the network moment statistic and making it analytically tractable. Our assumptions match the minimum requirements in related literature. For sparse networks, our theory shows that our empirical Edgeworth expansion and a simple normal approximation both achieve the same gradually depreciating Berry-Esseen type bound as the network becomes sparser. This result also significantly refines the best previous theoretical result.

For practitioners, our empirical Edgeworth expansion is highly accurate and computationally efficient. It is also easy to implement and convenient for parallel computing. We demonstrate the clear advantage of our method by several comprehensive simulation studies. As a byproduct, we also provide a finite-sample analysis of the network jackknife.

We showcase three applications of our results in network inference. We prove, to our knowledge, the first theoretical guarantee of higher-order accuracy for some network bootstrap schemes, and moreover, the first theoretical guidance for selecting the sub-sample size for network sub-sampling. We also derive a one-sample test and the Cornish-Fisher confidence interval for a given moment with higher-order accurate controls of confidence level and type I error, respectively.

## 1. Introduction.

1.1. *Overview.* *Network moments* are the frequencies of particular patterns, called *motifs*, that repeatedly occur in networks [102, 7, 114]. Examples include triangles, stars and wheels. They provide succinct and informative sketches of potentially very high-dimensional network population distributions. Pioneered by [20, 95], the *method of moments* for network data has become a powerful tool for frequentist nonparametric network inferences [8, 101, 131, 6, 99]. Compared to model-based network inference methods [91, 128, 94], moment method enjoys several unique advantages.

First, network moments play important roles in network modeling. They are the building blocks of the well-known exponential random graph models (ERGM) [78, 135]. More generally, under an exchangeable network assumption, the deep theory by [20] (Theorem 3) and [26] (Theorem 2.1) show that knowing all population moments can uniquely determine the network model up to weak isomorphism, despite no explicit inversion formula is yet available. From the perspective of statistical inference, evaluation of network moments

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is completely model-free, making them objective evidences for specification, validation and comparison of network models [27, 117, 125, 106]. Second, network moments can be very efficiently computed, easily allowing parallel computing. This is a crucial advantage in a big data era, where business and industry networks could contain  $10^5 \sim 10^7$  or even more nodes [43, 92] and computation efficiency becomes a substantive practicality concern. Model-fitting based network inference methods might face challenges in handling huge networks, while moment method equipped with proper sampling techniques [112, 46] will scale more comfortably (also see our comment in Section 6). Third, many network moments and their derived functionals are important structural features of great practical interest. Examples include clustering coefficient [76, 130], degree distribution [109, 122], transitivity [113], and more listed in Table A.1 in [114].

Despite the importance and raising interest in network moment method, the answer to the following core question remains under-explored:

*What is the sampling distribution of a network moment?*

For a given network motif  $R^1$ , let  $\hat{U}_n$  denote its sample relative frequency (see (2.3) for a formal definition) with expectation  $\mu_n := \mathbb{E}[\hat{U}_n]$ . Let  $\hat{S}_n^2$  be an estimator of  $\text{Var}(\hat{U}_n)$  that we shall specify later. We are mainly interested in finding the distribution of the studentized form  $\hat{T}_n := (\hat{U}_n - \mu_n)/\hat{S}_n$ . It is well-known that under the widely-studied *exchangeable network* model framework (see formal definition in Section 2.1), we have  $\hat{T}_n \xrightarrow{d} N(0, 1)$  uniformly for “not too sparse” networks [20, 17, 61], but usually,  $N(0, 1)$  only provides a rough characterization of the CDF  $F_{\hat{T}_n}$ , and one naturally yearns for a finer approximation. To this end, several network bootstrap methods have been recently proposed [20, 17, 61, 93] in an attempt to address this question. They quickly inspired many follow-up works [124, 123, 60, 37] that clearly reflect data analysts’ need of an accurate approximation method. However, compared to their empirical effectiveness, the theoretical foundation of network bootstraps remains weak. Almost all existing justifications of network bootstraps critically depend on the following type of results

$$|\hat{U}_n^* - \hat{U}_n| = o_p(n^{-1/2}), \quad \text{and} \quad |\hat{U}_n - U_n| = o_p(n^{-1/2});$$

$$\text{or similarly, } |\hat{T}_n^* - \hat{T}_n| = o_p(1), \quad \text{and} \quad |\hat{T}_n - T_n| = o_p(1);$$

where  $\hat{U}_n^*$  or  $\hat{T}_n^*$  are bootstrapped statistics and  $U_n$  or  $T_n$  are noiseless versions (see formal definitions in Section 2.2). Then the validity of network bootstraps is implied by the well-known asymptotic normality of  $U_n$  or  $T_n$  [17, 61]. However, this approach cannot show whether network bootstraps have any accuracy advantage over a simple normal approximation, especially considering the much higher computational costs of bootstraps.

In this paper, we propose the first provable *higher-order accurate* approximation to the sampling distribution of a given studentized network moment. Our paper uncovers, for the first time, that in fact the noisy  $\hat{U}_n$  and  $\hat{T}_n$  are usually more analytically tractable than the noiseless versions  $U_n$  and  $T_n$ . This enables our original analysis that sharply contrasts the common approach in existing network bootstrap literature that studies  $\hat{U}_n$  by approximately reducing it to  $U_n$ .

Now, we briefly summarize our main results by an informal theorem here. Before presenting the main results, we make a few preparatory definitions.

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<sup>1</sup>Without confusion, in this paper, we use  $R$  to represent both the motif as a subgraph pattern and its corresponding adjacency matrix representation.

92 DEFINITION 1.1 (Acyclic and cyclic motifs, see also [20, 17, 61, 93]). *A motif  $R$  is called*  
 93 *acyclic, if its edge set is a subset of an  $r$ -tree. The motif is called cyclic, if it is connected and*  
 94 *contains at least one cycle. In other words, a cyclic motif is connected but not a tree.*

95 DEFINITION 1.2. *To simplify the narration of our method's error bounds under different*  
 96 *motif shapes, especially in Table 2 and proof steps, define the following shorthand*

$$(1.1) \quad \mathcal{M}(\rho_n, n; R) := \begin{cases} (\rho_n \cdot n)^{-1} \cdot \log^{1/2} n + n^{-1} \cdot \log^{3/2} n, & \text{For acyclic } R \\ \rho_n^{-r/2} \cdot n^{-1} \cdot \log^{1/2} n + n^{-1} \cdot \log^{3/2} n, & \text{For cyclic } R \end{cases}$$

97 To simplify the narration of tail-probability control, we define the following symbol.

98 DEFINITION 1.3. *For a sequence of random variables  $\{Z_n\}$  and a deterministic se-*  
 99 *quence  $\{\alpha_n\}$ , define  $\tilde{O}_p(\cdot)$  as follows*

$$(1.2) \quad \text{We write } Z_n := \tilde{O}_p(\alpha_n), \text{ if } \mathbb{P}(|Z_n| \geq C\alpha_n) = O(n^{-1}) \text{ for some constant } C > 0.$$

101 Our " $\tilde{O}_p$ " is similar to " $\mathfrak{o}_p$ " in [96] (see the remark beneath its Lemma 2) and Assumption  
 102 (A1) in [90]. For technical reasons, in this paper, we do not need to define a  $\tilde{\mathfrak{o}}_p(\cdot)$  sign.

103 Now we are ready to present the informal statement of our main results.

104 THEOREM 1.1 (Informal statement of main results). *Assume the network is generated*  
 105 *by an exchangeable network model. Define the population Edgeworth expansion for a given*  
 106 *network moment  $R$  with  $r$  nodes and  $s$  edges as follows:*

$$(1.3) \quad G_n(x) := \Phi(x) + \frac{\varphi(x)}{\sqrt{n} \cdot \xi_1^3} \cdot \left\{ \frac{2x^2 + 1}{6} \cdot \mathbb{E}[g_1^3(X_1)] \right. \\ \left. + \frac{r-1}{2} \cdot (x^2 + 1) \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\},$$

109 where  $\Phi$  and  $\varphi$  are the CDF and PDF of  $N(0, 1)$ , respectively, and the estimable coefficients  
 110 components  $\xi_1$ ,  $\mathbb{E}[g_1^3(X_1)]$  and  $\mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$  will be defined in Section 3  
 111 and they only depend on the graphon  $f$  and the motif  $R$ . Let  $\rho_n$  denote the network sparsity  
 112 parameter. For dense networks, under the assumptions:

- 113 1.  $\rho_n^{-2s} \cdot \text{Var}(g_1(X_1)) \geq \text{constant} > 0$ ;
- 114 2. (**Dense regime**)  $\rho_n = \omega(n^{-1/2})$  for acyclic  $R$ , or  $\rho_n = \omega(n^{-1/r})$  for cyclic  $R$ ;
- 115 3. Either  $\rho_n \leq (\log n)^{-1}$ , or  $\limsup_{t \rightarrow \infty} |\mathbb{E}[e^{itg_1(X_1)/\xi_1}]| < 1$ ;

116 we have

$$(1.3) \quad \left\| F_{\hat{T}_n}(u) - G_n(u) \right\|_{\infty} = O(\mathcal{M}(\rho_n, n; R)),$$

117 where  $\|H(u)\|_{\infty} := \sup_{u \in \mathbb{R}} |H(u)|$ , and  $\mathcal{M}(\rho_n, n; R)$  (defined in (1.1)) satisfies  $\mathcal{M}(\rho_n, n; R) \ll$   
 118  $n^{-1/2}$ . Under the same conditions, the empirical Edgeworth expansion  $\hat{G}_n$  with estimated co-  
 119 efficients (see (3.14)) satisfies

$$(1.4) \quad \left\| F_{\hat{T}_n}(u) - \hat{G}_n(u) \right\|_{\infty} = \tilde{O}_p(\mathcal{M}(\rho_n, n; R)).$$

120 for a large enough absolute constant  $C$ .

121 For sparse networks, we replace condition 2 by:

122 2'. **(Sparse regime)**  $n^{-1} < \rho_n \leq n^{-1/2}$  for acyclic  $R$ , or  $n^{-2/r} < \rho_n \leq n^{-1/r}$  for cyclic  $R$ ,  
 123 The population Edgeworth expansion and a simple  $N(0, 1)$  approximation both achieve the  
 124 following Berry-Esseen bound<sup>2</sup>:

$$125 \quad (1.5) \quad \left\| F_{\hat{T}_n}(u) - G_n(u) \right\|_{\infty} = \left\| F_{\hat{T}_n}(u) - \Phi(u) \right\|_{\infty} = O(\mathcal{M}(\rho_n, n; R)) \bigwedge o(1).$$

The empirical Edgeworth expansion achieves

$$\left\| F_{\hat{T}_n}(u) - \hat{G}_n(u) \right\|_{\infty} = \tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \bigwedge o_p(1).$$

126 That is, in the sparse regime, the empirical Edgeworth expansion has the same proved error  
 127 rate bound as  $N(0, 1)$ .

128 1.2. *Our contributions.* Our contributions are three-fold. First, we establish the first  
 129 provably higher-order accurate distribution approximations for network moments (1.3) and  
 130 provide the first finite-sample error rate guarantee. The results originated from our discovery  
 131 of the surprisingly blessing roles that network noise and sparsity jointly play in this setting.  
 132 Our work reveals a new dimension to the understanding of these two components in net-  
 133 work analysis. Second, we propose a provably highly accurate and computationally efficient  
 134 empirical Edgeworth approximation (1.4) for practical use. Our method not only enjoys a sig-  
 135 nificantly improved error control than network bootstrap methods in existing literature, but  
 136 also computes much faster. Third, our results enable accurate and fast nonparametric network  
 137 inference procedures.

138 To understand the strength of our main results (1.3) and (1.4), notice that for dense net-  
 139 works (see Assumption (ii) of Lemma 3.1), we achieve *higher-order accuracy* in distribution  
 140 approximation *without non-lattice or smoothness assumption*. To our best knowledge, the  
 141 non-lattice assumption is universally required to achieve higher-order accuracy in all liter-  
 142 ature for similar settings. However, this assumption is violated by some popular network  
 143 models such as stochastic block model, arguably one of the most important and widely-used  
 144 network models. Waiving the graphon smoothness assumption makes our approach a pow-  
 145 erful tool for model-free exploratory network analysis and analyzing networks with high  
 146 complexity and irregularities, see our discussion in Section 3.4.

147 Apart from the first higher-order approximation for dense networks, for sparse networks,  
 148 we also establish a novel modified Berry-Esseen bound (1.5) for both our method and normal  
 149 approximation – this is also the sharpest result to date. These results significantly improve  
 150 over the previous best known  $o(1)$  bound in literature [20, 17, 61, 93] and fills a large blank  
 151 in the big picture. As the network sparsity  $\rho_n$  declines from  $n^{-1/2}$  towards  $n^{-1}$  for acyclic  $R$ ,  
 152 or from  $n^{-1/r}$  towards  $n^{-2/r}$  for cyclic  $R$ , our result reveals a gradually depreciating uniform  
 153 error bound. In the boundary case, where  $\rho_n = \omega(n^{-1})$  (acyclic), or  $\rho_n = \omega(n^{-2/r})$  (cyclic),  
 154 our result matches the uniform consistency result in classical literature.

155 The key insight of our method is to view the sample network moment  $\hat{U}_n$  as a *noisy*  
 156 *U-statistic*, where “noise” refers to edge-wise observational errors in the adjacency matrix  
 157  $A$ . Our analysis reveals the connection and differences between the noisy and the conven-  
 158 tional *noiseless* U-statistic settings. We discover the surprisingly blessing roles that the two  
 159 typically-hated factors, namely, *edge-wise observational errors* and *network sparsity* jointly  
 160 play in this setting, roughly summarized by the following intuitions:

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<sup>2</sup>Berry-Esseen bound for an asymptotically normal random variable  $Y_n \xrightarrow{d} N(\mu, \sigma^2)$  refers to the finite error bound  $\tau_n$  such that  $\|F_{Y_n}(u) - F_{N(\mu, \sigma^2)}(u)\|_{\infty} \leq \tau_n$ . This bound is typically discussed for CLT where  $Y_n$  is a centered and rescaled sample mean. Berry-Esseen bound for U-statistics: see [30].

- 161 1. The edge-wise errors behave like a smoother that tames potential distribution discontinuity  
 162 due to a lattice or discrete network population<sup>3</sup>;
- 163 2. Network sparsity elevates the smoothing effect of the observational error term to a suffi-  
 164 cient level, such that  $F_{\hat{T}_n}$  becomes analytically tractable.

165 At first sight, the smoothing effect of edge-wise errors is rather counter-intuitive. For in-  
 166 stance, generating a binary  $A$  from the probability matrix  $W$  is *discretizing* the edge proba-  
 167 bilities drawn from a *continuum*  $[0, 1]$  into *binary* entries. How could this eventually yield a  
 168 smoothing effect? In Section 3.1, we present two simple examples to illustrate the intuitive  
 169 reason. In our proofs, we present original analysis to carefully quantify the impact of such  
 170 smoothing effect. Our analysis techniques are very different from those in network bootstrap  
 171 papers [17, 61, 93]. Also, it seems unlikely that our assumptions can be substantially relaxed  
 172 since they match the well-known minimum conditions in related settings in [89].

173 Our empirical Edgeworth expansion (1.4) is model-free, assuming only weak regularity  
 174 conditions; has the sharpest finite-sample error bound guarantees to date; computes very fast,  
 175 much more scalable than network bootstraps; and easily permits parallel computing.

176 We showcase three applications of our main results. We present the first proof of the  
 177 higher-order accuracy of some mainstream network bootstrap techniques under certain condi-  
 178 tions, which their original proposing papers did not prove. Our results also enable rich future  
 179 works on accurate and computationally very efficient network inferences. We present two  
 180 immediate applications to testing and Cornish-Fisher type confidence interval for network  
 181 moments with explicit accuracy guarantees.

182 1.3. *Paper organization.* The rest of this paper is organized as follows. In Section 2, we  
 183 formally set up the problem and provide a detailed literature review. In Section 3, we present  
 184 our core ideas, derive the Edgeworth expansions and establish their uniform approximation  
 185 error bounds. We discuss different versions of the studentization form. We also present our  
 186 modified Berry-Esseen theorem for the sparse regime. In Section 4, we present three appli-  
 187 cations of our results: bootstrap accuracy, one-sample test, and one-sample Cornish-Fisher  
 188 confidence interval. In Section 5, we conduct three simulations to evaluate the performance  
 189 of our method from various aspects. Section 6 discusses interesting implications of our results  
 190 and future work.

191 1.4. *Big-O and small-o notation system.* In this paper, we will make frequent references  
 192 to the big-O and small-o notation system. We use the same definitions of  $O(\cdot)$ ,  $o(\cdot)$ ,  $\Omega(\cdot)$   
 193 and  $\omega(\cdot)$  as that in standard mathematical analysis, and the same  $O_p(\cdot)$  and  $o_p(\cdot)$  as that in  
 194 probability theory. For two deterministic series  $a_n$  and  $b_n$ , we write  $a_n \leq b_n$  to stand for  
 195  $a_n = O(b_n)$ ,  $n \rightarrow \infty$ ; and use  $a_n < b_n$  or  $a_n \ll b_n$  to stand for  $a_n = o(b_n)$ ,  $n \rightarrow \infty$ ; similarly  
 196 define  $\geq$ ,  $>$  and  $\gg$ .

## 197 2. Problem set up and literature review.

198 2.1. *Exchangeable networks and graphon model.* The base model of this paper is ex-  
 199 changeable network model [49, 19]. Exchangeability describes the unlabeled nature of many  
 200 networks in social, knowledge and biological contexts, where node indices do not carry  
 201 meaningful information. It is a very rich family that contains many popular models as spe-  
 202 cial cases, including the stochastic block model and its variants including degree-corrected

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<sup>3</sup>More precisely speaking, such irregularity is jointly induced by both the network population distribution and the shape of the motif, but the former is usually the determining factor.

203 stochastic block model and overlapping memberships<sup>4</sup> [75, 141, 139, 140, 3, 83, 137, 82, 57],  
 204 the configuration model [42, 103], latent space models [74, 62] and general smooth graphon  
 205 models [41, 56, 136]<sup>5</sup>. In this paper, we base our study on the following exchangeable net-  
 206 work model called *graphon model*. The framework is closely related to the Aldous-Hoover  
 207 representation for infinite matrices [5, 77]. Under a graphon model, the  $n$  nodes correspond  
 208 to latent space positions  $X_1, \dots, X_n \stackrel{\text{i.i.d.}}{\sim} \text{Uniform}[0, 1]$ . Network generation is governed by  
 209 a measurable latent graphon function  $f(\cdot, \cdot) : [0, 1]^2 \rightarrow [0, 1]$ ,  $f(x, y) = f(y, x)$  that encodes  
 210 all structures. The edge probability between nodes  $(i, j)$  is

$$(2.1) \quad W_{ij} = W_{ji} := \rho_n \cdot f(X_i, X_j); \quad 1 \leq i < j \leq n,$$

211 where the sparsity parameter  $\rho_n \in (0, 1)$  absorbs the constant factor, and we fix  $\int_{[0, 1]^2} f(u, v) dudv = \text{con-}$   
 212  $\text{stant}$ . We only observe the adjacency matrix  $A$  with conditionally independent edges:

$$(2.2) \quad A_{ij} = A_{ji} | W \sim \text{Bernoulli}(W_{ij}), \forall 1 \leq i < j \leq n.$$

213 The model defined by (2.1) and (2.2) has a well-known issue that both  $f$  and  $\{X_1, \dots, X_n\}$   
 214 are only identifiable up to equivalence classes [34]. This may pose significant challenges  
 215 for model-based network inference, especially those based on parameter estimations. On  
 216 the other hand, network moments are permutation-invariant and thus clearly immune to this  
 217 identification issue. This makes network moments attractive study objectives.

218 **2.2. Network moment statistics.** To formalize network moments, it is more convenient  
 219 to first define the sample version and then the population version. Each network moment is  
 220 indexed by the corresponding motif  $R$ . For simplicity, we focus on connected motifs. Slightly  
 221 abusing notation, here let  $R$  represent the adjacency matrix of a motif with  $r$  nodes and  $s$   
 222 edges. For any  $r$ -node sub-network  $A_{i_1, \dots, i_r}$ <sup>6</sup> of  $A$ , define

$$(2.3) \quad h(A_{i_1, \dots, i_r}) := \mathbb{1}_{[A_{i_1, \dots, i_r} \supseteq R]}^7, \quad \text{for all } 1 \leq i_1 < \dots < i_r \leq n,$$

223 Here, “ $A_{i_1, \dots, i_r} \supseteq R$ ” means there exists a permutation map  $\pi : \{1, \dots, r\} \rightarrow \{1, \dots, r\}$ , such  
 224 that  $A_{i_1, \dots, i_r} \supseteq R_\pi$ , where the “ $\supseteq$ ” is entry-wise and  $R_\pi$  is defined as  $(R_\pi)_{ij} := R_{\pi(i)\pi(j)}$ . Our  
 225 definition of  $h(A_{i_1, \dots, i_r})$  here corresponds to the “ $Q(R)$ ” defined in [20]. One can similarly  
 226 define

$$(2.4) \quad \tilde{h}(A_{i_1, \dots, i_r}) := \mathbb{1}_{[A_{i_1, \dots, i_r} \cong R]}, \quad \text{for all } 1 \leq i_1 < \dots < i_r \leq n,$$

227 where “ $A_{i_1, \dots, i_r} \cong R$ ” means there exists a permutation map  $\pi : \{1, \dots, r\} \rightarrow \{1, \dots, r\}$ , such  
 228 that  $A_{i_1, \dots, i_r} = R_\pi$ . The definition of  $\tilde{h}$  corresponds to the “ $P(R)$ ” studied in [20, 17], and  
 229 [61]. As noted by [20], each  $h$  can be explicitly expressed as a linear combination of  $\tilde{h}$  terms,  
 230 and vice versa. Therefore, they are usually treated with conceptual equivalence in literature,  
 231 and most existing papers would choose one of them to study. For technical cleanness, in this

<sup>4</sup>Here we adopt the convention of [3, 19, 1] and view community memberships and degree corrections as random samples from their respective fixed hyper-distributions. There is a distinct understanding that memberships and degree corrections are completely free unknown model parameters [59], which our study does not cover.

<sup>5</sup>Smooth graphon: we can simply think that a graphon is called “smooth” if  $f(\cdot, \cdot)$  is a smooth function. In the rigorous definition,  $f$  is smooth if  $f(\psi(\cdot), \psi(\cdot))$  is smooth under some measure-preserving map  $\psi : [0, 1] \rightarrow [0, 1]$ , see [19, 56, 136].

<sup>6</sup>We write  $A_{i_1, \dots, i_r}$  to denote the sub-matrix of  $A$  with rows and columns indexed by  $\{i_1, \dots, i_r\}$ .

<sup>7</sup>Since we consider an arbitrary but fixed  $R$  throughout this paper, without causing confusion, we drop the dependency on  $R$  in symbols such as  $h$  to simplify notation.

232 paper we focus on  $h$ . We believe our analysis also applies to  $\tilde{h}$ , but the analysis is much more  
 233 complicated and we leave it to future work. Define the *sample network moment* as

$$(2.5) \quad \hat{U}_n := \frac{1}{\binom{n}{r}} \sum_{1 \leq i_1 < \dots < i_r \leq n} h(A_{i_1, \dots, i_r}),$$

234 Then we define the *sample-population version* and *population version* of  $\hat{U}_n$  to be  $U_n :=$   
 235  $\mathbb{E}[\hat{U}_n|W]$  and  $\mu_n := \mathbb{E}[U_n] = \mathbb{E}[\hat{U}_n]$ , respectively. We refer to  $\hat{U}_n$  as the *noisy U-statistic*,  
 236 and call  $U_n := \binom{n}{r}^{-1} \sum_{1 \leq i_1 < \dots < i_r \leq n} h(W_{i_1, \dots, i_r}) = \binom{n}{r}^{-1} \sum_{1 \leq i_1 < \dots < i_r \leq n} h(X_{i_1}, \dots, X_{i_r})$ <sup>8</sup>  
 237 the conventional *noiseless U-statistic*, where we define  $h(W_{i_1, \dots, i_r}) = \mathbb{E}[h(A_{i_1, \dots, i_r})|W]$ , thus  
 238  $\mu_n = \mathbb{E}[h(X_1, \dots, X_r)]$ . Similar to the insight that studentization is key to achieve higher-  
 239 order accurate approximations in the i.i.d. setting (Section 3.5 of [129]), we study

$$\hat{T}_n := \frac{\hat{U}_n - \mu_n}{\hat{S}_n},$$

240 where  $\hat{S}_n$  will be specified later in (3.3) and (3.4). We can similarly standardize or studentize  
 241 the noiseless U-statistic  $U_n$  by  $\check{T}_n := (U_n - \mu_n)/\sigma_n$  and  $T_n := (U_n - \mu_n)/S_n$ , respectively,  
 242 where  $\sigma_n^2 := \text{Var}(U_n)$  and  $S_n^2$  is a  $\sqrt{n}$ -consistent estimator<sup>9</sup> for  $\sigma_n^2$ , for instance, a jackknife  
 243 variance estimator for the noiseless U-statistic  $U_n$ , c.f. [71, 96].

244 **2.3. Edgeworth expansions for i.i.d. data and noiseless U-statistics.** Edgeworth expansion  
 245 [51, 127] refines the central limit theorem. It is the supporting pillar in the justification of  
 246 bootstrap's higher-order accuracy, while itself is of great independent interest. In this subsection,  
 247 we review the literature on Edgeworth expansions for i.i.d. data and conventional noiseless  
 248 U-statistics, due to their close connection. Under mild conditions, the one-term Edgeworth  
 249 expansion for the sample mean of  $n$  i.i.d.  $X_1, \dots, X_n$  reads  $F_{n^{1/2}(\bar{X} - \mathbb{E}[X_1])/\sigma_{X_1}}(u) =$   
 250  $\Phi(u) - n^{-1/2} \cdot \mathbb{E}[X_1^3](u^2 - 1)\varphi(u)/(6\sigma_{X_1}^3) + O(n^{-1})$ , where  $\Phi$  and  $\varphi$  are the CDF and PDF  
 251 of  $N(0, 1)$ , respectively. Edgeworth terms of even higher orders can be derived [68] but are  
 252 not meaningful in practice unless we know a few true population moments. The minimax  
 253 rate for estimating  $\mathbb{E}[X_1^3]$  is  $O_p(n^{-1/2})$ , so  $O(n^{-1})$  is the best practical remainder bound for  
 254 an Edgeworth expansion. For further references, see [18, 115, 16, 66, 67, 10] and textbooks  
 255 [68, 47, 129].

256 The literature on Edgeworth expansions for U-statistics concentrates on the noiseless version.  
 257 In early 1980's, [30, 79, 32] established the asymptotic normality of the standardized  
 258 and the studentized U-statistics, respectively, both with  $O(n^{-1/2})$  Berry-Esseen type bounds.  
 259 Then [31, 21, 90] approximated degree-two (i.e.  $r = 2$ ) standardized U-statistics with an  
 260  $o(n^{-1})$  remainder with known population moments, and [14] established an  $O(n^{-1})$  bound  
 261 under relaxed conditions for more general symmetric statistics. Later, [71, 110] studied empirical  
 262 Edgeworth expansions (EEE) with estimated coefficients and established  $o(n^{-1/2})$   
 263 bounds. For finite populations, [11, 24, 25, 23] established the earliest results, and we will  
 264 use some of their results in our analysis of network bootstraps. An incomplete list of other  
 265 notable works on Edgeworth expansions for noiseless U-statistics with various finite moment  
 266 assumptions includes [13, 70, 80, 96, 15, 81].

<sup>8</sup>Here, without causing confusion, we slightly abused the notation of  $h(\cdot)$ , letting it take either  $W$  or  $X$  as its argument, noticing that  $W$  is determined by  $X_1, \dots, X_n$ . To elucidate  $h(W_{i_1, \dots, i_r})$ , we first explicitly re-express  $h(A_{i_1, \dots, i_r})$  as a polynomial of  $A_{i_1, \dots, i_r}$ 's edges, then replace "A" by " $W$ ". For example, with  $R = \text{triangle}$ , we have  $h(W_{123}) = W_{12}W_{13}W_{23} = \rho_n^3 f(X_1, X_2)f(X_1, X_3)f(X_2, X_3)$ . Notice that generally,  $h(W_{i_1, \dots, i_r}) \neq \mathbb{1}[W_{i_1, \dots, i_r} \supseteq R]$ .

<sup>9</sup> $\sqrt{n}$ -consistency of  $S_n^2$  means that  $\sqrt{n}(S_n^2 - \sigma_n^2) = o_p(1)$ , see [17, 93] for definition.

267 2.4. *The non-lattice condition and lattice Edgeworth expansions in the i.i.d. setting.* A  
 268 major assumption called the *non-lattice condition* is critical for achieving  $o(n^{-1/2})$  accu-  
 269 racy in Edgeworth expansions and is needed by all results in the i.i.d. setting without oracle  
 270 moment knowledge and all results for noiseless U-statistics, but this condition is clearly not  
 271 required for an  $O(n^{-1/2})$  accuracy bound<sup>10</sup>. A random variable  $X_1$  is called *lattice*, if it is  
 272 supported on  $\{a + bk : k \in \mathbb{Z}\}$  for some  $a, b \in \mathbb{R}$  where  $b \neq 0$ . General discrete distributions  
 273 are “nearly lattice”<sup>11</sup>. A distribution is essentially *non-lattice* if it contains a continuous  
 274 component. In many works, the non-lattice condition is replaced by the stronger Cramer’s  
 275 condition [45]:

$$\limsup_{t \rightarrow \infty} |\mathbb{E}[e^{itX_1}]| < 1.$$

276 For U-statistics, this condition is imposed on  $g_1(X_1) := \mathbb{E}[h(X_1, \dots, X_r)|X_1] - \mu_n$ .  
 277 Cramer’s condition can be relaxed [9, 100, 119, 120] towards a non-lattice condition, but  
 278 all existing relaxations come at the price of essentially depreciated error bounds<sup>12</sup>. There-  
 279 fore, for simplicity, in Theorems 3.1 and 4.1, we use Cramer’s condition to represent the  
 280 non-lattice setting.

281 However, in network analysis, Cramer’s condition may be a strong assumption, for the  
 282 following reasons. First, it is violated by stochastic block model, a very popular and im-  
 283 portant network model. In a block model,  $g_1(X_1)$  only depends on node 1’s community  
 284 membership, thus is discrete. Second, this condition is difficult to check in practice. Third,  
 285 some smooth models may even induce a lattice  $g_1(X_1)$  under certain motifs and a non-  
 286 lattice  $g_1(X_1)$  under a different motif. For example, under the graphon model  $f(x, y) :=$   
 287  $0.3 + 0.1 \cdot \mathbb{1}_{[x > 1/2; y > 1/2]} + 0.1 \sin(2\pi(x + y))$ ,  $g_1(X_1)$  is lattice when  $R$  is an edge, but it is  
 288 non-lattice when  $R$  is a triangle.

289 Next, we review existing treatments of Edgeworth expansion in the lattice case that will  
 290 spark the key inspiration to our work. In current literature, in the lattice case, we could ap-  
 291 proximate the CDF of an i.i.d. sample mean at higher-order accuracy, where the lattice Edge-  
 292 worth expansion would contain an order  $n^{-1/2}$  jump function; whereas to our best knowl-  
 293 edge, no analogous result exists for U-statistics. Available approaches can be categorized into  
 294 two mainstreams: (1) adding an artificial error term to the sample mean to smooth out lattice-  
 295 induced discontinuity [118, 89]; and (2) formulating the lattice version Edgeworth expansion  
 296 with a jump function [118]. The seminal work [118] adds a uniform error of bandwidth  $n^{-1/2}$ ,  
 297 and by inverting its impact on the smoothed distribution function, it explicitly formulates the  
 298 lattice Edgeworth expansion with an  $O(n^{-1})$  remainder. Another classical work [89] uses  
 299 a normal artificial error instead of uniform and shows that the Gaussian bandwidth must be  
 300  $\omega((\log n/n)^{1/2})$  and  $o(1)$  to provide sufficient smoothing effect without causing an  $\omega(n^{-1/2})$   
 301 distribution distortion. Other notable works include [132, 86, 12], in which, [132] and [86]  
 302 also formulate lattice Edgeworth expansions in the i.i.d. univariate setting, and [12] studies  
 303 Edgeworth expansions for the sample mean of i.i.d. random vectors, where some dimensions  
 304 are lattice and the others are non-lattice.

<sup>10</sup>Simply use a Berry-Esseen theorem.

<sup>11</sup>“A discrete distribution is nearly-lattice”: a discrete distribution, if not already lattice, can be viewed as a lattice distribution with diminishing periodicity.

<sup>12</sup>To our knowledge, existing results assuming only non-latticeness achieve no better than  $o(n^{-1/2})$  approximation errors. For example, [14] replaces the RHS “1” in Cramer’s condition by  $1 - q$  and assumes it holds for  $t \leq n^{1/2}$ . They obtain an error bound proportional to  $q^{-2}$ . Another example is [25]. It replaces [14]’s  $t$  range by  $t \leq \pi$  (their  $\pi$  is a variable) and obtains an error bound proportional to  $q^{-2}\pi^{-2}$ . Also see the comment beneath equation (4.7) of [110].

305 Despite the significant achievements of these treatments, latticeness remains an obstacle  
 306 in practice. The difficulties are two-fold. On one hand, if we introduce an artificial error  
 307 to smooth the distribution, it will unavoidably bring an  $\Omega(n^{-1/2})$  distortion to the original  
 308 distribution<sup>13</sup>. On the other hand, the exact formulation of a lattice Edgeworth expansion  
 309 contains an  $n^{-1/2}$  jump term. In many examples such as bootstrap, the jump locations depend  
 310 on the true population variance, laying an uncrossable  $\Omega(n^{-1/2})$  barrier for practical CDF  
 311 approximation. For more details, see page 91 of [68].

312 **3. Edgeworth expansions for network moments.** Our approach to formulate the Edge-  
 313 worth expansion can be summarized into the following progressive steps. We naturally start  
 314 with decomposing  $\hat{U}_n$  and study the stochastic variation of each term in its expansion. Based  
 315 on this understanding, we can design  $\hat{S}_n^2$  to estimate  $\text{Var}(\hat{U}_n)$ , studentize  $\hat{U}_n$  and formulate  
 316  $\hat{T}_n := (\hat{U}_n - \mu_n)/\hat{S}_n$ . But using  $\hat{S}_n$  on the denominator of  $\hat{T}_n$  introduces additional first order  
 317 (i.e.  $O(n^{-1/2})$ ) bias in the eventual distribution approximation formula and also alters  
 318 the approximately-Gaussian error term that contributes the key *self-smoothing* effect. Bear-  
 319 ing this in mind, we expand  $\hat{T}_n$  and study the impact of the terms in this decomposition.  
 320 The outcome of this part of analysis is the Edgeworth expansion formula. We then present  
 321 our main theoretical results on explicit uniform and finite-sample error bounds for population  
 322 and sample Edgeworth expansions, for dense and sparse networks, respectively. We conclude  
 323 this section by a comprehensive comparison table of our results to existing literature and fur-  
 324 ther discussions on the assumptions and results of our theory.

325 **3.1. Decomposition of the stochastic variations of  $\hat{U}_n$  and design of the variance estimator**  
 326  $\hat{S}_n^2$ . The starting point of all analysis is the decomposition of  $\hat{U}_n$ . This would allow us to  
 327 design a variance estimator of  $\hat{U}_n$  for studentization. The studentized form,  $\hat{T}_n$ , has a related  
 328 but different decomposition, which will be formulated and analyzed next in Section 3.2. Now  
 329 let us inspect  $\hat{U}_n$ .

330 The stochastic variations in  $\hat{U}_n - \mu_n = (U_n - \mu_n) + (\hat{U}_n - U_n)$  stem from two sources:  
 331 (1) the randomness in  $U_n - \mu_n$  due to  $W$  and ultimately  $X_1, \dots, X_n$ ; and (2) the ran-  
 332 domness in  $\hat{U}_n - U_n$  due to  $A|W$ , the edge-wise observational errors. In  $\text{Var}(\hat{U}_n) =$   
 333  $\mathbb{E}[\text{Var}(\hat{U}_n|W)] + \text{Var}(\mathbb{E}[\hat{U}_n|W])$ , by Lemma 3.1, we observe  $\text{Var}(\hat{U}_n|W) \asymp \rho_n^{2s-1} \cdot n^{-2}$   
 334 and  $\text{Var}(\mathbb{E}[\hat{U}_n|W]) = \text{Var}(U_n) \asymp \rho_n^{2s} \cdot n^{-1}$ . We shall universally assume  $\rho_n \cdot n \rightarrow \infty$ , so  
 335  $\sigma_n^2 = \text{Var}(U_n) = \text{Var}(\mathbb{E}[\hat{U}_n|W])$  dominates. Therefore, our design of the variance estimator  
 336  $\hat{S}_n^2$  for  $\text{Var}(\hat{U}_n)$  should align with the formulation of  $\text{Var}(U_n - \mu_n)$ .

337 Now we inspect the main term  $U_n - \mu_n$ . It is a conventional noiseless U-statistic that  
 338 admits the well-known Hoeffding's decomposition [73]:

$$339 \quad (3.1) \quad U_n - \mu_n = \underbrace{\frac{r}{n} \sum_{i=1}^n g_1(X_i)}_{\text{Linear part}} + \underbrace{\frac{r(r-1)}{n(n-1)} \sum_{1 \leq i < j \leq n} g_2(X_i, X_j)}_{\text{Quadratic part}} + \underbrace{\tilde{O}_p(\rho_n^s \cdot n^{-3/2} \log^{3/2} n)}_{\text{Higher-degree part}}$$

340 where  $g_1, \dots, g_r$  are defined as follows. To avoid complicated subscripts, without confu-  
 341 sion we define  $g_k$ 's for special indexes  $(i_1, \dots, i_r) = (1, \dots, r)$ . For indexes 1 and  $k \in$   
 342  $\{2, \dots, r-1\}$  (only when  $r \geq 3$ ) and  $r$ , define  $g_1(x_1) := \mathbb{E}[h(X_1, \dots, X_r)|X_1 = x_1] - \mu_n$ ,  
 343  $g_k(x_1, \dots, x_k) := \mathbb{E}[h(X_1, \dots, X_r)|X_1 = x_1, \dots, X_k = x_k] - \mu_n - \sum_{k'=1}^{k-1} \sum_{1 \leq i_1 < \dots < i_{k'} \leq r} g_{k'}(x_{i_1}, \dots, x_{i_{k'}})$   
 344 for  $2 \leq k \leq r-1$  and  $g_r(x_1, \dots, x_r) := h(x_1, \dots, x_r) - \mu_n$ . From classical literature,

<sup>13</sup>To see this, simply notice that the original distribution contains  $n^{-1/2}$  jumps, but the smoothed distribution does not, so an  $o(n^{-1/2})$  approximation error is impossible [21].

we know that  $\mathbb{E}[g_k(X_{i_1}, \dots, X_{i_k}) | \{X_i : i \in \mathcal{I}_k \subset \{i_1, \dots, i_k\}\}] = 0$ , where the strict subset  $\mathcal{I}_k$  could be  $\emptyset$ , and  $\text{Cov}(g_k(X_{i_1}, \dots, X_{i_k}), g_\ell(X_{j_1}, \dots, X_{j_\ell})) = 0$  unless  $k = \ell$  and  $\{i_1, \dots, i_k\} = \{j_1, \dots, j_\ell\}$ . Consequently, the linear part in the Hoeffding's decomposition makes dominating contribution to  $\text{Var}(U_n - \mu_n)$ <sup>14</sup>. Define

$$(3.2) \quad \xi_1^2 := \text{Var}(g_1(X_1)).$$

Now we are ready to design  $\widehat{S}_n$  and thus can fully specify  $\widehat{T}_n = (\widehat{U}_n - \mu_n) / \widehat{S}_n$ . There are two main choices of  $\widehat{S}_n$ . The conventional choice for studentizing noiseless U-statistics [32, 71, 110] uses the jackknife estimator

$$(3.3) \quad n \cdot \widehat{S}_{n;\text{jackknife}}^2 := (n-1) \sum_{i=1}^n \left( \widehat{U}_n^{(-i)} - \widehat{U}_n \right)^2,$$

where  $\widehat{U}_n^{(-i)}$  is  $\widehat{U}_n$  calculated on the induced sub-network of  $A$  with node  $i$  removed. Despite conceptual straightforwardness, the jackknife estimator unnecessarily complicates analysis. In this paper, we propose an estimator with a simpler formulation. In  $\text{Var}(\widehat{U}_n) = \sigma_n^2 + O(\rho_n^{2s-1}n^{-2}) = r^2\xi_1^2/n + O(\rho_n^{2s-1}n^{-2})$ , replace  $\xi_1$  by its moment estimator. Specifically, recall that  $\xi_1^2 = \text{Var}(g_1(X_1)) = \mathbb{E}[(\mathbb{E}[h(X_1, \dots, X_n) | X_1] - \mu_n)^2]$ . Replacing  $\mathbb{E}[h(X_1, \dots, X_n) | X_1]$  and  $\mu_n$  by their estimators based on observable data, we can design  $\widehat{S}_n$  as follows

$$(3.4) \quad n \cdot \widehat{S}_n^2 := \frac{r^2}{n} \sum_{i=1}^n \underbrace{\left\{ \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(A_{i, i_1, \dots, i_{r-1}}) - \widehat{U}_n \right\}^2}_{\text{Estimates } \xi_1^2 = \text{Var}(g_1(X_1))}.$$

We will show in Theorem 3.3 that the  $|\widehat{S}_n^2 - \widehat{S}_{n;\text{jackknife}}^2|$  is ignorable, but our estimator  $\widehat{S}_n$  is computationally much more efficient than the jackknife estimator. See our discussion right following Theorem 3.3.

*3.2. Expansion of  $\widehat{T}_n$  and self-smoothing phenomenon.* The studentization  $\widehat{T}_n$  can be expanded using a similar method to our study of  $\widehat{U}_n$ , but certain into a very different expression. The analysis in Section 3.1 already gives us a good understanding of the expansion of  $\widehat{T}_n$ 's numerator, namely, recall that

$$(3.5) \quad \widehat{U}_n - \mu_n = \frac{r}{n} \sum_{i=1}^n g_1(X_i) + \frac{r(r-1)}{n(n-1)} \sum_{1 \leq i < j \leq n} g_2(X_i, X_j) + (\widehat{U}_n - U_n) + \text{remainder}$$

where we shall prove that the remainder terms contributed by  $g_k, k \geq 3$  are dominated by  $\widehat{U}_n - U_n$ . Now, to handle  $\widehat{T}_n$ 's denominator, we follow the method in Maesono [96] and re-express  $\widehat{T}_n$  as:

$$(3.6) \quad \widehat{T}_n = \frac{\widehat{U}_n - \mu_n}{\widehat{S}_n} = \frac{\widehat{U}_n - \mu_n}{\sigma_n} \cdot \left\{ 1 + \frac{\widehat{S}_n^2 - \sigma_n^2}{\sigma_n^2} \right\}^{-1/2}$$

<sup>14</sup>Hoeffding's decomposition reveals that the asymptotic behavior of the noiseless U-statistic  $U_n$  is largely determined by the linear part and bears some similarity to the i.i.d. case. But we should also notice that the quadratic part, i.e.  $g_2$  terms, plays a non-ignorable role in the Edgeworth expansion of  $U_n$ . For more details, see [21, 71, 96, 110]

369 and use Taylor expansion  $(1+x)^{-1/2} \approx 1 - x/2 + O(x^2)$  with  $x := (\widehat{S}_n^2 - \sigma_n^2)/\sigma_n^2 =$   
 370  $\widetilde{O}_p(n^{-1/2})$ . In fact, just like our earlier decomposition of  $\widehat{U}_n - \mu_n$  into two parts that repre-  
 371 sent the random variations originated from  $W$  (or  $X_1, \dots, X_n$ ) and  $A|W$ , respectively; here,  
 372 it is also technically beneficial to do the same for  $\widehat{S}_n^2 - \sigma_n^2$ . Define an auxiliary intermediate  
 373 term  $\widehat{\sigma}_n^2$  to insert in between  $\widehat{S}_n^2$  and  $\sigma_n^2$ :

$$374 \quad n \cdot \widehat{\sigma}_n^2 := \frac{r^2}{n} \sum_{i=1}^n \left\{ \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(W_{i, i_1, \dots, i_{r-1}}) - U_n \right\}^2.$$

375 and also define the following convenience shorthand

$$376 \quad (3.7) \quad U_n^\# := \frac{1}{\sqrt{n} \cdot \xi_1} \sum_{i=1}^n g_1(X_i), \quad \Delta_n := \frac{r-1}{\sqrt{n}(n-1)\xi_1} \sum_{1 \leq i < j \leq n} g_2(X_i, X_j),$$

$$377 \quad \widehat{\Delta}_n := (\widehat{U}_n - U_n)/\sigma_n, \quad \delta_n := (\widehat{\sigma}_n^2 - \sigma_n^2)/\sigma_n^2, \quad \text{and} \quad \widehat{\delta}_n := (\widehat{S}_n^2 - \widehat{\sigma}_n^2)/\sigma_n^2,$$

378 Recall that in Section 3.1 we observed that  $\sigma_n^2 := \text{Var}(U_n) \asymp r^2 \xi_1^2/n$ . We now obtain the key  
 379 expansion of  $\widehat{T}_n$  as follows:

$$380 \quad \widehat{T}_n = \left( U_n^\# + \Delta_n + \widehat{\Delta}_n + \widetilde{O}_p(n^{-1} \log^{3/2} n) \right) \cdot \left( 1 + \widehat{\delta}_n + \delta_n \right)^{-1/2} \\
 381 \quad (3.8) \quad = \widetilde{T}_n + \check{\Delta}_n + \text{Remainder},$$

382 in which we define

$$383 \quad (3.9) \quad \widetilde{T}_n := U_n^\# + \Delta_n - \frac{1}{2} U_n^\# \cdot \delta_n,$$

$$384 \quad (3.10) \quad \check{\Delta}_n := \binom{n}{2}^{-1} \sum_{1 \leq i < j \leq n} \widehat{\Theta}_{ij} \cdot \eta_{ij},$$

385 where we define  $\eta_{ij} := A_{ij} - W_{ij}$ , and the formal definition of  $\widehat{\Theta}_{ij}$  is lengthy and sunk  
 386 to Supplemental Material (see (8.19)). The gist is that  $\widehat{\Theta}_{ij}$  is a function of  $W$  (thus all its  
 387 randomness comes from  $X_1, \dots, X_n$ ) and does *not* depend on the conditional randomness in  
 388  $A|W$ , and also that  $\widehat{\Theta}_{ij} \asymp \rho_n^{-1} \cdot n^{1/2}$ . The term  $\check{\Delta}_n$  encodes the ‘‘linear part’’ (linear in  $\eta_{ij}$ ’s)  
 389 of  $\widehat{\Delta}_n$  (see Lemma 3.1-(c)). The remainder in (3.8) consists of the remainder terms from the  
 390 two expansions of  $U_n - \mu_n$  and  $\widehat{U}_n - U_n$ , respectively. We will show that the remainder is  
 391  $\widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$ , where we recall the definition of  $\widetilde{O}_p$  from Section 1.4.

392 To give readers a quick preview of the roles of the main constituent terms in the expansion  
 393 of  $\widehat{T}_n$ , we present a summary table, see Table 1. The full quantitative justification of its  
 394 contents will be provided soon in Lemma 3.1. Notice that despite smoother  $\check{\Delta}_n$  is  $\Omega(n^{-1/2})$ ,  
 395 it does *not* distort any smooth order- $n^{-1/2}$  term in the Edgeworth expansion formula. Similar  
 396 phenomenon is observed in the i.i.d. setting, see [118] (equation (2.8)) and [89] (Section 2.2).

397 Our decomposition (3.8) is a renaissance of the spirits of [118] and [89], but with the  
 398 following crucial conceptual distinctions. First and most important, the error term  $\check{\Delta}_n$  in our  
 399 formula is *not* artificial, but a natural constituent component of  $\widehat{T}_n$ . Therefore, the smoother  
 400 does *not* distort the objective distribution, that is,  $\widehat{T}_n$  is *self-smoothed*. The second distinction  
 401 lies in the bandwidth of the smoothing error term. Since the smoothing error terms in [118]  
 402 and [89] are artificial, the user is at the freedom to choose these bandwidths. In our setting,  
 403 the bandwidth of the smoothing term  $(\rho_n \cdot n)^{-1/2}$  is not managed by the user, but governed  
 404 by the network sparsity. Therefore, when Cramer’s condition fails, we make the very mild

TABLE 1

Summary of the main components in  $\hat{T}_n$ 

Component	Order of std. dev.	Impacts Edgeworth formula	Smoothing effect
$U_n^\#$	1	Yes	No
$\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n$	$n^{-1/2}$	Yes	No
$\check{\Delta}_n$	$(\rho_n \cdot n)^{-1/2}$	No	Yes
Remainder	$\tilde{O}_p(\mathcal{M}(\rho_n, n; R))$	No	No

405 sparsity assumption that  $\rho_n = O((\log n)^{-1})$  to ensure enough smoothing effect. This echoes  
 406 the lower bound on Gaussian bandwidth in [89]. This upper bound can be easily enforced  
 407 by a data pre-processing step. See our discussion in Section 6. We also need  $\rho_n$  to be lower  
 408 bounded to effectively bound the remainder term, see Lemma 3.1-(b). Third, our error term  
 409  $\check{\Delta}_n$  is dependent on  $\tilde{T}_n$  through  $W$ . Last, the proof technique of [118] is inapplicable to our  
 410 setting due to the quadratic part ( $g_2(X_i, X_j)$  terms) in  $\tilde{T}_n$ ; and [89] obtains an  $o(n^{-1/2})$  error  
 411 bound<sup>15</sup>, while we aim at stronger results under a more complicated U-statistic setting with  
 412 degree-two terms. In our proofs, we carefully manage these challenges with original analysis.

413 A key difference between our *noisy* U-statistic setting and the conventional *noiseless* set-  
 414 ting is carried by the  $\check{\Delta}_n$  term, which is unique to network data. Prior to our paper, the typical  
 415 treatment in network bootstrap literature is to simply bound and ignore this component, such  
 416 as Lemma 7 in [61]. In sharp contrast, by carefully quantifying the impact of  $\check{\Delta}_n$ , we shall  
 417 reveal its key smoothing effect by a refined analysis. Therefore, before advancing to the state-  
 418 ment of our main lemma, we present two concrete examples to give the general audience an  
 419 intuitive impression of the asymptotic orders of each constituent term in (3.10). For sim-  
 420 plicity of illustration, in these examples, we would *standardize*  $\hat{U}_n$  using its *true* variance  
 421  $\sigma_n^2$ , rather than the estimator  $\hat{S}_n^2$ . The impact of this simplification is that the expansion of  
 422 the standardization would not have the  $-(1/2)U_n^\# \cdot \delta_n$  term, and an altered  $\hat{\Theta}_{ij}$  at the same  
 423 asymptotic order as the original  $\hat{\Theta}_{ij}$ , and a different remainder term; but all these differences  
 424 are non-essential for demonstrating our core ideas. For the moment, let us bear in mind that  
 425  $\sigma_n \asymp \hat{S}_n \asymp \rho_n^s \cdot n^{-1/2}$  by Lemma 3.1. We first study the simplest motif  $R = \text{Edge}$ .

EXAMPLE 3.1. *Let  $R$  be an edge with  $r = 2$  and  $s = 1$ , and  $\hat{U}_n$  is simply the sample edge density  $\hat{\rho}_n := \bar{A}$ . By definition, all  $h(A_{i_1, i_2}) - h(W_{i_1, i_2})$  terms are mutually conditionally independent given  $W$ . Then the asymptotic behavior of the self-smoother term is*

$$\frac{\hat{U}_n - U_n}{\sigma_n} \Big| W \xrightarrow{d} N \left( 0, \sigma_{\frac{\hat{U}_n - U_n}{\sigma_n} | W}^2 \asymp (\rho_n \cdot n)^{-1} \right)$$

426 *at a uniform  $O(\rho_n^{-1/2} \cdot n^{-1})$  Berry-Esseen CDF approximation error rate.*

427 The next example shows that the key insight of Example 3.1 also applies to general motifs.

428 EXAMPLE 3.2. *Let  $R$  be a triangular motif with  $r = 3, s = 3$ , and  $\hat{U}_n$  is the empirical*  
 429 *triangle frequency. We can decompose  $\hat{U}_n - U_n$  as follows:*

$$430 \quad \frac{\hat{U}_n - U_n}{\sigma_n} = \frac{1}{\binom{n}{3}} \sum_{1 \leq i_1 < i_2 < i_3 \leq n} \frac{\{h(A_{i_1, i_2, i_3}) - h(W_{i_1, i_2, i_3})\}}{\sigma_n}$$

<sup>15</sup>The  $o(n^{-1/2})$  error bound in [89] holds on some  $\mathfrak{B} \subset \mathbb{R}$  with ‘‘diminishing boundary’’, while our error bounds hold on the entire  $\mathbb{R}$ .

$$\begin{aligned}
 &= \frac{1}{\binom{n}{3}} \sum_{1 \leq i_1 < i_2 < i_3 \leq n} \frac{(W_{i_1 i_2} + \eta_{i_1 i_2})(W_{i_1 i_3} + \eta_{i_1 i_3})(W_{i_2 i_3} + \eta_{i_2 i_3}) - W_{i_1 i_2} W_{i_1 i_3} W_{i_2 i_3}}{\sigma_n} \\
 &= \frac{1}{\binom{n}{3}} \left\{ \sum_{\substack{1 \leq i_1 < i_2 \leq n \\ 1 \leq i_3 \leq n \\ i_3 \neq i_1, i_2}} \frac{W_{i_1 i_3} W_{i_2 i_3} \eta_{i_1 i_2} + W_{i_1 i_2} \eta_{i_1 i_3} \eta_{i_2 i_3}}{\sigma_n} + \sum_{1 \leq i_1 < i_2 < i_3 \leq n} \frac{\eta_{i_1 i_2} \eta_{i_1 i_3} \eta_{i_2 i_3}}{\sigma_n} \right\} \\
 &= \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \underbrace{\left( \frac{3 \sum_{\substack{1 \leq k \leq n \\ k \neq i, j}} W_{ik} W_{jk}}{(n-2)\sigma_n} \right) \eta_{ij}}_{\text{Linear part}} + \frac{1}{\binom{n}{3}} \sum_{\substack{1 \leq i < j \leq n \\ 1 \leq k \leq n \\ k \neq i, j}} \frac{W_{ij}}{\sigma_n} \eta_{ik} \eta_{jk} \\
 &\quad + \frac{1}{\binom{n}{3}} \sum_{1 \leq i < j < k \leq n} \frac{1}{\sigma_n} \eta_{ij} \eta_{ik} \eta_{jk} \\
 &\quad \underbrace{\hspace{10em}}_{\text{Cubic part}}
 \end{aligned}$$

where recall that we define  $\eta_{ij} := A_{ij} - W_{ij}$ . Recall that we are conditioning on  $W$ , so  $\sigma_n \asymp \rho_n^s \cdot n^{-1/2}$  is treated as a constant. The linear part is  $\asymp \rho_n^{-1/2} \cdot n^{-1/2}$ , the quadratic part is  $\tilde{O}_p(\rho_n^{-1} \cdot n^{-1} \log^{1/2} n)$  and the cubic part is  $\tilde{O}_p(\rho_n^{-3/2} \cdot n^{-1} \log^{1/2} n)$ . We make two observations. First, the linear part in this example has the same asymptotic order as the linear part in Example 3.1. This is not a coincidence and will be formalized by Lemma 3.1-(b). In other words, regardless of the shape of  $R$ , the linear part in such decomposition always provides smoothing effect at the same magnitude. Second, different from Example 3.1, we now have higher-degree remainder consisting of products of quadratic and cubic  $\eta$  terms. The linear part nicely always dominates the quadratic part; but it only dominates the cubic part when  $\rho_n = \omega(n^{-1/2} \log^{1/2} n)$ .

For readers' convenience, we now link the terms in the two examples to items in Table 1. The entire  $(\hat{U}_n - U_n)/\sigma_n$  in Example 3.1 and the linear part of the expansion in Example 3.2 both map to  $\Delta_n$  in Table 1; and the quadratic and cubic parts of the expansion in Example 3.2 correspond to the remainder part in Table 1.

Readers who are familiar with the martingale CLT (c.f. [69]) see immediately that the cubic part in Example 3.2 is also asymptotically normal and naturally question why our study would stick to  $\rho_n$  regimes such that this term is ignorable. In other words, when the network is very sparse that the cubic part dominates the linear part, can the asymptotic normality of the former take over the role of self-smoother? The reason why the cubic part is much more challenging to characterize than the linear part lies in its very slow convergence to its limiting normal distribution. In Example 3.2, the CDF of the linear part converges to its limiting distribution at a uniform rate of  $O(\rho_n^{-1/2} \cdot n^{-1})$  (See (3.12) in our Lemma 3.1-(b)). In sharp contrast, the convergence rate of the cubic part as a martingale is much slower. The reported uniform convergence rate for martingale CLT across various different settings in literature are all significantly slower than  $n^{-1/2}$ , see [72, 64, 107, 29] and so on. This is not surprising considering the lack of independence between summands in the scenarios that martingale CLT addresses. Our discussion here does not disprove the possibility that a sharper analysis might show that the higher-degree  $\eta$ -product terms in fact can serve as the self-smoother, but required analysis might be difficult. Considering the already existing complexity of this paper, we simply control the stochastic magnitude of the cubic part in Example 3.2.

On the other hand, however, the asymptotic normality of the cubic part provides a unrigorous but helpful intuitive understanding of the  $\log^{1/2} n$  factor in the first term of our error

bound (1.2). If we roughly treat  $Z_{\text{cubic}} :=$  the cubic part in Example 3.2 as normal, then  $\mathbb{P}(|Z_{\text{cubic}}| > C(\text{Var}(Z_{\text{cubic}})^{1/2} \cdot \log^{1/2} n)) = O(n^{-1})$  for a large enough constant  $C$ . The  $\log^{3/2} n$  factor in the second term of (1.2) comes from a different source, namely, the tail probability control of  $g_k : k \geq 3$  terms in the Hoeffding's decomposition of  $U_n - \mu_n$  (not presented by Example 3.2) in a similar spirit.

The insights of the two examples will be generalized in part (b) of our main lemma below. When the network is sufficiently dense, among the expansion terms of  $\hat{U}_n - U_n$ , the linear part dominates. Consequently, the overall contribution of the stochastic variations in  $A|W$  approximates Gaussian at an  $O(\rho_n^{-1/2} \cdot n^{-1})$  Berry-Esseen error rate. Now recalling the definition of acyclic and cyclic  $R$  shapes from Definition 1.1, the definition of  $\mathcal{M}(\rho_n, n; R)$  from definition 1.2 in Section 1, and the definition of  $\tilde{O}_p$ , we are ready to state our main lemma.

LEMMA 3.1. *Assume the following conditions hold:*

- (i).  $\rho_n^{-s} \cdot \xi_1 > C > 0$ , where  $C > 0$  is a universal constant,
- (ii).  $\rho_n = \omega(n^{-1})$  for acyclic  $R$ , or  $\rho_n = \omega(n^{-2/r})$  for cyclic  $R$ ,

We have the following results:

- (a)  $\frac{U_n - \mu_n}{\sigma_n} = U_n^\# + \Delta_n + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n)$ ,
- (b) We have

$$\hat{\Delta}_n = \frac{(\hat{U}_n - U_n)}{\sigma_n} = \check{\Delta}_n + \check{R}_n,$$

where  $\check{\Delta}_n$  and the remainder  $\check{R}_n$  satisfy

$$(3.11) \quad \check{R}_n = \tilde{O}_p(\mathcal{M}(\rho_n, n; R))$$

$$(3.12) \quad \|F_{\check{\Delta}_n|W}(u) - F_{N(0, (\rho_n \cdot n)^{-1} \sigma_w^2)}(u)\|_\infty = \tilde{O}_p(\rho_n^{-1/2} \cdot n^{-1})$$

where the order control in (3.12) is  $\tilde{O}_p(\cdot)$  rather than  $O(\cdot)$  due to the randomness in  $W$ .

The definition of  $\sigma_w$  is lengthy and formally stated in Section 7 in Supplemental Material.

As  $n \rightarrow \infty$ , we have  $\sigma_w \stackrel{p}{\rightarrow} 1$ .

- (c)  $\hat{\delta}_n = \tilde{O}_p(\mathcal{M}(\rho_n, n; R))$ ,
- (d) We have

$$\delta_n = \frac{1}{n} \sum_{i=1}^n \frac{g_1^2(X_i) - \xi_1^2}{\xi_1^2} + \frac{2(r-1)}{n(n-1)} \sum_{\substack{1 \leq \{i,j\} \leq n \\ i \neq j}} \frac{g_1(X_i)g_2(X_i, X_j)}{\xi_1^2} + \tilde{O}_p(n^{-1} \cdot \log n).$$

Overall, Lemma 3.1 clarifies the asymptotic orders of the leading terms in the expansion of  $\hat{T}_n$ . In fact, Lemma 3.1 has a parallel version for the jackknife  $\hat{S}_{n;\text{jackknife}}$  in view of Theorem 3.3, but we do not present it due to page limit. We spend the rest of this section on discussing the conditions and results of Lemma 3.1.

REMARK 3.1. *Assumption (i) is a standard non-degeneration assumption in literature. It is different from a smoothness assumption on graphon  $f$ <sup>16</sup>. A globally smooth Erdos-Renyi graphon leads to a degenerate  $g_1(X_1)$  that  $\xi_1^2 = \text{Var}(g_1(X_1)) = 0$ . In the degenerate setting,*

<sup>16</sup>Smooth graphon:  $f$  is called *smooth*, if there exists a measure-preserving map  $\varrho : [0, 1] \rightarrow [0, 1]$  such that  $f(\varrho(\cdot), \varrho(\cdot))$  is a smooth function. See [56, 136] for more details.

498 both the standardization/studentization and the analysis would be very different. Asymptotic  
 499 results for  $r = 2, 3$  motifs under an Erdos-Renyi graphon have been established by [54, 55].  
 500 Degenerate  $U$ -statistics are outside the scope of this paper.

501 **REMARK 3.2.** We note that Lemma 3.1 only requires the weak assumption on  $\rho_n$  (see  
 502 Assumption ii). This assumption matches the classical sparsity assumptions in network boot-  
 503 strap literature [20, 17, 61]. Using Lemma 3.1, we prove a higher-order error bound of the  
 504 Edgeworth expansion in Theorem 3.1 with a stronger density assumption; while in Theo-  
 505 rem 3.4 on sparse networks, we prove a novel modified Berry-Esseen bound for the normal  
 506 approximation. Both downstream theorems significantly improve over existing best results.

507 **REMARK 3.3.** Lemma 3.1-(a) and (d) are similar to results in classical literature on  
 508 Edgeworth expansion for noiseless  $U$ -statistics [71, 96], but here we account for  $\rho_n$ . Parts  
 509 (b) and (c) are new results unique to the network setting. Especially in the proof of part (b),  
 510 we significantly refine the analysis of the randomness in  $A|W$  in [17] and [61].

511 **3.3. Population and empirical Edgeworth expansions for network moments.** In this sub-  
 512 section, we present our main theorems.

513 **THEOREM 3.1 (Population network Edgeworth expansion).** Define

$$514 \quad G_n(x) := \Phi(x) + \frac{\varphi(x)}{\sqrt{n} \cdot \xi_1^3} \cdot \left\{ \frac{2x^2 + 1}{6} \cdot \mathbb{E}[g_1^3(X_1)] \right. \\ 515 \quad \left. + \frac{r-1}{2} \cdot (x^2 + 1) \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\}, \quad (3.13)$$

516 where  $\Phi(x)$  and  $\varphi(x)$  are the CDF and PDF of  $N(0, 1)$ , respectively. Assume condition (i)  
 517 of Lemma 3.1 hold, and replace condition (ii) by a stronger assumption that either  $R$  is  
 518 acyclic and  $\rho_n = \omega(n^{-1/2})$ , or  $R$  is cyclic and  $\rho_n = \omega(n^{-1/r})$ . Additionally, assume either  
 519  $\rho_n = O((\log n)^{-1})$  or Cramer's condition  $\limsup_{t \rightarrow \infty} \left| \mathbb{E} \left[ e^{itg_1(X_1) \cdot \xi_1^{-1}} \right] \right| < 1$  holds. We have

$$\left\| F_{\hat{T}_n}(x) - G_n(x) \right\|_{\infty} = O(\mathcal{M}(\rho_n, n; R)).$$

520 **REMARK 3.4.** The assumed  $\rho_n$ 's upper bound in absence of Cramer's condition serves  
 521 to sufficiently boost the smoothing power of  $\hat{\Delta}_n$ , quantified in Lemma 3.1-(3.12). This as-  
 522 sumption seems minimal in presence of a lattice  $g_1(X_1)$ , since it corresponds to a normal  
 523 smoother with variance  $(\rho_n \cdot n)^{-1} = \Omega(\log n \cdot n^{-1})$ . This matches the minimum standard  
 524 deviation requirement  $\Omega((\log n)^{1/2} \cdot n^{-1/2})$  in Remark 2.4 in [89] for the i.i.d. setting.

525 In (3.13), the Edgeworth coefficients depend on true population moments. In practice, they  
 526 need to be estimated from data. Define

$$527 \quad \hat{g}_1(X_i) := \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(A_{i, i_1, \dots, i_{r-1}}) - \hat{U}_n, \\ 528 \quad \hat{g}_2(X_i, X_j) := \frac{1}{\binom{n-2}{r-2}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-2} \leq n \\ i_1, \dots, i_{r-2} \neq i, j}} h(A_{i, j, i_1, \dots, i_{r-2}}) - \hat{U}_n - \hat{g}_1(X_i) - \hat{g}_1(X_j),$$

529 where we write “ $\widehat{g}_1(X_i)$ ” rather than “ $\widehat{g}_1(\widehat{X}_i)$ ” for cleanness. We stress that the evaluation of  
 530  $\widehat{g}_1(X_i)$  and  $\widehat{g}_2(X_i, X_j)$  does *not* require knowing the latent  $X_i, X_j$ . The Edgeworth coefficients  
 531 can be estimated by

$$532 \quad \widehat{\xi}_1^2 := \frac{n \cdot \widehat{S}_n^2}{r^2} = \frac{1}{n} \sum_{i=1}^n \widehat{g}_1^2(X_i), \quad \text{and} \quad \widehat{\mathbb{E}}[g_1^3(X_1)] := \frac{1}{n} \sum_{i=1}^n \widehat{g}_1^3(X_i),$$

$$533 \quad \widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] := \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \widehat{g}_1(X_i)\widehat{g}_1(X_j)\widehat{g}_2(X_i, X_j).$$

534 **THEOREM 3.2** (Empirical network Edgeworth expansion). *Define the empirical Edge-*  
 535 *worth expansion as follows:*

$$536 \quad \widehat{G}_n(x) := \Phi(x) + \frac{\varphi(x)}{\sqrt{n} \cdot \widehat{\xi}_1^3} \cdot \left\{ \frac{2x^2 + 1}{6} \cdot \widehat{\mathbb{E}}[g_1^3(X_1)] \right.$$

$$537 \quad (3.14) \quad \left. + \frac{r-1}{2} \cdot (x^2 + 1) \widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\},$$

538 *Under the conditions of Theorem 3.1, we have*

$$\left\| F_{\widehat{T}_n}(x) - \widehat{G}_n(x) \right\|_{\infty} = \widetilde{O}_p(\mathcal{M}(\rho_n, n; R)).$$

539 **REMARK 3.5.** *Another approach to estimate the unknown coefficients in Edgeworth ex-*  
 540 *pansion is bootstrap. The concentration of  $\widehat{G}_n \rightarrow G_n$  should not be confused with the con-*  
 541 *centration  $\widehat{G}_n^* \rightarrow \widehat{G}_n$ , where  $\widehat{G}_n^*$  is the expansion with bootstrap-estimated coefficients. See*  
 542 *literature regarding the i.i.d. setting [71, 96]. In  $\widehat{G}_n^* \rightarrow \widehat{G}_n$ , the convergence rate is not a*  
 543 *concern, because without constraining computation cost, one can let the number of bootstrap*  
 544 *samples grow arbitrarily fast. Hence, establishing consistency would suffice for the analysis*  
 545 *of  $\widehat{G}_n^* \rightarrow \widehat{G}_n$ , whereas our proof concerning  $\widehat{G}_n \rightarrow G_n$  requires careful rate calculations.*

546 Next, we show that different choices of the variance estimators for studentization represent  
 547 no essential discrepancy.

**THEOREM 3.3** (Studentizing by a jackknife variance estimator (3.3)). *Define*

$$\widehat{T}_{n;\text{jackknife}} := \frac{\widehat{U}_n - \mu_n}{\widehat{S}_{n;\text{jackknife}}}.$$

548 *Under the assumptions of Theorem 3.1, we have*

$$549 \quad (3.15) \quad |\widehat{S}_n - \widehat{S}_{n;\text{jackknife}}| = O(\widehat{S}_n \cdot n^{-1}),$$

$$550 \quad \left\| F_{\widehat{T}_{n;\text{jackknife}}}(x) - G_n(x) \right\|_{\infty} = O(\mathcal{M}(\rho_n, n; R)),$$

$$551 \quad \left\| F_{\widehat{T}_{n;\text{jackknife}}}(x) - \widehat{G}_n(x) \right\|_{\infty} = \widetilde{O}_p(\mathcal{M}(\rho_n, n; R)).$$

552 **Theorem 3.3** states that on statistical properties, one does not need to differentiate between  
 553  $\widehat{T}_n$  and  $\widehat{T}_{n;\text{jackknife}}$ . The evaluation of  $\widehat{S}_{n;\text{jackknife}}$  costs  $O(n^{r+1})$  time because each individual  
 554  $\widehat{U}_n^{(-i)}$  costs  $O(n^r)$ ; whereas our estimator  $\widehat{S}_n$  costs  $O(n^r)$ . Our estimator also has a more  
 555 convenient form for theoretical analysis.

556 3.4. *Remarks on non-smooth graphons.* Our results do not assume graphon smoothness  
 557 or low-rankness. This aligns with the literature on noiseless U-statistics but sharply con-  
 558 trasts network inferences based on model parameter estimation such as [74, 91] and network  
 559 bootstraps based on model estimation [61, 93]. Notice that the concept “non-smoothness”  
 560 usually emphasizes “not assuming smoothness” rather than explicitly describing irregularity.  
 561 It is a very useful tool for modeling networks with high structural complexity or unbalanced  
 562 observations, examples include: (1) a small group of *outlier* nodes that behave differently  
 563 from the main network patterns [28]; (2) in networks that exhibit “core-periphery” structures  
 564 [48, 138], we may wish to relax structural assumptions on periphery nodes due to scarcity of  
 565 observations; and (3) networks generated from a mixture model [104, 82] with many small-  
 566 probability mixing components may appear non-smooth in these parts. Unfortunately, ex-  
 567 isting research on practical methods for non-smooth graphons is rather limited due to the  
 568 obvious technical difficulty, but exceptions include [40].

569 Our results send the surprising message that under mild conditions, the sampling distri-  
 570 bution of a network moment is still *smooth* and can be *accurately* approximated, even if the  
 571 graphon is non-smooth.

572 3.5. *Sparse networks.* We have been focusing on discussing dense networks, but many  
 573 networks tend to be sparse [63]. In this section, we investigate the following sparsity regime

$$(3.16) \quad \rho_n : \begin{cases} n^{-1} < \rho_n \leq n^{-1/2}, & \text{for acyclic } R \\ n^{-2/r} < \rho_n \leq n^{-1/r}, & \text{for cyclic } R \end{cases}$$

574 It turns out that the Berry-Esseen bound in this setting would be slower than  $n^{-1/2}$ , unlike  
 575 that in i.i.d. and noiseless U-statistic settings. The exact reason is technical and will be better  
 576 seen in the proof of Theorem 3.4, but the intuitive explanation is that if  $\rho_n$  is too small,  
 577 the higher degree ( $\geq 2$ ) random errors in  $\hat{U}_n - U_n$  diminishes too slowly compared to the  
 578 scale of the demoninator of  $\hat{T}_n$ . If the network sparsity  $\rho_n$  falls below the typically assumed  
 579 lower bounds:  $n^{-1}$  for acyclic  $R$  and  $n^{-2/r}$  for cyclic  $R$  [20, 17, 61], no known consistency  
 580 guarantee exists. In fact, in this case we do not even know if  $\hat{T}_n$  is asymptotically normal.

581 **THEOREM 3.4.** *Under the conditions of Lemma 3.1, except replacing Condition (ii) by*  
 582 *(3.16), we have the following modified Berry-Esseen bound*

$$583 \quad \left\| F_{\hat{T}_n}(u) - G_n(u) \right\|_{\infty} \asymp \left\| F_{\hat{T}_n}(u) - \Phi(u) \right\|_{\infty} = O(\mathcal{M}(\rho_n, n; R)) \bigwedge o(1),$$

where recall that  $\Phi(\cdot)$  is the CDF of  $N(0, 1)$ . Moreover,

$$\left\| F_{\hat{T}_n}(u) - \hat{G}_n(u) \right\|_{\infty} = \tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \bigwedge o_p(1)$$

584 In the sparse regime, the current upper bound on the remainder terms would dominate  
 585 the  $n^{-1/2}$  leading term in the Edgeworth expansion. In other words, in sparse networks,  
 586 the Edgeworth expansion is guaranteed by the same error rate bound as a simple  $N(0, 1)$   
 587 approximation. On the other hand, the conclusion of Theorem 3.4 connects the error bound  
 588 results for dense and sparse regimes. Interestingly, as the order of  $\rho_n$  decreases from  $n^{-1/2}$  to  
 589  $n^{-1}$  for acyclic  $R$ , or from  $n^{-1/r}$  to  $n^{-2/r}$  for cyclic  $R$ , we see a gradual depreciation in the  
 590 uniform CDF approximation error from the order of  $n^{-1/2}$  to merely uniform consistency.  
 591 The classical literature only studied the boundary cases ( $\rho_n = \omega(n^{-1})$  or  $\rho_n = \omega(n^{-2/r})$ ),  
 592 depending on  $R$ ), and our result here reveals the complete picture.

593 A natural question is whether a higher-order approximation would be possible in the sparse  
 594 regime. We conjecture not. We also conjecture that the Berry-Esseen bound that both empiri-  
 595 cal Edgeworth expansion and  $N(0, 1)$  approximation achieve is either sharp or nearly sharp,  
 596 but we do not know an answer for sure. This would be an interesting future work.

597 3.6. *Comparison table of our method to benchmarks.* We conclude this section by com-  
598 paring our results to some representative works in classical and very recent literature.

TABLE 2  
*Comparison of CDF approximation methods for noisy/noiseless studentized U-statistics*

Method	U-stat. type	Popul. momt. <sup>17</sup>	Smooth graphon	Non lat. /Cramer	Network sparsity assumption on $\rho_n$ <sup>18</sup>	CDF approx. error rate
Our method (empirical Edgeworth)	Noisy	No	No	If yes	$\omega(n^{-2/r})(C); \omega(n^{-1})(Ac)$ <sup>19</sup>	$\tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \wedge o_p(1) (\mathbf{H})$ <sup>20</sup>
				If no	$\omega(n^{-2/r})(C); \omega(n^{-1})(Ac)$ and $O((\log n)^{-1})(C, Ac)$	$\tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \wedge o_p(1) (\mathbf{H})$
Node re-/sub- sampling justified by our theory	Noisy	No	No	Yes	$\omega(n^{-1/r})(C); \omega(n^{-1/2})(Ac)$	$o_p(n^{-1/2}) (\mathbf{H})$
Bickel, Chen and Levina [20]	Noisy	No <sup>21</sup>	No	No	$\omega(n^{-2/r})(C); \omega(n^{-1})(Ac)$	Consistency
Bhattacharyya and Bickel [17]	Noisy	No	No	No	$\omega(n^{-2/r})(C); \omega(n^{-1})(Ac)$	Consistency
Green and Shalizi [61]	Noisy	No	Mixed <sup>22</sup>	No	$R$ is Ac; or $\omega(n^{-1/(2r)})(C)$ <sup>23</sup>	Consistency
Levin and Levina [93]	Noisy	No	Low-rank <sup>24</sup>	No	$\omega(n^{-1} \cdot \log n) (Ac^*)$ <sup>25</sup>	Consistency
Bickel, Gotze and Zwet [21]	Noiseless	Yes	No	Yes	Not applicable	$o(n^{-1}) (\mathbf{H})$
Bentkus, Gotze and Zwet [14]	Noiseless	Yes	No	Yes	Not applicable	$O(n^{-1}) (\mathbf{H})$
Putter and Zwet [110]	Noiseless	No	No	Yes	Not applicable	$o_p(n^{-1/2}) (\mathbf{H})$
Bloznelis [23]	Noiseless	No	No	Yes	Not applicable	$o_p(n^{-1/2}) (\mathbf{H})$

#### 599 4. Theoretical and methodological applications.

600 4.1. *Higher-order accuracy of node sub- and re-sampling network bootstraps.* One im-  
601 portant corollary of our results is first higher-order accuracy proof of some network bootstrap  
602 schemes. For a network bootstrap scheme that produces an estimated  $\hat{U}_{n^*}^b$  and its jackknife<sup>26</sup>  
603 variance estimator  $\hat{S}_{n^*}^*$ , define  $\hat{T}_{n^*}^* = (\hat{U}_{n^*}^b - \hat{U}_n) / \hat{S}_{n^*}^*$ . We are going to establish the first  
604 explicit rate guarantees for following two schemes.

- 605 (a). Sub-sampling [17]: randomly sample  $n^*$  nodes from  $\{1, \dots, n\}$  *without replacement*,  
606 and compute  $\hat{T}_{n^*}^*$  from the induced sub-network of  $A$ .
- 607 (b). Re-sampling [61]: random sample  $n$  nodes from  $\{1, \dots, n\}$  *with replacement*, and com-  
608 pute  $\hat{T}_{n^*}^*$  from the induced sub-network of  $A$ .

<sup>17</sup>“Yes” means need to know the population moments that appear in Edgeworth coefficients, i.e.  $\xi_1, \mathbb{E}[g_1^3(X_1)]$  and  $\mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$ .

<sup>18</sup>To compare  $\rho_n$  assumptions, see our Remark 3.2

<sup>19</sup>(C): cyclic  $R$ ; (Ac): acyclic  $R$ .

<sup>20</sup>Recall  $\mathcal{M}(\rho_n, n; R)$  defined in (1.1) and  $\tilde{O}_p$  defined in Section 1.4. (H): higher-order accuracy results. “Consistency”: only convergence, no error rate.

<sup>21</sup>In [20, 17],  $\hat{U}_n - \mu_n$  was rescaled by  $\rho_n$  and  $n$ . Whether assuming the knowledge of the true  $\rho_n$  or not does not matter for their  $o_p(1)$  error bound, but it would make a difference if an  $o_p(n^{-1/2})$  or finer bound is desired.

<sup>22</sup>The bootstrap based on denoised  $A$  requires smoothness. See Theorem 2 of [61].

<sup>23</sup>It seems their assumption for cyclic  $R$  was a typo, and  $\rho_n = \omega(n^{-2/r})$  should suffice. Also, they used [17] in their proof, which requires  $\rho_n = \omega(n^{-1})$  for (Ac).

<sup>24</sup>[93] assumed the graphon rank is low and known.

<sup>25</sup>(Ac\*): They require the motif to be either acyclic or an  $r$ -cycle, see their Theorem 4. Their Theorem 3 requires condition (8) that only holds when  $R$  is a clique.

<sup>26</sup>Here, we use jackknife variance estimator in bootstraps to better connect with existing literature in the proof.

609 **REMARK 4.1.** Notice that [61] did not study the studentized form, and [17] proposed a  
 610 different variance estimator (what they call “ $\hat{\sigma}_{B_i}$ ”). Our justifications focus on the sampling  
 611 schemes combined with some natural formulation, not necessarily the same formulation as  
 612 in these two papers.

613 **REMARK 4.2.** As noted in [61], scheme (b) can be viewed as our data generation pro-  
 614 cedure described in Sections 2.1 and 2.2 but with the graphon  $f$  replaced by the adjacency-  
 615 induced graphon  $A(u, v) = A_{\lceil nu \rceil, \lceil nv \rceil}$ , where  $\lceil y \rceil := \text{Ceiling}(y)$ . This may seem similar to  
 616  $f$ -based data generaion, but in fact they are distinct. The graphon  $A(\cdot, \cdot)$  inherits the binary  
 617 nature of  $A$  and will necessarily yield a lattice  $g_1^*(X_1^*)$  regardless of the original graphon  $f$   
 618 and the motif  $\mathcal{R}$ , rendering most classical Edgeworth analysis techniques inapplicable. But  
 619 the real obstacle is that the bootstrapped network data from  $A(\cdot, \cdot)$  have no edge-wise ob-  
 620 servational errors (i.e. no counterpart to the randomness in  $A|W$ ). Consequently,  $\hat{T}_{n^*}$  loses  
 621 the self-smoothing feature that  $\hat{T}_n$  enjoys. For this reason, when justifying the higher-order  
 622 accuracy of network bootstraps, we cannot simply reproduce the proof of our main theorem  
 623 that crucially benefits from the self-smoothing effect. Aligned with this observation, the even-  
 624 tual error rates that we established for network bootstraps are significantly worse than our  
 625 population and empirical Edgeworth expansions. We conjecture that further improving the  
 626 error guarantee for network bootstraps beyond Theorem 4.1, if possible, might require much  
 627 more sophisticated analysis.

628 **THEOREM 4.1.** Assume  $g_1(X_1)$  satisfies a Cramer’s condition such that  $\limsup_{t \rightarrow \infty} \left| \mathbb{E} \left[ e^{itg_1(X_1) \cdot \xi_1^{-1}} \right] \right| <$   
 629 1. Under the conditions of Theorem 3.2, we conclude for the following bootstrap schemes:

630 (a). Sub-sampling: choosing  $n^* \asymp n$  and  $n - n^* \asymp n$ , we have

$$(4.1) \quad \left\| F_{\hat{T}_{n^*}}(u) - F_{\hat{T}_{n^*(1-n^*/n)}}(u) \right\|_{\infty} = o_p \left( (n^*)^{-1/2} \right) = o_p(n^{-1/2}).$$

631 (b). Re-sampling: choosing  $n^* = n$ , we have

$$(4.2) \quad \left\| F_{\hat{T}_{n^*}}(u) - F_{\hat{T}_n}(u) \right\|_{\infty} = o_p \left( (n^*)^{-1/2} \right) = o_p(n^{-1/2}).$$

632 **REMARK 4.3.** In the proof of Theorem 4.1, we combined our main results with the results  
 633 of [23] for finite population  $U$ -statistics. It is important to notice that all existing works  
 634 under the finite populations did assume non-lattice with population size growing to infinity,  
 635 see condition (1.13) in Theorem 1 of [23]. Consequently, the higher-order accuracy of some  
 636 network bootstraps is only proved under Cramer’s condition by so far.

637 Part (a) of Theorem 4.1 quantifies the effective sample size in the sub-sampling network  
 638 bootstrap: sampling  $n^*$  out of  $n$  nodes without replacement, the resulting bootstrap  $\hat{T}_{n^*}$  ap-  
 639 proximates the distribution of  $\hat{T}_m$  where  $m = \{n^*/n \cdot (1 - n^*/n)\} \times n$ . Consequently, in  
 640 order to approach the sampling distribution of  $\hat{T}_n$  with higher-order accuracy using sub-  
 641 sampling [17], one must have an observed network of at least  $4n$  nodes, from which she shall  
 642 repeatedly sub-sample  $2n$  nodes without replacement.

643 **4.2. One-sample  $t$ -test for network moments under general null graphon models.** In this  
 644 and the next subsections, we showcase how our results immediately lead to useful inference  
 645 procedures for network moments. For a given motif  $R$ , we test on its population mean fre-  
 646 quency  $\mu_n$ . Since  $\mu_n$  depends on  $n$  through  $\rho_n$ , we formulate the hypotheses as follows

$$647 \quad H_0 : \mu_n = c_n, \text{ versus } H_a : \mu_n \neq c_n.$$

648 where  $c_n$  is a speculated value of  $\mu_n = \mathbb{E}[h(A_{1,\dots,r})]$ . In practice,  $c_n$  may come from a prior  
 649 study on a similar data set or fitting a speculated model to the data (for concrete examples on  
 650  $c_n$  guesses, see Section 6.1 of [17]).

651 Here for simplicity we only discuss a two-sided alternative, and one-sided cases are exactly  
 652 similar. The p-value can be formulated using our empirical Edgeworth expansion  $\widehat{G}_n(\cdot)$  in  
 653 (3.14):

$$(4.3) \quad \text{Estimated p-value} = 2 \cdot \min \left\{ \widehat{G}_n(t^{(\text{obs})}), 1 - \widehat{G}_n(t^{(\text{obs})}) \right\}.$$

654 where  $t^{(\text{obs})} := (\widehat{u}_n^{(\text{obs})} - c_n) / \widehat{s}_n^{(\text{obs})}$ , and  $\widehat{u}_n^{(\text{obs})}$  and  $\widehat{s}_n^{(\text{obs})}$  are the observed  $\widehat{U}_n$  and  $\widehat{S}_n$ , re-  
 655 spectively. We have the following explicit Type-II error rate.

656 **THEOREM 4.2.** *Under the conditions of Theorem 3.2, we have the following results:*

- 657 1. *The Type-I error rate of test (4.3) is  $\alpha + O(\mathcal{M}(\rho_n, n; R))$ .*
- 658 2. *The Type-II error rate of this test is  $o(1)$  when  $|c_n - \mu_n| = \omega(\rho_n^s \cdot n^{-1/2})$ .*

659 **REMARK 4.4.** *The null model we study is complementary to the degenerate Erdos-  
 660 Renyi null model in [91, 54, 55]. The scientific questions are also different: they test model  
 661 goodness-of-fit whereas we test population moment values. Notice that distinct network mod-  
 662 els may possibly share some common population moments. These approaches also use very  
 663 different methods and analysis techniques.*

664 **4.3. Cornish-Fisher confidence intervals for network moments.** Noticing that  $\widehat{G}_n$  is al-  
 665 most never a valid CDF, in order to preserve the higher-order accuracy of  $\widehat{G}_n$ , we use the  
 666 Cornish-Fisher expansion [44, 53] to approximate the quantiles of  $F_{\widehat{T}_n}$ . A Cornish-Fisher  
 667 expansion is the inversion of an Edgeworth expansion, and its validity hinges on the validity  
 668 of its corresponding Edgeworth expansion. Using the technique of [65], we have

669 **THEOREM 4.3.** *For any  $\alpha \in (0, 1)$ , define the lower  $\alpha$  quantile of the distribution of  $\widehat{T}_n$   
 670 by*

$$(4.4) \quad q_{\widehat{T}_n; \alpha} := \arg \inf_{q \in \mathbb{R}} F_{\widehat{T}_n}(q) \geq \alpha$$

671 *and define the approximation*

$$(4.5) \quad \widehat{q}_{\widehat{T}_n; \alpha} := z_\alpha - \frac{1}{\sqrt{n} \cdot \widehat{\xi}_1^3} \cdot \left\{ \frac{2z_\alpha^2 + 1}{6} \cdot \widehat{\mathbb{E}}[g_1^3(X_1)] \right. \\ \left. + \frac{r-1}{2} \cdot (z_\alpha^2 + 1) \widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\},$$

674 where  $z_\alpha := \Phi^{-1}(\alpha)$ . Then under the conditions of Theorem 3.2, we have

675 (a). *The discrepancy between nominal and actual percentage-below for  $q_{\widehat{T}_n; \alpha}$  is bounded by*

$$(4.6) \quad |F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha}) - \alpha| = O(\mathcal{M}(\rho_n, n; R))$$

676 (b). *The ‘‘horizontal’’ error bound:*

$$(4.7) \quad \left| \widehat{q}_{\widehat{T}_n; \alpha} - q_{\widehat{T}_n; \alpha} \right| = \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$$

677 (c). *The uniform ‘‘vertical’’ error bound:*

$$(4.8) \quad \mathbb{P}(\widehat{T}_n \leq \widehat{q}_{\widehat{T}_n; \alpha}) = \alpha + O(\mathcal{M}(\rho_n, n; R)).$$

678 The vertical error bound describes the approximation error between the nominal and ac-  
 679 tual coverage probabilities, whereas the horizontal error bound governs the approximation  
 680 of quantiles. Using the vertical error bound, a  $1 - \alpha$  two-sided symmetric Cornish-Fisher  
 681 confidence interval for estimating  $\mu_n$  can be easily constructed as follows

$$(4.9) \quad \left( \hat{U}_n - \hat{q}_{\hat{T}_n; 1-\alpha/2} \cdot \hat{S}_n, \hat{U}_n - \hat{q}_{\hat{T}_n; \alpha/2} \cdot \hat{S}_n \right)$$

682 and by Theorem 4.3, we know this CI has a  $1 - \alpha + O(\mathcal{M}(\rho_n, n; R))$  coverage probability.  
 683 One-sided confidence intervals can be constructed exactly similarly, thus we omit them.

684 **5. Simulations.**

685 *5.1. Simulation 1: Higher-order accuracy of empirical Edgeworth expansion.* In the first  
 686 simulation, our numerical studies focus on the CDF of  $F_{\hat{T}_n}$ . In an illustrative example, we  
 687 simulate with a lattice  $g_1(X_1)$  and show the distinction between  $F_{\hat{T}_n}$  and  $F_{T_n}$  that clearly  
 688 illustrates the self-smoothing effect in  $\hat{T}_n$ . Then we systematically compare the performance  
 689 of our empirical Edgeworth expansion to benchmarks that demonstrates the clear advantage  
 690 of our method in both accuracy and computational efficiency.

We begin by describing the basic settings. We range the network size  $n$  in an exponen-  
 tially spaced set  $n \in \{10, 20, 40, 80, 160\}$ . Synthetic network data are generated from three  
 graphons marked by their code-names in our figures: (1). "BlockModel": This is an ordi-  
 nary stochastic block model with  $K = 2$  equal-sized communities and the following edge  
 probabilities  $B = (0.6, 0.2; 0.2, 0.2)$ ; (2). "SmoothGraphon": Graphon 4 in [136], i.e.  
 $f(u, v) := (u^2 + v^2)/3 \cdot \cos(1/(u^2 + v^2)) + 0.15$ . This graphon is smooth and full-rank [136];  
 (3). "NonSmoothGraphon"[40]: We set up a high-fluctuation area in a smooth  $f$  to emu-  
 late the sampling behavior of a non-smooth graphon, as follows

$$f(u, v) := 0.5 \cos \{0.1 / ((u - 1/2)^2 + (v - 1/2)^2)^{-1} + 0.01\} \max\{u, v\}^{2/3} + 0.4.$$

691 We test the four simplest motifs: *edge*, *triangle*, *V-shape*<sup>27</sup>, and a *three-star* among 4 nodes  
 692 with edge set  $\{(1, 2), (1, 3), (1, 4)\}$ . The main computation bottleneck lies in the evaluation  
 693 of  $F_{\hat{T}_n}$ . Network bootstraps also becomes costly as  $n$  increases.

694 The benchmarks are: 1.  $N(0, 1)$  (its computation time is deemed zero and not compared  
 695 to others); 2. sub-sampling by [17] with  $n^* = n/2$ ; 3. re-sampling  $A$  by [61]; 4. latent space  
 696 bootstrap called ‘‘ASE plug-in’’ defined in Theorem 2 of [93]. Notice that we equipped [93]  
 697 with an adaptive network rank estimation<sup>28</sup> by USVT [35].

698 For each (graphon, motif,  $n$ ) tuple, we first evaluate the true sampling distribution of  $\hat{T}_n$   
 699 by a Monte-Carlo approximation that samples  $n_{MC} := 10^6$  networks from the graphon. Next  
 700 we start 30 repeated experiments: in each iteration, we sample  $A$  from the graphon and ap-  
 701 proximate  $F_{\hat{T}_n}$  by all methods, in which we draw  $n_{boot} = 2000$  bootstrap samples for each  
 702 bootstrap method – notice that this is 10 times that in [93]. We compare

$$(5.1) \quad \text{Error}(\hat{F}_{\hat{T}_n}) := \sup_{u \in [-2, 2]; 10u \in \mathbb{Z}} \left| \hat{F}_{\hat{T}_n}(u) - F_{\hat{T}_n}(u) \right|.$$

703 **REMARK 5.1.** *We need many Monte-Carlo repetitions, because the uniform accuracy of*  
 704 *the empirical CDF of an i.i.d. sample is only  $O_p(n_{MC}^{-1/2})$  [50, 88], and for the noiseless and*

<sup>27</sup>A ‘‘V-shape’’ is the motif obtained by disconnecting one edge in a triangle. In the language of [20], it is a 2-star.

<sup>28</sup>Consequently, our enhanced version of this benchmark can decently denoise some smooth but high-rank graphons, in view of the remarks in [136] and the results of [134].

705 *noisy U-statistic setting, the bound might be worse than the i.i.d. setting due to dependency*<sup>29</sup>.  
 706 *Therefore, we set  $n_{\text{MC}} \gg \max(n^2) = 160^2$  to prevent the errors of the compared methods*  
 707 *being dominated by the error of the Monte-Carlo procedure; while keep our simulations*  
 708 *reproducible with moderate computation cost. We did find smaller  $n_{\text{MC}}$  such as  $10^5$  to cloud*  
 709 *the performance of our method.*

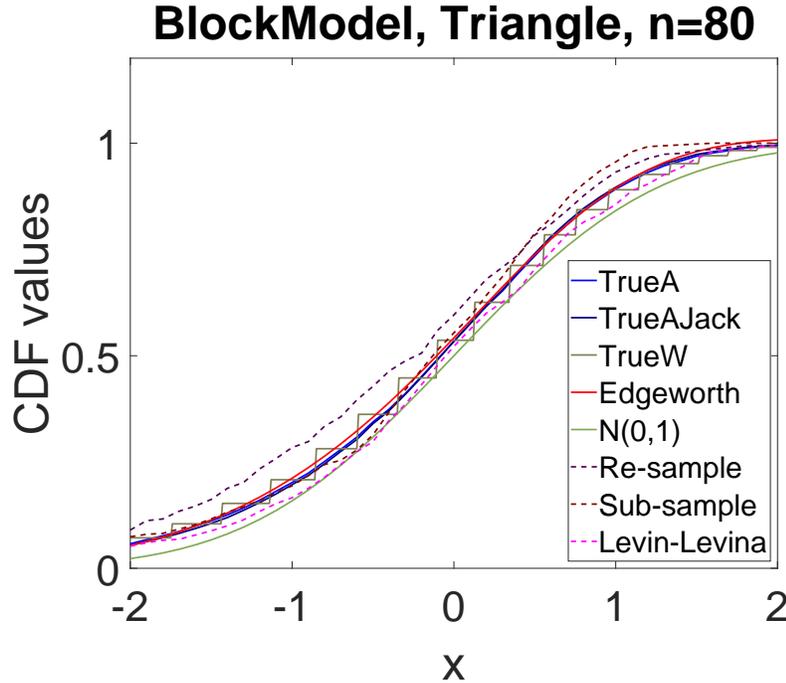


Fig 1: CDF curves of the studentization forms and approximations. Network size  $n = 80$ . The graphon is the “BlockModel” we described earlier in this section, and the motif is triangular. Each bootstrap method draws 500 random samples. TrueA is  $F_{\hat{T}_n}$ ; TrueAJack is  $F_{\hat{T}_n;\text{jackknife}}$ ; TrueW is  $F_{T_n}$ ; Edgeworth is our empirical Edgeworth expansion; Re-sample is node re-sampling A in [61]; Sub-sample is node sub-sampling A in [17]; Levin-Levina is the “ASE plug-in” bootstrap in [93].

710 Now we present the results. We first present the illustrative simulation for just one specific  
 711 setting. Figure 1 shows the distribution approximation curves under a block model graphon  
 712 that yields a lattice  $g_1(X_1)$ . Lines correspond to the population CDF of  $\hat{T}_n$ , its jackknife ver-  
 713 sion and noiseless version, all evaluated by Monte-Carlo procedures; our proposed empirical  
 714 Edgeworth expansion; and benchmarks. We make two main observations. First, TrueA and  
 715 TrueAJack are almost indistinguishable, echoing our Theorem 3.3; meanwhile, they are  
 716 both smooth and rather different from the step-function TrueW. This clearly demonstrates  
 717 the self-smoothing feature of  $\hat{T}_n$  in the lattice case. If we change the graphon to a smooth  
 718 one, these curves would all be smooth and close to each other. Second, we observe the higher  
 719 accuracy of our empirical Edgeworth expansion compared to competing methods. In fact,

<sup>29</sup>This is not to be confused with the Edgeworth approximation error bound. In this Monte Carlo procedure, both the true and approximate  $F_{\hat{T}_n}$  are oracle.

720 repeating this experiment multiple times, our method shows significantly stabler approxima-  
721 tions than bootstraps.

722 Next, we conduct a systematic comparison of the performances of all methods across many  
723 settings. We mainly varied three factors: graphon type, motif type and network size, over the  
724 previously described ranges. Our experiment results under different network sparsity levels  
725 would have to sink to Supplemental Material due to page limit, and here we keep  $\rho_n = 1$ .  
726 Results are shown in Figure 2 (error) and Figure 3 (time cost), where error bars show standard  
727 deviations.

728 In all experiments, our empirical Edgeworth expansion approach exhibited clear advan-  
729 tages over benchmark methods in all aspects: the absolute values of errors, the diminishing  
730 rates of errors, and computational efficiency. Our method shows a higher-order accuracy by  
731 slopes steeper than  $-1/2$  and much steeper than other methods. On computation efficiency,  
732 our method is the second cheapest after the simple  $N(0, 1)$  approximation (that does not need  
733 computation) and much faster than network bootstraps. It typically costs about  $e^{-5} \approx 1/150$   
734 the time of sub-sampling and about  $e^{-7} \approx 1/1000$  the time of re-sampling. Our method only  
735 needs one run and does not require repeated sampling.

736 Notice that there is no simple rule to judge the difficulty of different scenarios, which  
737 jointly depends on the graphon and the motif through implicit and complex relationship. In  
738 our experience, triangle may be more difficult than V-shape under some graphons, but easier  
739 under some others, and this comparison may vary from method to method. Answering this  
740 question requires calculation of the population Edgeworth expansion up to  $o(n^{-1})$  remainder,  
741 and the leading term in the remainder of the one-term Edgeworth expansion would then  
742 quantify the real difficulty. But the calculation is very complicated and outside the scope of  
743 this paper.

744 We did not observe the higher-order accuracy of bootstrap methods as our results pre-  
745 dicted. One likely reason is the numerical accuracy limited by the  $n_{\text{boot}}$  that our computing  
746 servers can afford. We did see an observable improvement in the performances of network  
747 bootstraps as we increased  $n_{\text{boot}}$  from 200 suggested by [93] to the current 2000. But further  
748 increasing  $n_{\text{boot}}$  will also increase their time costs and potentially memory usage. We ran  
749 each experiment on 36 parallel Intel(R) Xeon(R) X5650 CPU cores at 2.67GHz with 12M  
750 cache and 2GB RAM. It took roughly 3~8 hours to run each experiment that produces one  
751 individual plot in Figures 2 and 3.

752 *5.2. Simulation 2: Finite-sample performance of Cornish-Fisher confidence interval.* In  
753 this simulation, we numerically assess the performance of our Cornish-Fisher confidence in-  
754 terval, compared to benchmark methods. Throughout this subsection, we set  $\alpha = 0.2$  and  
755 focus on symmetric two-sided confidence intervals. We inherit most simulation settings from  
756 Section 5.1 with some modifications we now clarify. The main difference is that in this sim-  
757 ulation, we must conduct many repeated experiments in order to accurately evaluate the cov-  
758 erage probability (each iteration produces a binary outcome of whether the CI contains the  
759 population parameter). We repeated the experiment 10000 times for our method and normal  
760 approximation, and 500 times for the much slower bootstrap methods. Due to the computer  
761 limitations, while we can keep the same number of Monte Carlo evaluations, in order to re-  
762 peat the entire experiment 500 times to accurately evaluate the actual CI coverage rates of  
763 bootstraps, we have to reduce their numbers of bootstrap samples to 500 (still more than the  
764 200 in [93]). We evaluate three performance measures: `coverage`: *actual coverage proba-*  
765 *bility*; `length`: *confidence intereval length*; and `time`: *computation time in seconds*.

766 Due to page limit, in the main text, we only present the results for the setting  $n = 80$  and  
767  $\rho_n = 1$  in Tables 3 (block model), 4 (smooth graphon) and 5 (non-smooth graphon). Each en-  
768 try is formatted “mean(standard deviation)”. We sink the remaining results to Supplemental

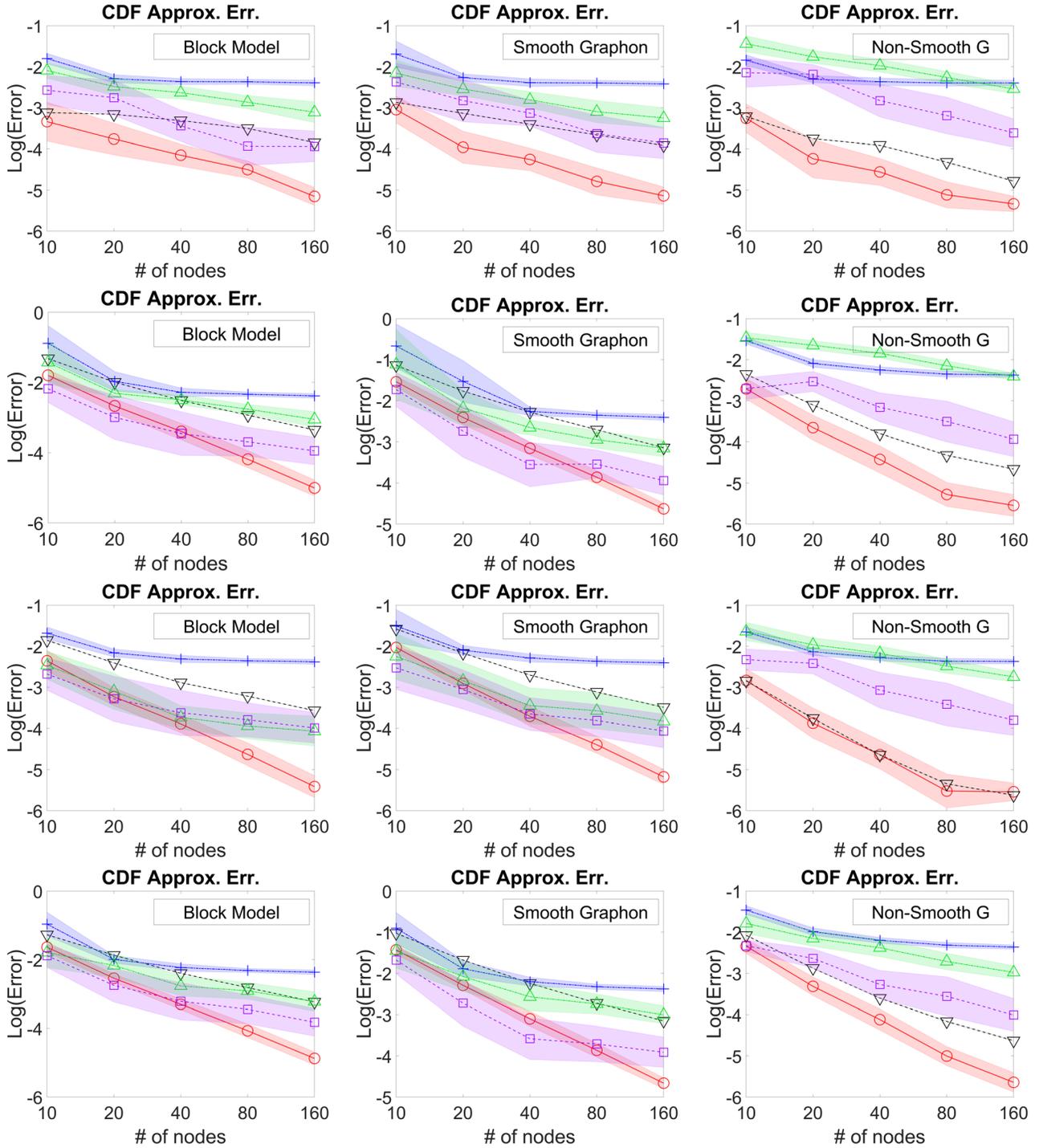


Fig 2: CDF approximation errors. Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\approx n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

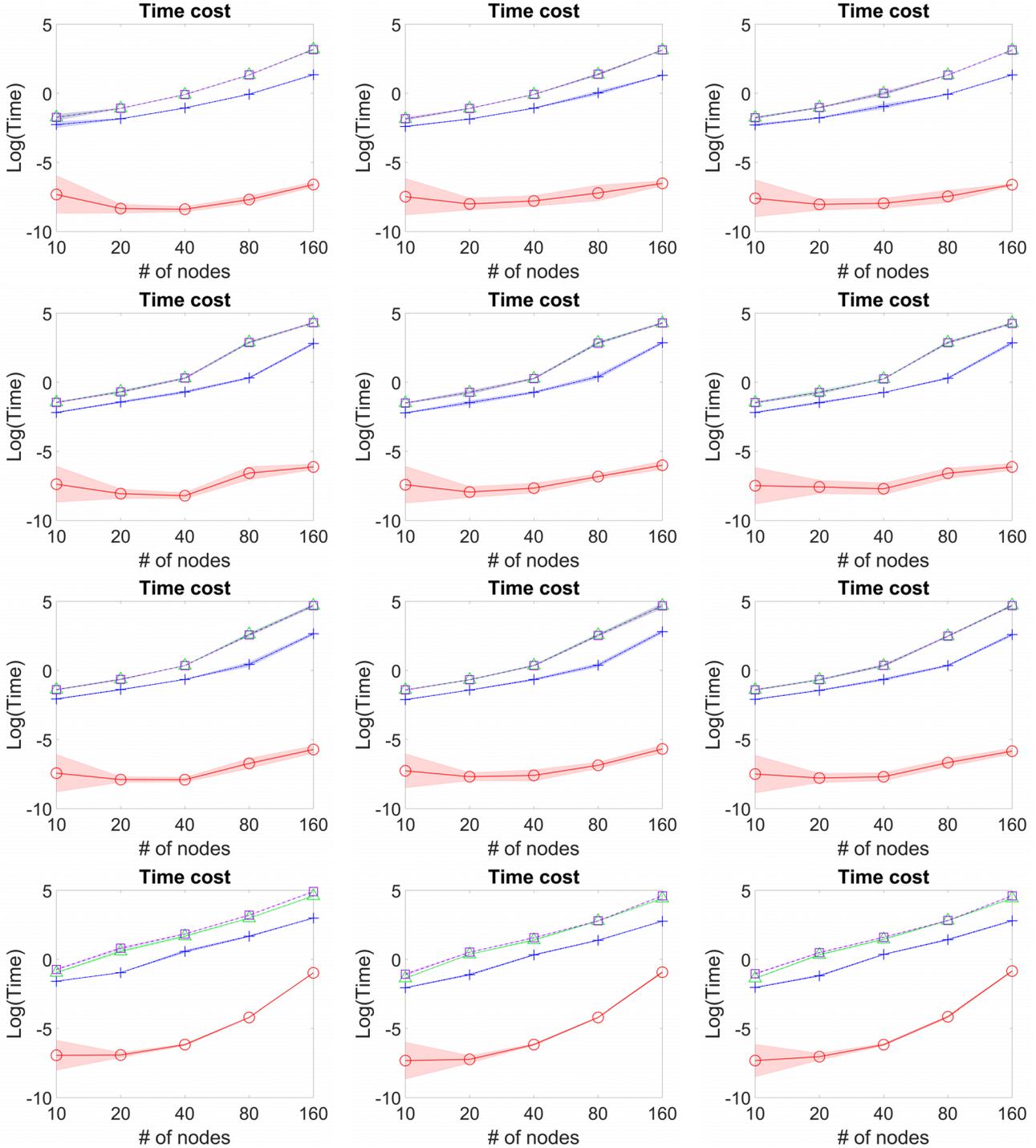


Fig 3: Time costs (in seconds) of all methods. Both axes are  $\log(e)$ -scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93]. We regarded  $N(0, 1)$  as zero time cost so does not appear in the time cost plot.

TABLE 3  
*Performance measures of 95% confidence intervals*  
 $n = 80, \rho_n \asymp 1$ , graphon: block model

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.957(0.202)	0.953(0.211)	0.956(0.205)	0.952(0.213)
	Length = 0.097(0.010)	0.040(0.008)	0.200(0.033)	0.145(0.033)
	LogTime = -8.448(0.110)	-7.214(0.083)	-7.165(0.082)	-7.180(0.353)
Norm. Approx.	0.950(0.218)	0.934(0.248)	0.942(0.235)	0.932(0.251)
	0.097(0.010)	0.040(0.008)	0.200(0.033)	0.145(0.033)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.842(0.365)	0.870(0.337)	0.852(0.355)	0.852(0.355)
	0.068(0.009)	0.031(0.007)	0.147(0.026)	0.113(0.025)
	-2.591(0.008)	-2.160(0.026)	-2.127(0.024)	-0.992(0.006)
Green and Shalizi [61]	0.938(0.241)	0.944(0.230)	0.934(0.249)	0.938(0.241)
	0.096(0.013)	0.044(0.010)	0.204(0.038)	0.150(0.037)
	-1.198(0.007)	0.499(0.032)	0.142(0.035)	0.383(0.010)
Levin and Levina [93]	0.942(0.234)	0.942(0.234)	0.942(0.234)	0.942(0.234)
	0.099(0.013)	0.043(0.010)	0.209(0.039)	0.155(0.038)
	-1.188(0.004)	0.507(0.028)	0.142(0.027)	0.489(0.004)

TABLE 4  
*Performance measures of 95% confidence intervals*  
 $n = 80, \rho_n \asymp 1$ , graphon: smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.958(0.201)	0.940(0.238)	0.951(0.216)	0.942(0.235)
	Length = 0.092(0.009)	0.021(0.005)	0.141(0.025)	0.083(0.021)
	LogTime = -8.225(0.113)	-7.363(0.066)	-7.278(0.086)	-6.974(0.541)
Norm. Approx.	0.951(0.216)	0.920(0.271)	0.938(0.242)	0.923(0.266)
	0.092(0.009)	0.021(0.005)	0.141(0.025)	0.083(0.021)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.816(0.388)	0.840(0.367)	0.826(0.379)	0.852(0.355)
	0.066(0.009)	0.018(0.005)	0.110(0.021)	0.072(0.018)
	-2.554(0.010)	-2.124(0.026)	-2.139(0.026)	-1.020(0.027)
Green and Shalizi [61]	0.928(0.259)	0.946(0.226)	0.938(0.241)	0.948(0.222)
	0.092(0.012)	0.025(0.007)	0.147(0.029)	0.090(0.024)
	-1.144(0.009)	0.497(0.042)	0.157(0.054)	0.334(0.025)
Levin and Levina [93]	0.948(0.222)	0.948(0.222)	0.950(0.218)	0.958(0.201)
	0.095(0.012)	0.024(0.007)	0.153(0.030)	0.095(0.025)
	-1.138(0.005)	0.507(0.031)	0.172(0.030)	0.447(0.019)

769 Materials. Our method exhibits very accurate actual coverage probabilities consistently close  
770 to the nominal confidence level. Our method is the only method that can always achieve a  
771  $\leq 0.010$  coverage error across all settings. It also produces competitively short confidence  
772 interval lengths, again, reflecting the high accuracy of the method. The comparison of com-  
773 putational efficiency between different methods echoes the qualitative results in Figure 3  
774 despite slightly different settings and confirms our method's huge speed advantage over all  
775 bootstrap methods.

776 It is interesting to observe that under the setting of this simulation, our empirical Edge-  
777 worth expansion method always produces the same interval length as the normal approxima-  
778 tion. This is not a coincidence in view of (4.5), (4.9) and that  $z_{\alpha/2}^2 = z_{1-\alpha/2}^2$ . In other words,  
779 as long as the studentization form  $\hat{T}_n$  that  $N(0, 1)$  approximates is equipped with the same  
780 variance estimator  $\hat{S}_n$  as our method, our two-sided Edgeworth confidence interval is a bias-  
781 corrected version (by mean-shift) of the corresponding ordinary CLT confidence interval.

TABLE 5  
*Performance measures of 95% confidence intervals*  
 $n = 80$ ,  $\rho_n \asymp 1$ , graphon: non-smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.956(0.205)	0.957(0.203)	0.957(0.202)	0.957(0.203)
	Length = 0.116(0.009)	0.135(0.010)	0.422(0.027)	0.531(0.040)
	LogTime = -8.291(0.076)	-7.345(0.103)	-7.817(0.153)	-7.045(0.373)
Norm. Approx.	0.952(0.215)	0.949(0.220)	0.951(0.215)	0.950(0.218)
	0.116(0.009)	0.135(0.010)	0.422(0.027)	0.531(0.040)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.830(0.376)	0.832(0.374)	0.830(0.376)	0.836(0.371)
	0.081(0.010)	0.096(0.011)	0.297(0.031)	0.379(0.041)
	-2.569(0.012)	-2.105(0.051)	-2.116(0.035)	-1.011(0.005)
Green and Shalizi [61]	0.940(0.238)	0.938(0.241)	0.944(0.230)	0.944(0.230)
	0.112(0.013)	0.135(0.014)	0.415(0.041)	0.529(0.055)
	-1.201(0.011)	0.547(0.075)	0.169(0.037)	0.328(0.015)
Levin and Levina [93]	0.954(0.210)	0.956(0.205)	0.956(0.205)	0.954(0.210)
	0.116(0.013)	0.138(0.013)	0.427(0.039)	0.544(0.052)
	-1.190(0.003)	0.534(0.049)	0.162(0.033)	0.436(0.014)

782 In the Supplemental Materials, we present simulation results for the remaining config-  
783 urations among  $n \in \{80, 160\}$  and  $\rho_n \asymp \{1, n^{-1/4}, n^{-1/2}, n^{-1}\}$ . For very sparse networks,  
784 our method and  $N(0, 1)$  approximation produce similar conservative confidence intervals for  
785 the  $R = \text{Edge}$ . On the other hand, all methods fail spectacularly for more complex motifs.  
786 Despite the required  $\rho_n$  lower bounds for all acyclic motifs are identically  $\omega(n^{-1})$  for our  
787 method and  $N(0, 1)$  approximation, the results are not surprising for two reasons: (i) the the-  
788 ory requires  $\rho_n \gg n^{-1}$ , so the simulation setting  $\rho_n \asymp n^{-1}$  is the boundary case and sensible  
789 outputs are not guaranteed; and (ii) the constant factor may matter a lot, and different acyclic  
790 motif shapes may require different minimum constants factor in  $\rho_n$  to show sensible results.

791 *5.3. Simulation 3: Numerical evaluation of the finite-sample impact of sparsity.* In this  
792 part, we conduct numerical studies to evaluate the finite sample performances of our method  
793 compared to benchmarks as the network grows sparser under fixed  $n$ . Despite in Simulation  
794 5.1, we tested different network sparsity settings (see Supplemental Material), it would still  
795 be interesting to more directly illustrate the impact of  $\rho_n$  for each fixed network size. The  
796 simulation set up carries over the same set of graphon models, motif shapes and compared  
797 methods from Simulation 5.1. Here, for simplicity, we only tested  $n = 80, 160$  and varied  $\rho_n$   
798 in a wider range of sparsity as follows:  $\{1$  (“dense”),  $n^{-1/4}, n^{-1/2}, n^{-1}\}$ .

799 Figure 4 shows the CDF approximation errors under different  $\rho_n$  settings for  $n = 160$ .  
800 Aligned with our theory’s prediction, we observe that as the network grows sparser, our  
801 method’s performance depreciates and gradually regresses to the performance of normal ap-  
802 proximation. Due to page limit, we sink the approximate error plots for  $n = 80$  and the time  
803 cost plots for both  $n$  settings to Supplemental Materials.

804 *5.4. Simulation 4: Degree-corrected stochastic block models.* We also tested our method  
805 on networks with degree heterogeneity. Our method maintains significant advantage in both  
806 accuracy and speed. Due to page limit, we sink all results and interpretation to Supplemental  
807 Materials (See Section 9.4).

808 *5.5. Simulation 5: Scalability of our method to large networks.* In this experiment, we  
809 test the scalability of our method. We find that all the three benchmark methods that we tested  
810 in previous simulations would fail to finish running on our high performance computing  
811 servers within 24 hours. Therefore, only the time costs of our method are recorded.

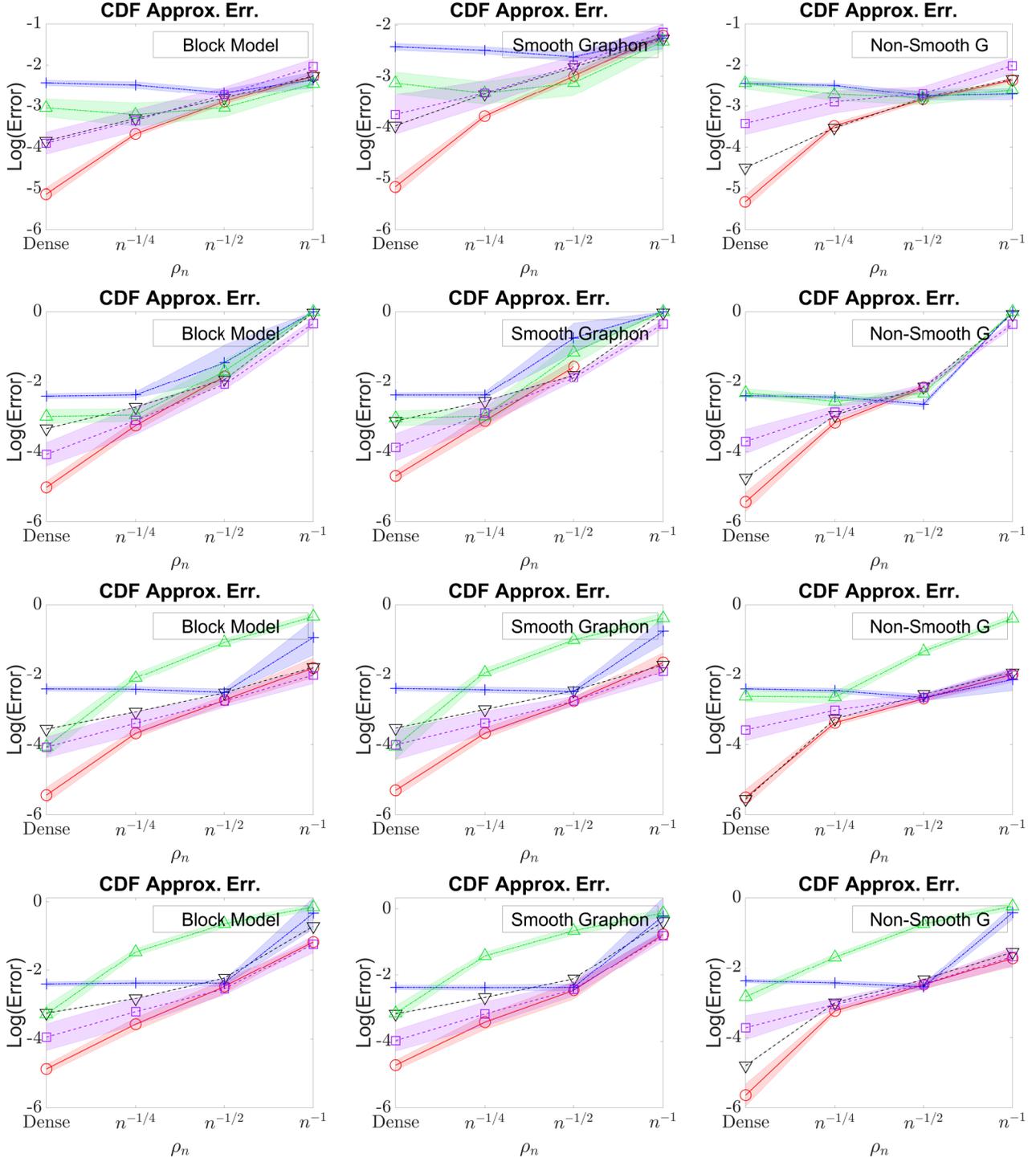


Fig 4: Impact of sparsity on approximation errors,  $n = 160$ . Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); black dashed curve marked down-triangle:  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

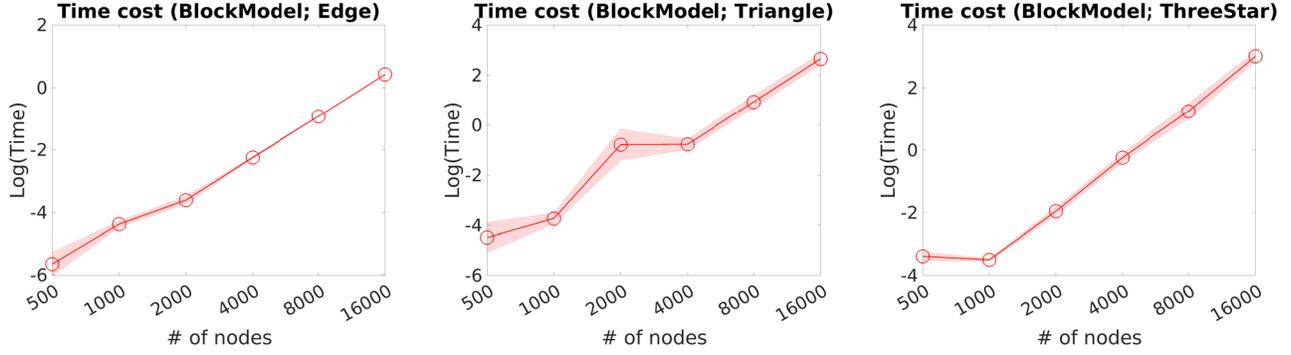


Fig 5: Scalability of our method on large networks. Bootstrap methods [17, 61, 93] with  $n_{\text{boot}} = 200$  bootstrap iterations did not finish in 24 hours, thus are not shown. In all experiments, our method took less than 20 seconds at the longest to run.

812 Figure 5 shows the results. Notably, our method shows a clear uniform slope in their  
 813 log-time cost growth rates for all the three motifs with  $r = 2, 3, 4$ , respectively. This echoes  
 814 our discussion in Section 6 that some “nicely shaped” motifs only cost  $O(n^2)$  or  $O(n^3)$   
 815 to count, regardless of motif size  $r$ . On the other hand, we recognize that counting a large and  
 816 “irregularly-shaped” motif could cost significantly more time.

**6. Discussion.** Our results do not cover the case where  $g_1(X_1)$  is lattice and  $\rho_n \asymp 1$ . An ad-hoc remedy is to simply introduce artificial missing links by sparsifying  $A$ :

$$\tilde{A}_{ij} := \tilde{A}_{ji} := \begin{cases} A_{ij} = A_{ji}, & \text{with probability } 1 - \tilde{\rho}_n \\ 0, & \text{with probability } \tilde{\rho}_n \end{cases}$$

817 where we set  $\tilde{\rho}_n \asymp (\log n)^{-1}$ . One can then make inferences about the population network  
 818 moment  $\tilde{\rho}_n \cdot \mu_n$  using  $\tilde{A}$  as the input data (notice  $\tilde{\rho}_n$  is known). This reinstates the  $(\log n)^{-1}$   
 819 sparsification that we need to overcome the latticeness at the price of a very minor information  
 820 loss.

The Edgeworth expansion we derived for Bernoulli  $A_{ij}|W_{ij}$  distributions can be readily extended to general weighted networks formulated by

$$A_{ij} := W_{ij} + \varepsilon_{ij},$$

821 where  $\varepsilon_{ij}$  may either depend on  $W_{ij}$  or not. A distinct feature of our setting is that the edge-  
 822 wise observational errors are a contributing component of  $\hat{T}_n$  that smooths the distribution.  
 823 In contrast to matrix estimation problems, where such noise is to be suppressed [33, 133], a  
 824 moderate amount of tailedness can strengthen the smoothing effect in  $A|W$  and might improve  
 825 finite sample performances. Notice that similar to [17, 61, 93], throughout this paper,  
 826 we work under the assumption inherited from well-known network analysis literature includ-  
 827 ing [19, 34, 56, 40] that  $\rho_n \cdot f(\cdot) \in [0, 1]$ , which also yields the boundedness of  $h(\cdot)$ . Thus,  
 828 the bounded-moment conditions in the classical literature of Edgeworth expansions for noise-  
 829 less U-statistics would be satisfied. There are at least two directions of potential relaxations:  
 830 relaxing the boundedness of the distribution of  $A_{ij}|W_{ij}$  and study a weighted network, or  
 831 consider unbounded graphons like that in [26]. The extension of our algorithm and analysis  
 832 to some light-tailed  $A_{ij}|W_{ij}$  distributions is straightforward. For instance, our proofs remain  
 833 valid for weighted network models with bounded graphon and an sub-exponential edge error  
 834 distribution, where  $W_{ij} \asymp \text{Var}(\varepsilon_{ij}|W_{ij}) \asymp \rho_n$ , by simply replacing Bernstein’s inequality by  
 835 generalized Hoeffding’s inequality (Theorem 1.2.2 in [126]).

836 On the other hand, in fact there is a simple universal strategy to handle heavy-tailed  $A_{ij}$   
 837 distributions, regardless of whether this is due to a heavy-tailed  $A_{ij}|W_{ij}$  distribution, or an  
 838 unbounded graphon such as  $f(x, y) = (xy)^{-\alpha}$  for  $\alpha \in (0, 1)$  in [26], or even both. As pointed  
 839 out in [111], we can perform a one-to-one transformation, such as the widely-used *sigmoid*  
 840 or *tanh* functions in machine learning, on each  $A_{ij}$ , tame it into a bounded  $\mathcal{T}(A_{ij})$ , and work  
 841 with the transformed data. This also guarantees that the population network moments of the  
 842 transformed network are always well-defined.

843 This paper focuses on studying the marginal randomness in  $A$  jointly contributed by the  
 844 randomness in  $W$  and  $A|W$ . In this study, we take the sparsity-scaled graphon  $\rho_n \cdot f$  as  
 845 the population and the graphon feature  $\mu_n$  as the ultimate inference goal. Our approach is  
 846 nonparametric and directly approximates  $F_{\hat{T}_n}$  without requiring a graphon estimation  $\widehat{W}$ . If  
 847 one regards  $W$  as the population and wants to make inference for  $U_n$ , she would need a CDF  
 848 approximation to  $(\widehat{U}_n - U_n)|X_1, \dots, X_n$ . This distribution is asymptotically normal as has  
 849 been described by (3.12) in our Lemma 3.1-(b). However, estimating the normal variance  
 850 typically *requires* a graphon estimation  $\widehat{W}$ <sup>30</sup>. Meanwhile, a practically meaningful graphon  
 851 estimation would typically require that  $f$  is smooth and/or low-rank, see [136, 134]. In other  
 852 words, the bootstrapping of  $\widehat{T}_n|X_1, \dots, X_n$  would (seemingly unavoidably) be a *parametric*  
 853 bootstrap. In view of Lemma 3.1-(b), asymptotically

$$(6.1) \quad (\rho_n \cdot n)^{1/2} \cdot \frac{\widehat{U}_n - U_n}{\sigma_n} \stackrel{d}{\approx} N(0, \sigma_w^2)$$

given  $X_1, \dots, X_n$ , where recall that  $\sigma_w \asymp 1$ . However, the minimax rate for sparse graphon  
 estimation (see [58, 85]) is

$$\text{Rescaled MSE: } (\rho_n \cdot n)^{-2} \cdot \|\widehat{W} - W\|_F^2 \asymp (\rho_n \cdot n)^{-1} \cdot \log n$$

854 If we use this error bound to control the estimation error of  $\sigma_w^2$ , then this yields an error of  
 855  $|\widehat{\sigma}_w^2 - \sigma_w^2| \asymp n^{-1} \|\widehat{W} - W\|_F \asymp \rho_n^{1/2} \cdot n^{-1/2} \cdot \log^{1/2} n$ . This error may dominate the  $n^{-1/2}$  cor-  
 856 rection term in an Edgeworth expansion even for dense networks (e.g., Cramer’s condition  
 857 holds and  $\rho_n \asymp 1$ ). Moreover, the minimax graphon estimation rate has not yet been achieved  
 858 by any polynomial-time algorithm (see [136, 134] for comments), and using a practically  
 859 feasible  $\widehat{W}$  would cause an error  $\gg n^{-1/2}$ , ignoring  $\rho_n$  and  $\log$ . Therefore, it might be chal-  
 860 lenging to accurately approximate the distribution of the LHS of (6.1) beyond asymptotic  
 861 normality. Our observation here echos the common practice in network bootstrap literature  
 862 [17, 61, 93] that they unanimously focus on the marginal distribution of  $\widehat{U}_n$ , rather than  
 863  $(\widehat{U}_n - U_n)|X_1, \dots, X_n$ <sup>31</sup>.

864 A retrospection on our simulation setting provides an interesting insight. In fact,  
 865 the population Edgeworth expansion provides a much more efficient Monte Carlo proce-  
 866 dure for simulating the true distribution  $F_{\hat{T}_n}$ . Indeed, estimating  $\xi_1$ ,  $\mathbb{E}[g_1^3(X_1)]$  and  
 867  $\mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$  with  $n_{\text{MC}} \asymp n$  Monte Carlo samples yields a CDF approxima-  
 868 tion rate of  $O(\mathcal{M}(\rho_n, n; R)) = o(n^{-1/2})$  when  $\rho_n$  satisfies the conditions of Theorem 3.1.  
 869 This is much more efficient than the empirical CDF, which requires  $n_{\text{MC}} \geq n^2$  to achieve the  
 870 same accuracy order.

In the application of our results, we focus on node sampling network bootstraps. It is an  
 interesting future work to investigate the higher-order accuracy properties of other schemes,

<sup>30</sup>The expression of  $\sigma_w^2$  contains  $W_{ij}^2$  terms originated from “ $W_{ij}(1 - W_{ij})$ ” terms, which could not be  
 estimated without a graphon estimation.

<sup>31</sup>For example, in Levin and Levina [93], the authors used a low-rank decomposition of  $A$ , which directly  
 leads to an estimated  $\widehat{W}$ . But they also solely focused the marginal distribution of  $U_n$  (in our notation system).

such as sub-graph sampling [17] and (artificially) weighted bootstrap [93]. Also comprehensive numerical comparisons of different schemes under various settings would certainly be interesting for practitioners. As a closely related point, this paper studies the *complete* noisy U-statistic, “complete” in the sense that  $(i_1, \dots, i_r)$  ranges over all  $\binom{n}{r}$  possible  $r$ -tuples. As one of the anonymous referees pointed out, evaluating the moment corresponding to an  $r$ -node motif would cost  $O(n^r)$ , which is still expensive for large  $n$ . Even for sparse networks, the counting may still need  $O(\rho_n^{r-1} \cdot n^r)$  time using cutting-edge algorithms, see Section III.A of [2]. To accelerate the computation, papers [22, 105, 36, 87, 121] investigated this topic for the conventional noiseless U-statistic setting and formulated the Edgeworth expansion for “incomplete” U-statistics. They study noiseless incomplete U-statistics, and [17] proposed a “subgraph subsampling” scheme (their scheme (a)) that computes noisy incomplete U-statistics, which we call  $\hat{U}_n^{(\text{Incomplete})}$  for the network setting. Formally, define

$$\hat{U}_n^{(\text{Incomplete})} := \frac{\sum_{1 \leq i_1 < \dots < i_r \leq n} I_{i_1, \dots, i_r} \cdot h(A_{i_1, \dots, i_r})}{\sum_{1 \leq i_1 < \dots < i_r \leq n} I_{i_1, \dots, i_r}}$$

871 where  $I_{i_1, \dots, i_r}$ ’s are random variables independent of the network data. These  $I_{i_1, \dots, i_r}$ ’s can  
 872 be i.i.d. Bernoulli, multinomial (if a given proportion of sub-sampled motifs is desired), or  
 873 other reasonable sampling scheme distributions. It would be an interesting future research to  
 874 carefully explore and quantify the self-smoothing effect for  $\hat{U}_n^{(\text{Incomplete})}$ .

875 On the other hand, however, some particular motifs, such as *cycles*, *stars* and *wheels*, can  
 876 be very efficiently evaluated, and the computational complexity may remain at most  $O(n^3)$ ,  
 877 instead of  $O(n^r)$ . For instance, the  $\hat{U}_n$  for *star* motifs can be approximately counted with  
 878 ignorable error in just  $O(n^2)$  time by averaging over  $\{A_{(i, \cdot)}^r\}_{i=1, \dots, n}$ . Another example is  
 879 that a  $(k, \ell)$ -*wheel* (see [20] for definition) can be evaluated in at most  $O(n^3)$  time using the  
 880 sample version of  $Q(R)$  in Equation (2.9) in [20]. More readings along this line include [98],  
 881 which provides detailed formula table for parallel computing up to  $r = 4$  motifs, and [4] that  
 882 studies fast-counting triangles in very large graphs. A recent paper [38] points out another  
 883 promising direction of distributed computation.

884 **Code.** The MATLAB code for our method (empirical Edgeworth expansion) is available  
 885 at <https://github.com/yzhanghf/NetworkEdgeworthExpansion>.

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## SUPPLEMENTARY MATERIAL

894 **Supplement for: “Edgeworth expansions for network moments”**  
 895 ([URL to be added](#)). The supplementary material contains: (1). Definition of  $\sigma_w$  in Lemma  
 896 3.1-(b); (2). All proofs; and (3). Additional simulation results and accompanying interpreta-  
 897 tions.

**SUPPLEMENTAL MATERIAL FOR:  
EDGEWORTH EXPANSIONS FOR NETWORK MOMENTS**

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903 **7. Definition of  $\sigma_w$  in Lemma 3.1-(b).** The formal definition of  $\sigma_w$  we present here  
904 would complete the statement of this lemma. To start, we express  $h(A_{i_1, \dots, i_r}) := \mathbb{1}_{[A_{i_1, \dots, i_r} \supseteq R]}$   
905 more explicitly as a sum of indicator product-terms, in which, each term checks if  $A_{i_1, \dots, i_r} \supseteq$   
906  $R_\pi$ , where the  $\supseteq$  is entry-wise and  $R_\pi$  is defined as  $(R_\pi)_{ij} := R_{\pi(i)\pi(j)}$  with  $\pi$  ranging over  
907 all permutations. To formalize this, let  $\text{Perm}(R) := \{\pi^{(\ell)}, \ell = 1, \dots, L\}$  denote the permutation  
908 group of  $R$ , where  $\pi^{(1)} = \text{id}$  is the identity map and  $\pi^{(\ell_1)}(R) \neq \pi^{(\ell_2)}(R)$  for any  $\ell_1 \neq \ell_2$ .  
909 For simplicity, for all  $1 \leq k_1 < k_2 \leq r$ , define

$$J^{(k_1, k_2)}(x) = \begin{cases} x & \text{if } R_{k_1 k_2} = 1 \\ 1 & \text{if } R_{k_1 k_2} = 0 \end{cases}$$

910 Then  $h(A_{i_1, \dots, i_r})$  can be formally represented as  
911

$$912 \quad h(A_{i_1, \dots, i_r}) = \sum_{\ell=1}^L \mathbb{1}_{[A_{i_1, \dots, i_r} \supseteq R_{\pi^{(\ell)}}]} = \sum_{\ell=1}^L \prod_{1 \leq k_1 < k_2 \leq r} J^{(\pi^{(\ell)}(k_1), \pi^{(\ell)}(k_2))}(A_{i_{k_1}, i_{k_2}})$$

913 Define

$$914 \quad (7.1) \quad \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} := J^{(\pi^{(\ell)}(j_1), \pi^{(\ell)}(j_2))}(W_{i_{j_1}, i_{j_2}})$$

$$915 \quad (7.2) \quad \mathfrak{S}_{j_1, j_2}^{(\ell)} := \text{Sign} \left\{ J^{(\pi^{(\ell)}(j_1), \pi^{(\ell)}(j_2))}(W_{i_{j_1}, i_{j_2}}) \right\}$$

where

$$\text{Sign}(J) := \frac{dJ(x)}{dx} = \begin{cases} +1 & \text{if } J(x) = x \\ 0 & \text{if } J(x) = 1 \end{cases}$$

and define

$$\hat{\Theta}_{ij} := \frac{2r(r-1)}{\sigma_n \cdot \binom{n-2}{r-2}} \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ \{i, j\} \subseteq \{i_1, \dots, i_r\}}} \sum_{\ell=1}^L \left\{ \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \neq (i, j)}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \right\} \cdot \mathfrak{S}_{\substack{j_1, j_2 \\ (i_{j_1}, i_{j_2}) = (i, j)}}^{(\ell)}$$

916 Define  $\sigma_w$  as follows

$$(7.3) \quad \sigma_w^2 := \frac{\rho_n \cdot n}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{\Theta}_{ij}^2 \cdot W_{ij}(1 - W_{ij})$$

917 This completes the statement of Lemma 3.1-(b).

918 **8. Proofs.**

919 8.1. *Bernstein-type concentration bound for multilinear polynomials of centered error*  
 920 *terms.* Our proof would need the main result Theorem 1.3 of [116], which is also used in  
 921 the proofs in [84]. To state this theorem, we first define a few preliminary quantities.

DEFINITION 8.1 (See Section 1.1 of [116]). *A hypergraph is formed by a node set*  
 $\mathcal{V}(H) := \{1, \dots, N\} = [N]$  *and a set*  $\mathcal{H}(H)$  *of hyperedges, where a hyperedge*  $\mathfrak{h}$  *of degree*  $q$  *is defined to be a subset of nodes*  $\mathcal{V}(\mathfrak{h}) \subset \mathcal{V}(H)$  *satisfying*  $|\mathcal{V}(\mathfrak{h})| \leq q$ . *We study the following*  
*multilinear polynomial*

$$f(\mathfrak{X}) := \sum_{\mathfrak{h} \in \mathcal{H}(H)} \mathfrak{W}_{\mathfrak{h}} \prod_{\mathfrak{v} \in \mathcal{V}(\mathfrak{h})} \mathfrak{X}_{\mathfrak{v}}$$

922 where  $\mathfrak{X} = (\mathfrak{X}_1, \dots, \mathfrak{X}_N)$  and  $\mathfrak{W}_{\mathfrak{h}}$  is an weight multiplier on each hyperedge  $\mathfrak{h}$ . Suppose on  
 923 each node we have a random variable,  $Y := \{Y_1, \dots, Y_N\}$  and a natural number  $\tau \geq 0$ . Let  
 924  $\mathfrak{W}$  denote the set of all edge weights. We define:

$$(8.1) \quad \Xi_{\tau} := \Xi_{\tau}(Y, H, \mathfrak{W}) := \max_{\mathcal{S} \subseteq [N]: |\mathcal{S}| = \tau} \left( \sum_{\mathfrak{h} \in \mathcal{H}(H): \mathcal{S} \subseteq \mathcal{V}(\mathfrak{h})} |\mathfrak{W}_{\mathfrak{h}}| \prod_{\mathfrak{v} \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|Y_{\mathfrak{v}}|] \right)$$

925 where to avoid symbol conflict we replaced “ $\mu$ ” in [116] by “ $\Xi$ ”.

926 Next we cite the main assumption.

DEFINITION 8.2. *A random variable*  $Z$  *is called central moment bounded with parameter*  $\mathcal{L} > 0$ , *if for any integer*  $i \geq 1$ , *we have*

$$\mathbb{E}[|Z - \mathbb{E}[Z]|^i] \leq i\mathcal{L} \cdot \mathbb{E}[|Z - \mathbb{E}[Z]|^{i-1}]$$

927 Now we are ready to cite their main result.

928 THEOREM 8.1 (Theorem 1.3 of [116], also cited as Lemma 15 in [84]). *Suppose all*  
 929  $Y_1, \dots, Y_N$  *are independent (but not necessarily identically distributed) and they all satisfy*  
 930 *the central moment bounded condition with a common parameter*  $\mathcal{L}$ . *Then we have*

$$(8.2) \quad \mathbb{P}(|f(Y) - \mathbb{E}[f(Y)]| \geq u) \leq e^2 \cdot \max \left\{ \exp \left\{ -\frac{u^2}{C^q \cdot \text{Var}(f(Y))} \right\}, \max_{1 \leq \tau \leq q} \exp \left\{ -\left( \frac{u}{\Xi_{\tau} \mathcal{L}^{\tau} C^q} \right)^{1/\tau} \right\} \right\}$$

933 where  $C$  is a universal constant.

934 **8.2. Proof of Lemma 3.1.**

935 8.2.1. *Proof of Lemma 3.1-(a).* By the decomposition in [96], we have

$$936 \quad \sigma_n^2 = \frac{r^2 \xi_1^2}{n} + O(\rho_n^{2s} \cdot n^{-2})$$

937 Therefore,  $\sigma_n \asymp n^{-1/2} \cdot \xi_1 \asymp \rho_n^s \cdot n^{-1/2}$ . Combining this fact with the Hoeffding's decompo-  
 938 sition of  $U_n - \mu_n$  in (3.1), we have

$$939 \quad \frac{U_n - \mu_n}{\sigma_n} = \frac{\frac{r}{n} \sum_{i=1}^n g_1(X_i) + \frac{r(r-1)}{n(n-1)} \sum_{1 \leq i < j \leq n} g_2(X_i, X_j) + \tilde{R}_{3:r}}{\frac{r\xi_1}{\sqrt{n}} + O(\rho_n^s \cdot n^{-3/2})}$$

940 where

$$941 \quad \tilde{R}_{3:r} := \binom{n}{3}^{-1} \binom{n-3}{r-3} \sum_{1 \leq i_1 < i_2 < i_3 \leq n} g_3(X_{i_1}, X_{i_2}, X_{i_3})$$

$$942 \quad + \sum_{k=4}^r \binom{n}{k}^{-1} \binom{n-k}{r-k} \sum_{1 \leq i_1 < \dots < i_k \leq n} g_k(X_{i_1}, \dots, X_{i_k})$$

$$943 \quad (8.3) \quad =: \tilde{R}_3 + \tilde{R}_{4:r}$$

944 and we also recall the definitions of  $U_n^\#$  and  $\Delta_n$  from (3.7) and the  $O(\rho_n^s \cdot n^{-3/2})$  remainder  
 945 control on the denominator is due to

$$946 \quad \sigma_n = \frac{r\xi_1}{\sqrt{n}} \sqrt{1 + O(n^{-1})} = \frac{r\xi_1}{\sqrt{n}} + O(\rho_n^s \cdot n^{-3/2}).$$

947 Recall that for simplicity, throughout this paper we assume  $f$  is bounded, which implies  
 948 the boundedness of the induced kernel function  $h(\cdot)$ . Therefore, the moment conditions of  
 949 Lemma 1 of [96] are satisfied. By Lemma 1 of [96], we know that  $\mathbb{E}[\tilde{R}_{4:r}] = O_p(\rho_n^s \cdot n^{-2})$ ,  
 950 thus by the remark below Lemma 2 in [96], this term is also  $\tilde{O}_p(\rho_n^s \cdot n^{-3/2})$ . Now for  $\tilde{R}_3$ ,  
 951 using Theorem 1 in [97], we know that  $\tilde{R}_3 = \tilde{O}_p(\rho_n^s \cdot n^{-3/2} \cdot \log^{3/2} n)$ . Therefore, we have

$$952 \quad \frac{U_n - \mu_n}{\sigma_n} = U_n^\# + \Delta_n(1 + O(n^{-1})) + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n)$$

$$953 \quad = U_n^\# + \Delta_n + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n)$$

954 This completes the proof of Lemma 3.1-(a).

955 Note that we use  $\tilde{R}$ ,  $\check{R}$  and  $\hat{R}$  to denote the remainder terms, where the ‘‘R’’ means ‘‘remain-  
 956 der’’. The properties of  $\tilde{R}$ ,  $\check{R}$  and  $\hat{R}$  certainly depend on the shape of the motif  $R$ , where we  
 957 inherit the tradition of using ‘‘R’’ to represent the motif from past network moment method  
 958 literature [20], but  $\tilde{R}$ ,  $\check{R}$  and  $\hat{R}$  are distinct notions from  $R$ .

959 8.2.2. *Proof of Lemma 3.1-(b).* We have

$$960 \quad \binom{n}{r} \cdot \hat{U}_n = \sum_{1 \leq i_1 < \dots < i_r \leq n} h(A_{i_1, \dots, i_r})$$

$$961 \quad = \sum_{1 \leq i_1 < \dots < i_r \leq n} \left\{ \sum_{\ell=1}^L \prod_{1 \leq j_1 < j_2 \leq r} \left( \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} + \mathfrak{S}_{j_1, j_2}^{(\ell)} \cdot \eta_{i_{j_1}, i_{j_2}} \right) \right\}$$

$$962 \quad =: \sum_{1 \leq k_1 < k_2 \leq n} \tilde{\Theta}_{k_1, k_2} \cdot \eta_{k_1, k_2} + \sum_{1 \leq i_1 < \dots < i_r \leq n} \sum_{\ell=1}^L \prod_{1 \leq j_1 < j_2 \leq r} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} + \hat{R}$$

$$963 \quad (8.4) \quad = \sum_{1 \leq k_1 < k_2 \leq n} \tilde{\Theta}_{k_1, k_2} \cdot \eta_{k_1, k_2} + \binom{n}{r} \cdot U_n + \hat{R},$$

964 where we denote

$$965 \quad \eta_{i,j} = A_{ij} - W_{ij},$$

$$966 \quad \tilde{\Theta}_{k_1, k_2} := \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ \{k_1, k_2\} \subseteq \{i_1, \dots, i_r\}}} \sum_{\ell=1}^L \left( \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \neq (k_1, k_2)}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \right) \mathfrak{S}_{\substack{j'_1, j'_2: \\ (i_{j'_1}, i_{j'_2}) = (k_1, k_2)}}^{(\ell)},$$

where we recall the definitions of  $\mathfrak{E}$  and  $\mathfrak{S}$  from (7.1) and (7.2), respectively, and  $\mathring{R} := \binom{n}{r} \sigma_n \cdot \check{R}_n$  is the remainder that contains all unmentioned terms. Referring to the later formal definition of  $\check{\Delta}_n$  in (8.19) and recalling the definition of  $\hat{\Delta}_n$  in (3.7), one can also easily verify that by definition  $\hat{\Delta}_n = \check{\Delta}_n + \mathring{R}_n$ . For clarity, we first verify that the coefficient in front of  $\eta_{k_1, k_2}$  is indeed  $\tilde{\Theta}_{k_1, k_2}$ . For each  $\{i_1, \dots, i_r\} : 1 \leq i_1 < \dots < i_r \leq n$  and each  $\ell$ , the term

$$\prod_{1 \leq j_1 < j_2 \leq r} \left( \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} + \mathfrak{S}_{j_1, j_2}^{(\ell)} \cdot \eta_{j_1, j_2} \right)$$

contributes to the coefficient of  $\eta_{k_1, k_2}$  if and only if  $\{k_1, k_2\} \subseteq \{i_1, \dots, i_r\}$ . Now if  $(j'_1, j'_2)$  is the index pair from  $\{1, \dots, r\}$  such that  $(i_{j'_1}, i_{j'_2}) = (k_1, k_2)$ , then itself contributes a multiplicative factor of  $\mathfrak{S}_{j'_1, j'_2}^{(\ell)}$  and every other pair  $(i_{j_1}, i_{j_2}) \neq (k_1, k_2)$  among  $\{1, \dots, r\}$  contributes a multiplicative factor of  $\mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)}$ , both into the term:

$$\left( \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \neq (k_1, k_2)}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \right) \mathfrak{S}_{\substack{j'_1, j'_2: \\ (i_{j'_1}, i_{j'_2}) = (k_1, k_2)}}^{(\ell)}$$

967 as an additive term in the expression of  $\tilde{\Theta}_{k_1, k_2}$ . This confirms that the coefficient of  $\eta_{k_1, k_2}$  is  
968 indeed  $\tilde{\Theta}_{k_1, k_2}$ .

969 The main content of this proof is to show the finite sample convergence rate of the linear  
970 part to its asymptotic distribution, and to bound the remainder  $\mathring{R}$ .

### 971 Concentration inequality for the remainder term $\mathring{R}$

In this part of the proof, our focus is to bound the remainder term  $\mathring{R}$ . Without loss of generality, we inspect the coefficient in front of the term

$$\eta_{(k_1^{(1)}, k_2^{(1)})} \cdots \eta_{(k_1^{(v)}, k_2^{(v)})}$$

972 where  $(k_1^{(1)}, k_2^{(1)}), \dots, (k_1^{(v)}, k_2^{(v)})$  are mutually different pairs from the set of node pairs  
973 formed by the first  $r$  indices  $\left\{ \left( \tilde{k}_1, \tilde{k}_2 \right) : \tilde{k}_1 < \tilde{k}_2, \{ \tilde{k}_1, \tilde{k}_2 \} \subseteq \{ i_1, \dots, i_r \} \right\}$ . This coefficient  
974 can be denoted and explicitly expanded as follows

$$975 \quad \tilde{\Theta}_{\mathcal{K}} := \{ (k_1^{(1)}, k_2^{(1)}), \dots, (k_1^{(v)}, k_2^{(v)}) \}$$

$$976 \quad := \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ (\cup_{j=1}^v \{k_1^{(j)}, k_2^{(j)}\}) \subset \{i_1, \dots, i_r\}}} \sum_{\ell=1}^L \left( \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \notin \mathcal{K}}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \right) \left( \prod_{\substack{(j'_1, j'_2): \\ (i_{j'_1}, i_{j'_2}) \in \mathcal{K}}} \mathfrak{S}_{j'_1, j'_2}^{(\ell)} \right)$$

$$977 \quad (8.5) \quad =: \sum_{\ell=1}^L \tilde{\Theta}_{\mathcal{K}}^{(\ell)}$$

Here we note a crucially important property of  $\tilde{\Theta}_{\mathcal{K}}^{(\ell)}$  that it is nonzero if and only if all of the node pairs in  $\mathcal{K}$  are edges in the  $\ell$ -th permuted version of the motif  $\pi^{(\ell)}(R)$ . This will be the key for us to effectively bound  $\tilde{\Theta}_{\mathcal{K}}$  and  $\hat{\Delta}^{(v,p)}$  in (8.7). We now upper bound  $\tilde{\Theta}_{(k_1^{(1)}, k_2^{(1)}), \dots, (k_1^{(v)}, k_2^{(v)})}$  for all  $v \geq 2$ , and this is an important step in upper bounding  $\mathring{R}$ . Define  $p$

$$p := \left| \{k_1^{(1)}, k_2^{(1)}\} \cup \dots \cup \{k_1^{(v)}, k_2^{(v)}\} \right|$$

978 to be the number of distinct indexes among  $(k_1^{(1)}, k_2^{(1)}), \dots, (k_1^{(v)}, k_2^{(v)})$ . Clearly, for  $v \geq 2$ ,  
979 we have

$$980 \quad 3 \leq p \leq r, \quad \text{and} \quad \frac{p}{2} \leq v \leq \begin{cases} p-1, & \text{for acyclic } R, \\ p(p-1)/2, & \text{for cyclic } R, \end{cases}$$

It suffices to bound inside part of the right hand side of (8.5) for each fixed set of indices  $\{i_1, \dots, i_r\}$  and  $\ell$ , because multiplying such upper bound by  $\binom{n-p}{r-p}$  gives an upper bound on  $\tilde{\Theta}_{\mathcal{K}}$ , ignoring constant factors including  $L$ . For each fixed  $\ell$  and given  $i_1, \dots, i_r$  and  $\mathcal{K}$ , we see that the number of  $(k_1^{(j)}, k_2^{(j)})$  that correspond to edges under the permutation mapping  $\pi^{(\ell)}$  must be  $v$ , otherwise at least one  $\mathfrak{S}$  term is zero and the summand is zero. So we have

$$\left| \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \notin \mathcal{K}}} \mathfrak{E}_{\{i_1, \dots, i_r, j_1, j_2\}}^{(\ell)} \right| = \rho_n^{s-v}$$

981 and consequently,

$$982 \quad (8.6) \quad \left| \tilde{\Theta}_{(k_1^{(1)}, k_2^{(1)}), \dots, (k_1^{(v)}, k_2^{(v)})} \right| \leq \rho_n^{s-v} \cdot \binom{n-p}{r-p} = \rho_n^{s-v} n^{r-p}$$

983 We can express the remainder term  $\mathring{R}$  in terms of  $\tilde{\Theta}$  and  $\eta$  terms. To facilitate detailed  
984 discussion and bounding, we group these terms. Define

$$985 \quad (8.7) \quad \hat{\Delta}^{(v,p)} := \sum_{\substack{\mathcal{K} \subseteq \{(k_1, k_2) : 1 \leq k_1 < k_2 \leq n\} \\ \text{Unique nodes}(\mathcal{K})=p \\ |\mathcal{K}|=v}} \left( \tilde{\Theta}_{\mathcal{K}} \prod_{(k_1, k_2) \in \mathcal{K}} \eta_{k_1, k_2} \right)$$

986 to be the collection of the terms in the remainder  $\mathring{R}$  corresponding to the product over  $v$   
987 different  $\eta$ -terms with exactly  $p$  unique participating nodes in these  $\eta$ -terms' indexes. Then

$$988 \quad (8.8) \quad \mathring{R} = \sum_{\substack{\text{All possible } (v,p) \\ v \geq 2, p \geq 3}} \hat{\Delta}^{(v,p)}$$

989 Obviously,  $v$ ,  $p$  and the total number of possible  $(v, p)$  pairs are all universally bounded,  
990 because the motif  $R$  is fixed. Therefore, in order to bound  $\mathring{R}$ , it suffices to bound  $\hat{\Delta}^{(v,p)}$  for  
991 every  $(v, p)$  pair. We need to bound not only the asymptotic magnitude of  $\hat{\Delta}^{(v,p)}$ , but also  
992 its tail probability. Notice that  $\hat{\Delta}^{(v,p)}$  is mean zero both conditional on  $W$  and marginally.

993 In order to bound its tail probability, it suffice to show a proper concentration inequality for  
 994  $\widehat{\Delta}^{(v,p)}$  conditional on  $W$ .

995 For this goal, we are going to apply Theorem 8.1, which derives a Bernstein inequality for  
 996 polynomials of independent centered random variables. Notice that  $\widehat{\Delta}^{(v,p)}$  can be rewritten  
 997 in the form of (8.1), where the nodes of the hypergraph  $\mathcal{V}(H) = \{(i, j) : 1 \leq i < j \leq n\}$  are  
 998 defined to be the following set of *node pairs* (notice that they are *not* nodes but node pairs in  
 999 our network). Define the set of hyperedges  $\mathcal{H}(H)$  as follows

$$\begin{aligned}
 1000 \quad \mathcal{H}(H) &:= \left\{ \mathcal{K} := \left\{ \left( k_1^{(1)}, k_2^{(1)} \right), \dots, \left( k_1^{(v)}, k_2^{(v)} \right) \right\} \subseteq \{(i, j) : 1 \leq i < j \leq n\} \right. \\
 1001 \quad &\text{s.t. } \left| \bigcup_{v'=1}^v \{k_1^{(v')}, k_2^{(v')}\} \right| = p, \text{ and} \\
 1002 \quad &\text{there exists } 1 \leq i_1 < \dots < i_r \leq n \text{ and } 1 \leq \ell \leq L, \text{ s.t.} \\
 1003 \quad &\mathcal{K} \in \{(i_{k'_1}, i_{k'_2}) : 1 \leq k'_1 < k'_2 \leq r\}, \\
 1004 \quad &\text{and } \left( \pi^{(\ell)}(R) \right)_{k'_1, k'_2} = 1, \text{ for all } (k'_1, k'_2) : (i_{k'_1}, i_{k'_2}) \in \mathcal{K} \left. \right\}
 \end{aligned}$$

1005 In other words, using the notation in Theorem 8.1,  $\mathcal{H}(H)$  is the collection of all size- $v$   
 1006 subsets of  $\mathcal{V}(H)$  that span across  $p$  nodes and are subset to some  $\ell$ th permuted version of  
 1007 the motif  $\pi^{(\ell)}(R)$ , edge weights being  $\mathfrak{W}_{\mathfrak{h}} = \tilde{\Theta}_{\mathcal{K}}$ , and each individual node-wise random  
 1008 variable is  $\{Y_{v'}\} := \left\{ \eta_{k_1^{(v')}, k_2^{(v')}} \right\}$ . Clearly, centered Bernoulli random variables satisfy the  
 1009 ‘‘bounded central moment’’ assumption with parameter  $\mathcal{L} = 1$ . In our context,  $q = v$ . In order  
 1010 to apply Theorem 8.1, now we bound the key quantities  $\Xi_1, \dots, \Xi_v$ . For each  $q' : 1 \leq q' \leq v$ ,  
 1011 bounding  $\Xi_{q'}$  consists of two sub-tasks:

- 1012 (i). Bounding  $\prod_{v' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{v'}|]$ ,  
 1013 (ii). Bounding  $\sum_{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{S} \subseteq V(\mathfrak{h})} |\mathfrak{W}_{\mathfrak{h}}| \cdot \prod_{v' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{v'}|]$ ,

1014 where in both bounds,  $\mathcal{S} \subseteq \mathcal{V}(H) : |\mathcal{S}| = q'$ . Bounding (i) is easy since it is just a product  
 1015 over  $v - q'$  independent  $\eta$  terms, each of which has an absolute expectation of  $\rho_n$ . We have

$$(8.9) \quad \prod_{v' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{v'}|] = \prod_{v' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[\mathbb{E}[|\eta_{v'}| | W]] \leq \rho_n^{v-q'}$$

1016 Now we bound (ii). This requires more detailed calculations to count the number of  $\tilde{\Theta}$  terms  
 1017 involved in the summation. It turns out the bound would differ for acyclic and cyclic motifs,  
 1018 which we discuss as follows.

- 1019 • When  $R$  is acyclic, in the summation  $\sum_{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{S} \subseteq V(\mathfrak{h})} |\mathfrak{W}_{\mathfrak{h}}|$ , we are summing over at most  
 1020  $p - q' - 1$  free indices. To see this fact, recall that in order for an individual summand to  
 1021 be nonzero, its corresponding hyperedge  $\mathfrak{h}$ , or equivalently, the corresponding  $\mathcal{K}$ , must be  
 1022 a subset of some permuted version of the motif  $R$ . Therefore the requirement that it must  
 1023 contain  $\mathcal{S} : |\mathcal{S}| = q'$  would pin down at least  $q' + 1$  indices, leaving us at most  $p - q' - 1$   
 1024 free indices. Therefore, recalling that  $|\mathfrak{W}_{\mathfrak{h}}| := |\tilde{\Theta}_{\mathcal{K}}| \leq \rho_n^{s-v} n^{r-p}$ , we obtain the following  
 1025 bound for (1)

$$\begin{aligned}
 1026 \quad &\sum_{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{S} \subseteq V(\mathfrak{h})} |\mathfrak{W}_{\mathfrak{h}}| \cdot \prod_{v' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{v'}|] \leq \binom{n}{p - q' - 1} \cdot \rho_n^{s-v} n^{r-p} \cdot \rho_n^{v-q'} \\
 1027 \quad (8.10) \quad &\leq \rho_n^{s-q'} \cdot n^{r-q'-1}
 \end{aligned}$$

1028 Since (8.10) holds for any  $\mathcal{S} : |\mathcal{S}| = q'$ , by the definition of  $\Xi_{q'}$ , we have

$$1029 \quad (8.11) \quad \frac{\Xi_{q'}}{\binom{n}{r} \cdot \sigma_n} \leq \frac{\rho_n^{s-q'} \cdot n^{r-q'-1}}{\rho_n^s \cdot n^{r-1/2}} = (\rho_n \cdot n)^{-q'} \cdot n^{-1/2}$$

1030 Notice that under the weak sparsity assumption  $\rho_n = \omega(n^{-1})$  for acyclic  $R$ , the RHS of  
1031 (8.11) is decreasing in  $q'$ . The interpretation of the result (8.11) says that in fact, our choice  
1032 of “ $u$ ” in the second term inside “max” in Theorem 8.1 for  $\tau : 1 \leq \tau \leq q = v$  is bottlenecked  
1033 by the case  $\tau = 1$ .

- Then we discuss the more complicated case that  $R$  is cyclic. Now, we consider those  $\mathcal{S} \subset \mathcal{V}(H)$  whose numbers of unique nodes are  $q'' + 1$  for some  $q'' \in \{2, \dots, p-1\}$  ( $q'' = 1$  cannot form a cyclic  $R$ ). For such  $\mathcal{S}$ , we have

$$\sum_{\substack{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{S} \subseteq V(\mathfrak{h}) \\ \text{Unique nodes}(\mathcal{S}) = q'' + 1}} |\mathfrak{W}_{\mathfrak{h}}| \leq \binom{n}{p - q'' - 1} \cdot \rho_n^{s-v} n^{r-p} \leq \rho_n^{s-v} \cdot n^{r-q''-1}$$

since we have  $p - q'' - 1$  free indices to sum over. Meanwhile, regardless of the number of unique nodes in  $\mathcal{S}$ , we always have

$$\prod_{\mathfrak{v}' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{\mathfrak{v}'}|] \leq \rho_n^{v-|\mathcal{S}|},$$

1034 Now using the simple relationship  $|\mathcal{S}| \leq q''(q'' + 1)/2$ , we have

$$\begin{aligned} 1035 & \sum_{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{S} \subseteq V(\mathfrak{h})} |\mathfrak{W}_{\mathfrak{h}}| \cdot \prod_{\mathfrak{v}' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{\mathfrak{v}'}|] \\ 1036 &= \sum_{\text{All possible } q''} \sum_{\substack{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{S} \subseteq V(\mathfrak{h}) \\ \text{Unique nodes}(\mathcal{S}) = q'' + 1}} |\mathfrak{W}_{\mathfrak{h}}| \prod_{\mathfrak{v}' \in \mathcal{V}(\mathfrak{h}) \setminus \mathcal{S}} \mathbb{E}[|\eta_{\mathfrak{v}'}|] \\ 1037 &\leq \sum_{q''} \binom{n}{p - q'' - 1} \cdot \rho_n^{s-v} n^{r-p} \cdot \rho_n^{v-q''(q''+1)/2} \\ 1038 \quad (8.12) &\leq \sum_{q''} \rho_n^{s-q''(q''+1)/2} \cdot n^{r-q''-1} \end{aligned}$$

1039 Therefore, we have

$$\begin{aligned} 1040 & \frac{\Xi_{q'}}{\binom{n}{r} \cdot \sigma_n} \leq \max_{q'' : q''(q''+1)/2 \geq q'} \frac{\rho_n^{s-q''(q''+1)/2} \cdot n^{r-q''-1}}{\rho_n^s \cdot n^{r-1/2}} \\ 1041 \quad (8.13) &= \max_{q'' : q''(q''+1)/2 \geq q'} (\rho_n^{-(q''+1)/2} \cdot n^{-1})^{q''} \cdot n^{-1/2} \end{aligned}$$

1042 Recall that by definition  $q'' \leq p - 1 \leq r - 1$ . Under the weak sparsity assumption that  
1043  $\rho_n = \omega(n^{-2/r})$ , we know that  $\rho_n^{-(q''+1)/2} \cdot n^{-1} \ll 1$ , so the maximum asymptotic order on  
1044 the RHS of (8.13) is achieved at the minimum possible  $q''$  value of 2.

1045 Now we have bounded the  $\Xi$  terms. In fact, as we will see, in Theorem 8.1, the concentra-  
1046 tion error bound terms due to  $\Xi$ 's are dominated by the term due to variance. In order to  
1047 apply Theorem 8.1, it only remains to bound  $\text{Var}(\widehat{\Delta}^{(v,p)})$ . We shall do this by bounding  
1048  $\text{Var}(\widehat{\Delta}^{(v,p)}|W)$  for each individual  $(v, p)$ , since the number of such terms is a fixed number.

1049 We have

$$\begin{aligned}
 1050 \quad \text{Var} \left( \widehat{\Delta}^{(v,p)} | W \right) &= \sum_{1 \leq i_1 < \dots < i_p \leq n} \tilde{\Theta}_{\mathcal{K}}^2 \cdot \text{Var} \left( \prod_{(k_1, k_2) \in \mathcal{K}} \eta_{k_1, k_2} | W \right) \\
 1051 \quad (8.14) \quad &\leq n^p \cdot \rho_n^{2s-2v} \cdot n^{2r-2p} \cdot \rho_n^v = \rho_n^{2s-v} \cdot n^{2r-p} \leq \rho_n^{2s-v} \cdot n^{2r-p}
 \end{aligned}$$

1052 where we used (8.6). Since  $v \leq s$  and  $p \geq 3$ , this yields the following upper bound.

$$\begin{aligned}
 1053 \quad \frac{\left\{ \text{Var} \left( \widehat{\Delta}^{(v,p)} | W \right) \right\}^{1/2}}{\binom{n}{r} \cdot \sigma_n} &\preceq \rho_n^{-s} \cdot n^{1/2-r} \cdot \left\{ \text{Var} \left( \widehat{\Delta}^{(v,p)} | W \right) \right\}^{1/2} \\
 1054 &\leq \left( \rho_n^{-s} \cdot n^{1/2-r} \right) \cdot \left( \rho_n^{s-v/2} \cdot n^{r-p/2} \right) \\
 1055 \quad (8.15) \quad &= \rho_n^{-v/2} \cdot n^{-(p-1)/2}
 \end{aligned}$$

1056 Next we discuss different upper bounds of the RHS of (8.15) based on different motif  $R$   
1057 shapes.

1058 • **Case 1:** if  $R$  is acyclic, we have  $v \leq p-1$ . Combining this with the fact that  $p \geq 3$  and  
1059 Assumption (ii) of Lemma 3.1 that  $\rho_n = \omega(n^{-1})$ , we have

$$1060 \quad (8.16) \quad \rho_n^{-v/2} \cdot n^{-(p-1)/2} \leq (\rho_n \cdot n)^{-(p-1)/2} \leq (\rho_n \cdot n)^{-1}$$

1061 • **Case 2:** if  $R$  is cyclic, we have  $v \leq p(p-1)/2$ . Combining this with the fact that  $3 \leq p \leq r$   
1062 and Assumption (ii) of Lemma 3.1 that  $\rho_n = \omega(n^{-2/r})$ , we have

$$\begin{aligned}
 1063 \quad \rho_n^{-v/2} \cdot n^{-(p-1)/2} &\leq (\rho_n^{-p(p-1)/2} \cdot n^{-(p-1)})^{1/2} \\
 1064 \quad (8.17) \quad &= \left( \rho_n^{-p/2} \cdot n^{-1} \right)^{(p-1)/2} \leq \rho_n^{-r/2} \cdot n^{-1}
 \end{aligned}$$

1065 Repeating this argument for every  $(v, p)$  pair, and plug (8.11), (8.13), (8.16) and (8.17)  
1066 back into Theorem 8.1, we have

$$\begin{aligned}
 1067 \quad (8.18) \quad \mathbb{P} \left( \check{R}_n := \frac{\check{R}}{\binom{n}{r} \cdot \sigma_n} \geq C \cdot \mathcal{M}(\rho_n, n; R) \right) \\
 1068 \quad \leq \begin{cases} \max \left\{ \exp \left( -\frac{((\rho_n \cdot n)^{-1} \log^{1/2} n)^2}{(\rho_n \cdot n)^{-2}} \right), \exp \left( -\frac{(\rho_n \cdot n)^{-1} \log^{1/2} n}{(\rho_n \cdot n)^{-1} \cdot n^{-1/2}} \right) \right\}, & \text{for acyclic } R; \\ \max \left\{ \exp \left( -\frac{((\rho_n^{-r/2} \cdot n^{-1}) \log^{1/2} n)^2}{(\rho_n^{-r/2} \cdot n^{-1})^2} \right), \exp \left( -\frac{(\rho_n^{-r/2} \cdot n^{-1}) \log^{1/2} n}{\rho_n^{-3} n^{-5/2}} \right) \right\}, & \text{for cyclic } R; \end{cases} \\
 1069 \quad = O(n^{-1})
 \end{aligned}$$

1070 for a large enough universal constant  $C$ .

### 1071 Asymptotic normality of the linear part $\check{\Delta}_n$ and Berry-Esseen bound

1072 Now, we focus on  $\check{\Delta}_n$ , the linear part of  $(\widehat{U}_n - U_n)/\sigma_n$  and show the uniform rate of its  
1073 normal approximation. Recalling the definitions of  $\check{\Delta}_n$ ,  $\tilde{\Theta}_{ij}$  and  $\hat{\Theta}_{ij}$ , ignoring the remainder  
1074 term, we have

$$1075 \quad \check{\Delta}_n := \text{Linear part of} \left( \frac{\widehat{U}_n - U_n}{\sigma_n} \right) = \frac{1}{\binom{n}{r} \cdot \sigma_n} \sum_{1 \leq i < j \leq n} \tilde{\Theta}_{ij} \cdot \eta_{ij}$$

$$(8.19) \quad =: \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{\Theta}_{ij} \cdot \eta_{ij}$$

We are going to show the asymptotic normality of  $\check{\Delta}_n$  and its concentration speed by applying the Berry-Esseen bound for independent but differently-distributed random variables [39] conditioning on  $W$ . In this derivation, the key terms are the asymptotic orders of the second and third central moments of each individual  $\hat{\Theta}_{ij}\eta_{ij}$  term. We first show that with respect to the randomness in  $W$ , we have  $\hat{\Theta}_{ij} \asymp \rho_n^{-1} \cdot n^{1/2}$ . Then when we condition on  $W$  and apply the generalized Berry-Esseen bound with respect to the randomness of  $A$  given  $W$ , we can think of  $\hat{\Theta}_{ij}$  as its asymptotic order  $\rho_n^{-1} \cdot n^{1/2}$ . Recall that

$$\prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \neq (i, j)}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \cdot \mathfrak{S}_{\substack{j'_1, j'_2: \\ (i_{j'_1}, i_{j'_2}) = (i, j)}}^{(\ell)} \asymp \begin{cases} \rho_n^{s-1}, & \text{if } \mathfrak{S}_{\substack{j'_1, j'_2: \\ (i_{j'_1}, i_{j'_2}) = (i, j)}}^{(\ell)} = 1 \text{ or equivalently } (\pi^{(\ell)}(R))_{j'_1, j'_2} = 1 \\ 0, & \text{if } \mathfrak{S}_{\substack{j'_1, j'_2: \\ (i_{j'_1}, i_{j'_2}) = (i, j)}}^{(\ell)} = 0 \text{ or equivalently } (\pi^{(\ell)}(R))_{j'_1, j'_2} = 0 \end{cases}$$

We have

$$\sigma_n \cdot \hat{\Theta}_{ij} \asymp \rho_n^{s-1}$$

This is because

$$(8.20) \quad \frac{\sigma_n \cdot \hat{\Theta}_{ij}}{2r(r-1)} = \frac{1}{\binom{n-2}{r-2}} \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ \{i, j\} \subset \{1, \dots, i_r\}}} \sum_{\ell=1}^L \left\{ \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (j_1, j_2) \neq (i, j)}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \right\} \cdot \mathfrak{S}_{i, j}^{(\ell)}$$

Since for each given  $\{i_1, \dots, i_r\}$  that contains  $\{i, j\}$ , the summation over  $\ell$  ranges among all  $\pi^{(\ell)}$  that keep  $(i, j)$  an edge in  $\pi^{(\ell)}(R)$ , so the outcome of this summation over  $\ell$  is symmetric in  $\{i_1, \dots, i_r\} \setminus \{i, j\}$ . Consequently,  $\hat{\Theta}_{ij}$  is also symmetric in  $\{1, 2, \dots, n\} \setminus \{i, j\}$ . Applying Hoeffding's decomposition to each  $\hat{\Theta}_{ij}$  viewed as a U-statistic with index set  $\{1, \dots, n\} \setminus \{i, j\}$  and using [97] to bound the remainder, we have

$$(8.21) \quad \frac{\sigma_n \cdot \hat{\Theta}_{ij}}{2r(r-1)} = \frac{\mathbb{E}[\sigma_n \cdot \hat{\Theta}_{ij} | X_i, X_j]}{2r(r-1)} + \frac{r-2}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i, j}} \check{g}_{1; i, j}(X_k) + \tilde{O}_p(\rho_n^{s-1} \cdot n^{-1} \cdot \log^{3/2} n)$$

where

$$\check{g}_{1; i, j}(X_k) := \mathbb{E} \left[ \prod_{\substack{1 \leq j_1 < j_2 \leq r \\ (i_{j_1}, i_{j_2}) \neq (i, j)}} \mathfrak{E}_{\{i_1, \dots, i_r\}, j_1, j_2}^{(\ell)} \middle| X_k, X_i, X_j \right] - \frac{\mathbb{E}[\sigma_n \cdot \hat{\Theta}_{ij} | X_i, X_j]}{2r(r-1)}$$

where the indexes  $i_1, \dots, i_r$  satisfy  $\{i, j, k\} \subseteq \{i_1, \dots, i_r\} \subseteq \{1, \dots, n\}$ . Since the linear part of a Hoeffding's decomposition are averaging over  $n$  i.i.d. terms with  $\mathbb{E}[\check{g}_{1; i, j}(X_k) | X_i, X_j] = 0$ ,  $|\check{g}_{1; i, j}(X_k)| = O(\rho_n^{s-1})$  a.s. and  $\text{Var}(\check{g}_{1; i, j}(X_k) | X_i, X_j) \leq \rho_n^{s-1}$ , by Bernstein's inequality combined with a union bound, we have

$$\mathbb{P} \left( \max_{1 \leq i < j \leq n} \frac{\sigma_n \cdot \left| \hat{\Theta}_{ij} - \mathbb{E}[\hat{\Theta}_{ij}] \right|}{2r(r-1)} \geq \rho_n^{s-1} \cdot t \middle| X_i, X_j \right) \leq C_1 \binom{n}{2} \cdot \left\{ e^{-C_2 n t^2} + e^{-C_3 n t} \right\}$$

1092 which yields that conditioning on  $X_i, X_j$ , we have

$$(8.22) \quad \max_{1 \leq i < j \leq n} \frac{\sigma_n \cdot \left| \hat{\Theta}_{ij} - \mathbb{E} \left[ \hat{\Theta}_{ij} \right] \right|}{2r(r-1)} = \tilde{O}_p \left( \rho_n^{s-1} \cdot n^{-1/2} \cdot \log n \right)$$

Since

$$\rho_n^{-(s-1)} \cdot \mathbb{E} \left[ \sigma_n \cdot \hat{\Theta}_{ij} \right] \asymp C > 0$$

1093 for a universal constant  $C$ , when discussing the concentration of  $\check{\Delta}_n$ , it suffices to prove  
 1094 the Berry-Esseen bound for the asymptotic normality of  $\check{\Delta}_n$  with respect to the random-  
 1095 ness in  $A|W$ , conditioning on a “nicely-behaved”  $W$  such that  $C/2 < \rho_n^{-(s-1)} \sigma_n \cdot \hat{\Theta}_{ij} \asymp$   
 1096  $\rho_n^{-(s-1)} \sigma_n \cdot \hat{\Theta}_{ij} < 3C/2$  holds for all  $1 \leq i < j \leq n$  simultaneously, because the probability  
 1097 that  $W$  behaves “badly” is exponentially small and ignorable. We write

$$(8.23) \quad \frac{(\rho_n \cdot n)^{1/2} \cdot \check{\Delta}_n}{\sigma_w} = \sum_{1 \leq i < j \leq n} \frac{(\rho_n \cdot n)^{1/2} \cdot \hat{\Theta}_{ij}}{\sigma_w \cdot \binom{n}{2}} \cdot \eta_{ij}$$

1098 where we notice that each individual coefficient in front of  $\eta_{ij}$  is at the order of  $\rho_n^{-1/2} \cdot n^{-1}$ .  
 1099 Using Theorem 2.1 of [39]

$$(8.24) \quad \left\| \frac{F_{\left( \frac{(\rho_n \cdot n)^{1/2} \cdot \check{\Delta}_n}{\sigma_w} \right) \Big| W} (u) - F_{N(0,1)}(u)}{\infty} \right\| \leq C \left\{ 0 + \sum_{1 \leq i < j \leq n} \left( \frac{(\rho_n \cdot n)^{1/2} \cdot \hat{\Theta}_{ij}}{\sigma_w \cdot \binom{n}{2}} \right)^3 \mathbb{E} \left[ |\eta_{ij}|^3 \right] \Big| W \right\}$$

$$\leq n^2 \cdot \rho_n^{-3/2} \cdot n^{-3} \cdot \rho_n \asymp \rho_n^{-1/2} \cdot n^{-1}$$

1102 where we used

$$1103 \quad \mathbb{E} \left[ |\eta_{ij}|^3 \mid W \right] = W_{ij}(1 - W_{ij})^3 + (1 - W_{ij})W_{ij}^3 \leq 2W_{ij} \asymp \rho_n$$

1104 Recall that the above result was obtained under “nicely-bahaved”  $W$ , but the probability of  
 1105 “bad”  $W$  is exponentially small. Therefore, we have

$$(8.25) \quad \left\| \frac{F_{\left( \frac{(\rho_n \cdot n)^{1/2} \cdot \check{\Delta}_n}{\sigma_w} \right) \Big| W} (u) - F_{N(0,1)}(u)}{\infty} \right\| = \tilde{O}_p(\rho_n^{-1/2} \cdot n^{-1})$$

1108 Combining (8.25) and (8.18) with Lemma 8.2 finishes the proof of Lemma 3.1-(b).

1109 8.2.3. *Proof of Lemma 3.1-(c).* Define the following shorthand that will be used in not  
 1110 only this proof but also others

$$1111 \quad (8.26) \quad \hat{a}_i := \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(A_{i, i_1, \dots, i_{r-1}})$$

$$1112 \quad (8.27) \quad a_i := \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(W_{i, i_1, \dots, i_{r-1}})$$

$$1113 \quad = \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(X_i, X_{i_1}, \dots, X_{i_{r-1}})$$

1114 A simple but useful property is as follows:

$$(8.28) \quad \frac{1}{n} \sum_{i=1}^n \hat{a}_i = \hat{U}_n \quad \text{and} \quad \frac{1}{n} \sum_{i=1}^n a_i = U_n$$

1115 To see (8.28), notice that

$$\sum_{i=1}^n \hat{a}_i \cdot \binom{n-1}{r-1} = r \sum_{1 \leq i_1 < \dots < i_r \leq n} h(A_{i_1, \dots, i_r}) = r \cdot \binom{n}{r} \hat{U}_n$$

1116 because each  $h(A_{i_1, \dots, i_r})$  is counted  $r$  times by  $\hat{a}_{i_1}, \dots, \hat{a}_{i_r}$ , respectively, on the LHS. The  
1117 relationship between  $a_i$  and  $U_n$  is verified exactly similarly.

Next, we start to decompose  $\hat{\delta}_n$ . By definition, we have

$$\hat{\delta}_n = \frac{\hat{S}_n^2 - \hat{\sigma}_n^2}{\sigma_n^2} = \frac{n\hat{S}_n^2 - n\hat{\sigma}_n^2}{\frac{n\sigma_n^2}{r^2}}$$

1118 in which,

$$\begin{aligned} 1119 \quad \frac{n\hat{S}_n^2}{r^2} &= \frac{1}{n} \sum_{i=1}^n (\hat{a}_i - \hat{U}_n)^2 = \frac{1}{n} \sum_{i=1}^n \{(\hat{a}_i - U_n) + (U_n - \hat{U}_n)\}^2 \\ 1120 \quad &= \frac{1}{n} \sum_{i=1}^n (\hat{a}_i - U_n)^2 + \frac{2}{n} \sum_{i=1}^n (\hat{a}_i - U_n) (U_n - \hat{U}_n) + (U_n - \hat{U}_n)^2 \\ 1121 \quad (8.29) \quad &= \frac{1}{n} \sum_{i=1}^n (\hat{a}_i - U_n)^2 - (U_n - \hat{U}_n)^2 \end{aligned}$$

1122 By the earlier proof steps, we know that

$$(8.30) \quad (U_n - \hat{U}_n)^2 = O_p(\rho_n^{2s-1} n^{-2})$$

1123 According to the remark under Lemma 2 in [96], this term is  $\tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1})$  and thus  
1124 ignorable. We focus on decomposing the first term on the RHS of (8.29). We have

$$\begin{aligned} 1125 \quad \frac{1}{n} \sum_{i=1}^n (\hat{a}_i - U_n)^2 &= \frac{1}{n} \sum_{i=1}^n \{(\hat{a}_i - a_i) + (a_i - U_n)\}^2 \\ 1126 \quad (8.31) \quad &= \frac{1}{n} \sum_{i=1}^n (\hat{a}_i - a_i)^2 + \frac{2}{n} \sum_{i=1}^n (\hat{a}_i - a_i) (a_i - U_n) + \frac{1}{n} \sum_{i=1}^n (a_i - U_n)^2 \end{aligned}$$

1127 Term 3 on the RHS of (8.31) is the constituting part of  $\hat{\sigma}_n^2$ , so we only need to bound the  
1128 first two terms. The key component is to study  $\hat{a}_i - a_i$ . Similar to the proof of part (b), starting  
1129 from re-expressing the definition of  $\hat{a}_i$  and  $a_i$ , we have

$$\begin{aligned} 1130 \quad \hat{a}_i - a_i &= \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ i \in \{i_1, \dots, i_r\}}} \{h(A_{i_1, \dots, i_r}) - h(W_{i_1, \dots, i_r})\} \\ 1131 \quad (8.32) \quad &= \frac{1}{\binom{n-1}{r-1}} \sum_{\text{All possible } (v,p)} \hat{\Delta}^{(i;v,p)} \end{aligned}$$

1132 where recall that we use the shorthand  $\mathcal{K} := \{(k_1^{(1)}, k_2^{(1)}), \dots, (k_1^{(v)}, k_2^{(v)})\}$ , and define

(8.33)

$$1133 \quad \widehat{\Delta}^{(i;v,p)} := \sum_{\substack{\mathcal{K} \subseteq \{(k_1, k_2) : 1 \leq k_1 < k_2 \leq n\} \\ \text{Unique nodes}(\mathcal{K}) = p \\ |\mathcal{K}| = v}} \widetilde{\Theta}_{\mathcal{K}}^{(i)} \prod_{(k_1, k_2) \in \mathcal{K}} \eta_{k_1, k_2}$$

(8.34)

$$1134 \quad \widetilde{\Theta}_{\mathcal{K}}^{(i)} := \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ i \in \{i_1, \dots, i_r\} \\ \mathcal{K} \subseteq \{(i_{j_1}, i_{j_2}) : 1 \leq j_1 < j_2 \leq r\}}} \sum_{\ell=1}^L \left( \prod_{1 \leq j_1 < j_2 \leq r} \mathfrak{E}_{\{i_1, \dots, i_r\}; j_1, j_2}^{(\ell)} \right) \left( \prod_{(j'_1, j'_2) : (i_{j'_1}, i_{j'_2}) \in \mathcal{K}} \mathfrak{E}_{j'_1, j'_2}^{(\ell)} \right)$$

1135 Here we stress a crucial point that although in these definitions we always have  $i \in$   
 1136  $\{i_1, \dots, i_r\}$ , the node  $i$ , however, might or might not appear in  $\mathcal{K}$ . This is because  $\mathcal{K}$  is a sub-  
 1137 set of  $\{(i_{j_1}, i_{j_2}) : 1 \leq j_1 < j_2 \leq r\}$ . Conceptually assisted by this understanding, by counting  
 1138 the number of indexes over which the first summation in the definition of  $\widetilde{\Theta}_{\mathcal{K}}^{(i)}$  is running, we  
 1139 have

$$(8.35) \quad \left| \widetilde{\Theta}_{\mathcal{K}}^{(i)} \right| \leq \begin{cases} \rho_n^{s-v} \cdot n^{r-p}, & \text{if } i \in \text{Unique nodes}(\mathcal{K}) \\ \rho_n^{s-v} \cdot n^{r-p-1}, & \text{if } i \notin \text{Unique nodes}(\mathcal{K}) \end{cases}$$

1140 Next, we separate the linear  $\widehat{\Delta}^{(i;v,p)}$  terms, “linear” in the sense they are linear in  $\eta_{(k_1, k_2)}$   
 1141 terms, from those terms quadratic and higher degree in “ $\eta$ ”. The linear term corresponds to  
 1142  $(v, p) = (1, 2)$ , and the higher degree terms correspond to  $v \geq 2$  and  $p \geq 3$ . For the linear  
 1143 part, we have

$$1144 \quad (8.36) \quad \widehat{\Delta}^{(i;1,2)} = \sum_{1 \leq j \leq n : j \neq i} \widetilde{\Theta}_{(i,j)} \eta_{i,j} + \sum_{\substack{1 \leq j_1 < j_2 \leq n \\ j_1, j_2 \neq i}} \widetilde{\Theta}_{(j_1, j_2)} \eta_{j_1, j_2}$$

1145 Conditioned on  $W$ , applying Bernstein’s inequality and (8.35) to the second term on the RHS  
 1146 of (8.36), respectively, we have

$$1147 \quad (8.37) \quad \widehat{\Delta}^{(i;1,2)} = \sum_{1 \leq j \leq n : j \neq i} \widetilde{\Theta}_{(i,j)} \eta_{i,j} + \widetilde{O}_p(\rho_n^{s-1/2} n^{r-2} \cdot \log n)$$

1148 where the first term on the RHS of (8.37) is  $\widetilde{O}_p(\rho_n^{s-1/2} n^{r-3/2} \cdot \log^{1/2} n)$ .

1149 Now we study the higher degree  $\widehat{\Delta}^{(i;v,p)}$  terms. We are going to apply Theorem 8.1. We  
 1150 first upper bound “ $\Xi_{q'}$ ” for all  $q' = 1, \dots, s$  as follows

1151 • If  $R$  is acyclic:

- (i). If  $i \in \mathcal{K}$ : with “ $|S| = q'$ ”, we are summing over  $(p-1) - q' - 1$  node indices in the  
 summation  $\sum_{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{V}(\mathfrak{h}) : \mathcal{S} \subseteq \mathcal{V}(H)}$  – compared to the derivation of (8.11), here we have  
 “ $p-1$ ” instead of “ $p$ ” because the index  $i$  is fixed and cannot vary in the summation.  
 Therefore

$$\sum_{\mathfrak{h} \in \mathcal{H}(H) : \mathcal{V}(\mathfrak{h}) : \mathcal{S} \subseteq \mathcal{V}(H)} |\mathfrak{W}_{\mathfrak{h}}| \leq \rho_n^{s-v} n^{r-p} \cdot n^{p-q'-2} = \rho_n^{s-v} \cdot n^{r-q'-2}$$

1152 and consequently

$$1153 \quad (8.38) \quad \Xi_{q'} \leq \rho_n^{s-v} \cdot n^{r-q'-2} \cdot \rho_n^{v-q'} = \rho_n^{s-q'} \cdot n^{r-q'-2} \leq \rho_n^{s-1} \cdot n^{r-3},$$

1154 under the weak sparsity assumption  $\rho_n = \omega(n^{-1})$ .

1155 (ii). If  $i \notin \mathcal{K}$ : with “ $|S| = q'$ ”, we are summing over  $p - q' - 1$  node indices in the sum-  
 1156 mation  $\sum_{\mathfrak{h} \in \mathcal{H}(H): \mathcal{V}(\mathfrak{h}): \mathcal{S} \subseteq \mathcal{V}(H)}$ . But compared to the “ $i \in \mathcal{K}$ ” case, here we lose an “ $n$ ”  
 1157 factor in  $|\tilde{\Theta}_{\mathcal{K}}^{(i)}|$  according to (8.35). Therefore, we arrive at the identical upper bound  
 1158 for  $\sum_{\mathfrak{h} \in \mathcal{H}(H): \mathcal{V}(\mathfrak{h}): \mathcal{S} \subseteq \mathcal{V}(H)}$  as the above case, namely,  $\Xi_{q'} \leq \rho_n^{s-1} \cdot n^{r-3}$  under the weak  
 1159 sparsity assumption  $\rho_n = \omega(n^{-1})$ .

- The proof for the case where  $R$  is cyclic can be obtained by revising the proof of (8.13). If  $i \in \mathcal{K}$ , we are summing over  $(p - 1) - q'' - 1$  instead of  $p - q'' - 1$  node indices in  $\sum_{\mathfrak{h} \in \mathcal{H}(H): \mathcal{V}(\mathfrak{h}): \mathcal{S} \subseteq \mathcal{V}(H)}$ , if  $i \notin \mathcal{K}$ , then we sum over  $p - q'' - 1$  node indices but will lose an  $n$  factor in the upper bound of  $|\Theta_{\mathcal{K}}^{(i)}|$  according to (8.35). Therefore, both cases would eventually lead to the same upper bound

$$\sum_{\mathfrak{h} \in \mathcal{H}(H): \mathcal{V}(\mathfrak{h}): \mathcal{S} \subseteq \mathcal{V}(H)} |\mathfrak{W}_{\mathfrak{h}}| \leq \rho_n^{s-v} n^{r-p-1} \cdot n^{p-q''-1} = \rho_n^{s-v} \cdot n^{r-q''-2}$$

1160 Similar to the proof of (8.13), it suffices to upper bound those  $\Xi_{q'}$  where  $q' = q''(q'' + 1)/2$ ,  
 1161 and we have

$$1162 \quad (8.39) \quad \Xi_{q''(q''+1)/2} \leq \rho_n^{s-v} \cdot n^{r-q''-2} \cdot \rho_n^{v-q''(q''+1)/2} = \rho_n^{s-q''(q''+1)/2} \cdot n^{r-q''-2}$$

1163 Same as before, the RHS is still monotone in  $q''$  under the assumption  $\rho_n = \omega(n^{-2/r})$  and  
 1164 thus it is bottlenecked by the  $q'' = 1$  case.

1165 Now in order to apply Theorem 8.1 to the higher degree  $\hat{\Delta}^{(i;v,p)}$  terms ( $v \geq 2$  and  $p \geq 3$ ),  
 1166 it only remains to calculate their conditional variances given  $W$ . Notice that given  $W$ , all  
 1167  $\hat{\Delta}^{(i;v,p)}$  terms with different  $(v, p)$  configurations are mutually uncorrelated. We can bound  
 1168 each of them. Straightforward calculations show that

$$1169 \quad \text{Var} \left\{ \hat{\Delta}^{(i;v,p)} | W \right\} \leq \underbrace{\binom{n-1}{p-1} \cdot \rho_n^{2s-2v} \cdot n^{2r-2p} \cdot \rho_n^v}_{\text{sum over } \mathcal{K}\text{-indexed terms: } i \in \mathcal{K}} + \underbrace{\binom{n}{p} \cdot \rho_n^{2s-2v} \cdot n^{2r-2p-2} \cdot \rho_n^v}_{\text{sum over } \mathcal{K}\text{-indexed terms: } i \notin \mathcal{K}}$$

$$1170 \quad = O(\rho_n^{2s-v} \cdot n^{2r-p-1})$$

- For acyclic  $R$ , if  $\rho_n = \omega(n^{-1})$ , we have

$$1172 \quad \text{Var} \left\{ \hat{\Delta}^{(i;v,p)} | W \right\} = O(\rho_n^{2s-v} \cdot n^{2r-p-1}) \leq O(\rho_n^{2s-(p-1)} \cdot n^{2r-p-1})$$

$$1173 \quad (8.40) \quad \leq O(\rho_n^{2s} \cdot n^{2r-2} \cdot (\rho_n \cdot n)^{-2})$$

- For cyclic  $R$ , if  $\rho_n = \omega(n^{-2/r})$ , we have

$$1175 \quad \text{Var} \left\{ \hat{\Delta}^{(i;v,p)} | W \right\} = O(\rho_n^{2s-v} \cdot n^{2r-p-1}) \leq O(\rho_n^{2s} \cdot n^{2r-1} \cdot \rho_n^{-p(p-1)/2} \cdot n^{-p})$$

$$1176 \quad (8.41) \quad \leq O(\rho_n^{2s} \cdot n^{2r-2} \cdot (\rho_n^{-r/2} \cdot n^{-1})^2)$$

1177 Combining (8.38), (8.39), (8.40) and (8.41) with Theorem 8.1, we see that the sum of all  
 1178 higher degree  $\hat{\Delta}^{(i;v,p)}$  terms into  $\hat{\delta}_n$  is at the order of

$$1179 \quad (8.42) \quad \sum_{\substack{\text{All possible } (v,p): \\ v \geq 2, p \geq 3}} \hat{\Delta}^{(i;v,p)} = \tilde{O}_p(\rho_n^s \cdot n^{r-1} \cdot \mathcal{M}(\rho_n, n; R))$$

1180 Compared to the order of the linear  $\hat{\Delta}^{(i;v,p)}$  terms as the leading term on the RHS of (8.37),  
 we see that the higher degree terms are ignorable.

1181 Therefore, for the rest of the proof of Lemma 3.1-(c), we can replace  $\hat{a}_i - a_i$  by

$$\begin{aligned}
 1182 \quad \hat{a}_i - a_i &= \frac{1}{\binom{n-1}{r-1}} \left\{ \sum_{\substack{1 \leq j \leq n \\ j \neq i}} \tilde{\Theta}_{(i,j)} \eta_{ij} + \tilde{O}_p(\rho_n^s \cdot n^{r-1} \cdot \mathcal{M}(\rho_n, n; R)) \right\} \\
 1183 \quad (8.43) \quad &=: \frac{1}{n-1} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} \check{\Theta}_{ij} \eta_{ij} + \tilde{O}_p(\rho_n^s \cdot \mathcal{M}(\rho_n, n; R))
 \end{aligned}$$

where

$$\check{\Theta}_{ij} := \frac{r-1}{\binom{n-2}{r-2}} \tilde{\Theta}_{(i,j)} \leq \rho_n^{s-1}$$

1184 according to (8.35).

1185 Now we are ready to bound the first two terms on the RHS of (8.31) and finish the proof  
 1186 of Lemma 3.1-(c). For term 1, by Bernstein inequality and  $\rho_n = \omega(n^{-1})$ , we have

$$1187 \quad (8.44) \quad \frac{1}{n} \sum_{i=1}^n (\hat{a}_i - a_i)^2 = \tilde{O}_p(\rho_n^{2s} \cdot \mathcal{M}(\rho_n, n; R))$$

1188 For term 2, recalling  $|a_i - U_n| \leq \rho_n^s$ , we have

$$\begin{aligned}
 1189 \quad (8.45) \quad \frac{2}{n} \sum_{i=1}^n (a_i - U_n)(\hat{a}_i - a_i) &= \frac{2}{n(n-1)} \sum_{\substack{1 \leq i \leq n \\ 1 \leq j \leq n \\ i \neq j}} (a_i - U_n) \check{\Theta}_{ij} \eta_{ij} + \tilde{O}_p(\rho_n^{2s} \cdot \mathcal{M}(\rho_n, n; R))
 \end{aligned}$$

1190 conditioned on  $W$ . Applying Bernstein's inequality to the first on the RHS of (8.45) yields a  
 1191 bound of  $\tilde{O}_p(\rho_n^{2s-1/2} \cdot n^{-1} \log n)$ . This completes the proof of Lemma 3.1-(c).

1192 8.2.4. *Proof of Lemma 3.1-(d).* By definition, we have

$$\begin{aligned}
 1193 \quad \frac{n\hat{\sigma}_n^2}{r^2} &= \frac{1}{n} \sum_{i=1}^n \left\{ \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(W_{i, i_1, \dots, i_{r-1}}) - U_n \right\}^2 \\
 1194 \quad &= \frac{1}{n} \sum_{i=1}^n (a_i - U_n)^2 = \frac{1}{n} \sum_{i=1}^n \{ (a_i - \mu_n)^2 + 2(a_i - \mu_n)(\mu_n - U_n) + (\mu_n - U_n)^2 \} \\
 1195 \quad &= \frac{1}{n} \sum_{i=1}^n (a_i - \mu_n)^2 - (U_n - \mu_n)^2
 \end{aligned}$$

1196 Recalling Hoeffding's decomposition for  $U_n$  and applying Theorem 1 of [97] to bound the  
 1197 high-order canonical U-statistics, we have

$$(U_n - \mu_n)^2 = \left\{ \frac{r}{n} \sum_{i=1}^n g_1(X_i) + \tilde{O}_p(\rho_n^s n^{-1} \cdot \log n) \right\}^2 = \tilde{O}_p(\rho_n^{2s} n^{-1} \cdot \log n)$$

1198 We focus on the first term. For notation convenience, define

$$\tilde{a}_i := \mathbb{E} [h(X_i, X_{i_1}, \dots, X_{i_{r-1}}) | X_i] = g_1(X_i) + \mu_n$$

1199 where  $i_1, \dots, i_{r-1} \neq i$  are distinct indexes. We have

$$\begin{aligned}
 1200 \quad \frac{1}{n} \sum_{i=1}^n (a_i - \mu_n)^2 &= \frac{1}{n} \sum_{i=1}^n \{(a_i - \tilde{a}_i) + (\tilde{a}_i - \mu_n)\} \\
 1201 \quad (8.46) \quad &= \frac{1}{n} \sum_{i=1}^n (a_i - \tilde{a}_i)^2 + \frac{2}{n} \sum_{i=1}^n (a_i - \tilde{a}_i)(\tilde{a}_i - \mu_n) + \frac{1}{n} \sum_{i=1}^n (\tilde{a}_i - \mu_n)^2
 \end{aligned}$$

1202 First, we realize that term 3 on the RHS of (8.46) is simply

$$(8.47) \quad \frac{1}{n} \sum_{i=1}^n (\tilde{a}_i - \mu_n)^2 = \frac{1}{n} \sum_{i=1}^n g_1^2(X_i)$$

1203 Now we focus on handling terms 1 and 2. The key part is to handle  $a_i - \tilde{a}_i$ . By applying the  
 1204 Hoeffding's ANOVA decomposition of an arbitrary symmetric statistic (1.1)–(1.3) in [14]  
 1205 onto each single  $h(X_i, X_{i_1}, \dots, X_{i_{r-1}})$  term, we can see that

$$\begin{aligned}
 1206 \quad a_i - \tilde{a}_i &= \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} \{h(X_i, X_{i_1}, \dots, X_{i_{r-1}}) - \mathbb{E}[h(X_i, X_{i_1}, \dots, X_{i_{r-1}}) | X_i]\} \\
 1207 \quad &= \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} \left\{ \sum_{k=1}^{r-1} \sum_{\substack{1 \leq j_1 < \dots < j_k \leq n \\ \{j_1, \dots, j_k\} \subseteq \{i_1, \dots, i_{r-1}\}}} g_{k+1}(X_i, X_{j_1}, \dots, X_{j_k}) \right\} \\
 1208 \quad (8.48) \quad &= \frac{1}{\binom{n-1}{r-1}} \sum_{k=1}^{r-1} \binom{n-k-1}{r-k-1} \sum_{\substack{1 \leq j_1 < \dots < j_k \leq n \\ j_1, \dots, j_k \neq i}} g_{k+1}(X_i, X_{j_1}, \dots, X_{j_k})
 \end{aligned}$$

1209 Now we apply Theorem 1 of [97] to the RHS of (8.48), we see that

$$\begin{aligned}
 1210 \quad (8.49) \quad a_i - \tilde{a}_i &= \frac{r-1}{n-1} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} g_2(X_i, X_j) + \tilde{O}_p(\rho_n^s n^{-1} \cdot \log n) \\
 1211 \quad (8.50) \quad &= \tilde{O}_p(\rho_n^s n^{-1/2} \cdot \log^{1/2} n)
 \end{aligned}$$

1212 Now we are ready to continue bounding the RHS of (8.46). Using (8.50), term 1 on the RHS  
 1213 of (8.46) is

$$(8.51) \quad (a_i - \tilde{a}_i)^2 = \tilde{O}_p(\rho_n^{2s} n^{-1} \cdot \log n)$$

1214 Using (8.49), term 2 on the RHS of (8.46) is

$$\begin{aligned}
 1215 \quad &\frac{2}{n} \sum_{i=1}^n (a_i - \tilde{a}_i)(\tilde{a}_i - \mu_n) \\
 1216 \quad &= \frac{2}{n} \sum_{i=1}^n \left\{ \frac{r-1}{n-1} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} g_2(X_i, X_j) + \tilde{O}_p(\rho_n^s n^{-1} \cdot \log n) \right\} g_1(X_i) \\
 1217 \quad (8.52) \quad &= \frac{2(r-1)}{n(n-1)} \sum_{\substack{1 \leq i \leq n \\ 1 \leq j \leq n \\ i \neq j}} g_1(X_i) g_2(X_i, X_j) + \tilde{O}_p(\rho_n^{2s} n^{-1} \cdot \log n)
 \end{aligned}$$

1218 Finally, combining (8.51), (8.52) and (8.47) completes the proof of Lemma 3.1-(d).

1219 8.3. *Proof of Theorem 3.1.* We mainly prove for the case  $\rho_n = O((\log n)^{-1})$  without  
 1220 non-lattice condition. We will explain how this proof can be revised for the other case with  
 1221 Carmer’s condition but without a  $\rho_n$  upper bound.

1222 LEMMA 8.1 (Esseen’s smoothing lemma ([52], Section XVI.3)). *For any distribution*  
 1223 *function  $F$  and a general function  $G$  that has universally bounded derivative and satisfy*  
 1224  $G(-\infty) = 0, G(\infty) = 1$ , we have

$$(8.53) \quad \|F(u) - G(u)\|_\infty \leq C_1 \int_{-\gamma}^\gamma \left| \frac{\text{Ch.f.}(F; t) - \text{Ch.f.}(G; t)}{t} \right| dt + \frac{C_2 \sup_u |G'(u)|}{\gamma}$$

for universal constants  $C_1, C_2 > 0$ , where  $\text{Ch.f.}(G; t)$  is defined to be the characteristic function of  $G$  as follows

$$\text{Ch.f.}(G; t) := \int_{-\infty}^\infty e^{itx} dG(x)$$

Recall the definition of  $\tilde{T}_n$  from (3.9) that

$$\tilde{T}_n = U_n^\# + \Delta_n - \frac{U_n^\#}{2} \delta_n \quad \text{and} \quad \hat{T}_n = \tilde{T}_n + \hat{\Delta}_n + \tilde{O}_p(\mathcal{M}(\rho_n, n; R)).$$

1225 We define a random variable  $\tilde{\Delta}_n | W \sim N(0, (\rho_n \cdot n)^{-1} \sigma_\omega^2)$ , that  $\tilde{\Delta}_n$  is conditionally in-  
 1226 dependent of  $A$ , given  $W$ . By Lemma 3.1-(b), we have  $\sup_{u \in \mathbb{R}} |F_{\tilde{\Delta}_n}(u) - F_{\tilde{\Delta}_n}^\sim(u)| =$   
 1227  $O(\rho_n^{-1/2} \cdot n^{-1})$ . We are going to show that

$$1228 \quad (8.54) \quad \|F_{\hat{T}_n}(u) - F_{\tilde{T}_n + \tilde{\Delta}_n}(u)\|_\infty = O(\mathcal{M}(\rho_n, n; R))$$

$$1229 \quad (8.55) \quad \|F_{\tilde{T}_n + \tilde{\Delta}_n}(u) - F_{\tilde{T}_n + \tilde{\Delta}_n}^\sim(u)\|_\infty = O(\rho_n^{-1/2} \cdot n^{-1})$$

$$1230 \quad (8.56) \quad \|F_{\tilde{T}_n + \tilde{\Delta}_n}^\sim(u) - G_n(u)\|_\infty = O((\rho_n \cdot n)^{-1} + n^{-1} \log n)$$

1231 where  $G_n(u)$  is defined in (3.13). To proceed, we need the following smoothing lemma.

LEMMA 8.2. *Suppose we have random variables  $X, Y, Z$  satisfying*

$$X = Y + Z$$

1232 *such that the CDF of  $Y$  is smooth, and there exists a universal constant  $0 < M < \infty$  such*  
 1233 *that  $F_Y(u + a) - F_Y(u) \leq M \cdot a + O(\zeta_n)$  for any  $u \in \mathbb{R}$  and  $a > 0$ . Also assume that  $\mathbb{P}(|Z| \geq$   
 1234  $\tilde{\zeta}_n) \leq n^{-1}$ , that is,  $Z = \tilde{O}_p(\tilde{\zeta}_n)$ . We have*

$$\|F_X(u) - F_Y(u)\|_\infty = O(\zeta_n + \tilde{\zeta}_n + n^{-1})$$

1235 **Remark.** We emphasize that Lemma 8.2 does *not* require any independence between  $X$ ,  
 1236  $Y$  and  $Z$ .

1237 PROOF OF LEMMA 8.2. Since “ $Y + Z > u$ ” implies the union of the following two  
 1238 events: “ $Y > u - a$ ” and “ $Z > a$ ”, we have  $1 - \mathbb{P}(Y + Z \leq u) \leq 1 - \mathbb{P}(Y \leq u - a) + \mathbb{P}(|Z| \geq$   
 1239  $a)$ , which further implies that

$$1240 \quad \mathbb{P}(Y + Z \leq u) \geq \mathbb{P}(Y \leq u - a) - \mathbb{P}(|Z| > a)$$

$$1241 \quad \geq \mathbb{P}(Y \leq u) - M \cdot a - O(\zeta_n) - \mathbb{P}(|Z| > a)$$

$$1242 \quad (\text{Setting } a = \tilde{\zeta}_n) \geq \mathbb{P}(Y \leq u) - O(\zeta_n + \tilde{\zeta}_n + n^{-1})$$

1243 On the other hand, we have

$$\begin{aligned}
1244 \quad & \mathbb{P}(Y + Z \leq u) = \int_z \mathbb{P}(Y \leq u - z | Z = z) dP_Z(z) \\
1245 \quad & = \int_{z: |z| \leq a} \mathbb{P}(Y \leq u - z | Z = z) dP_Z(z) + \int_{z: |z| > a} \mathbb{P}(Y \leq u - z | Z = z) dP_Z(z) \\
1246 \quad & \leq \int_z \mathbb{P}(Y \leq u + a | Z = z) dP_Z(z) + \int_{z: |z| > a} 1 dP_Z(z) \\
1247 \quad & \leq \mathbb{P}(Y \leq u + a) + \mathbb{P}(|Z| \geq a)
\end{aligned}$$

1248 Setting  $a = \tilde{\zeta}_n$ , the RHS is upper bounded by  $\mathbb{P}(Y \leq u) + O(\zeta_n + \tilde{\zeta}_n + n^{-1})$ . Combining the  
1249 two inequalities proves Lemma 8.2.  $\square$

1250 Now we return to the main proof of Theorem 3.1. Our proof would proceed as follows.  
1251 We shall use Lemma 3.1-(b) to prove (8.55); then with the assistance of Lemma 8.2, we use  
1252 (8.56) and (8.55) to prove (8.54); finally, we state the proof of (8.56) without needing (8.54)  
1253 or (8.55).

1254 • Proof of “Lemma 3.1-(b)  $\Rightarrow$  (8.55)”. Noticing that  $\tilde{T}_n$  does not depend on the random  
1255 variations of  $A|W$  given  $W$ , but it is determined if  $W$  is given, we have

$$\begin{aligned}
1256 \quad & F_{\tilde{T}_n + \check{\Delta}_n}(u) = \mathbb{P}(\tilde{T}_n + \check{\Delta}_n \leq u) \\
1257 \quad & = \mathbb{E} \left[ \mathbb{P}(\tilde{T}_n + \check{\Delta}_n \leq u | W) \right] \\
1258 \quad & = \mathbb{E} \left[ \mathbb{P}(\check{\Delta}_n \leq u - \tilde{T}_n | W) \right] \\
1259 \quad & \stackrel{\text{Lemma 3.1-(b)}}{=} \mathbb{E} \left[ \mathbb{P}(\check{\Delta}_n \leq u - \tilde{T}_n | W) + \tilde{O}_p(\rho_n^{-1/2} \cdot n^{-1}) \right] \\
1260 \quad & = \mathbb{E} \left[ \mathbb{P}(\tilde{T}_n + \check{\Delta}_n \leq u | W) \right] + O(\rho_n^{-1/2} \cdot n^{-1}) \\
1261 \quad & = F_{\tilde{T}_n + \check{\Delta}_n}(u) + O(\rho_n^{-1/2} \cdot n^{-1})
\end{aligned}$$

1262 • Proof of “(8.55), (8.56) and Lemma 8.2  $\Rightarrow$  (8.54)”. We set  $Y = \tilde{T}_n + \check{\Delta}_n$  and  $Z = \hat{T}_n - Y$ .  
1263 We notice that by Lemma 3.1-(b), we have  $Z = \tilde{O}_p(\mathcal{M}(\rho_n, n; R))$  meaning that  $\mathbb{P}(|Z| \geq$   
1264  $C_1 \mathcal{M}(\rho_n, n; R)) = O(n^{-1})$ . Next we verify that  $Y$  satisfies the condition of Lemma 8.2,  
1265 we notice that (8.56) implies that for any  $u \in \mathbb{R}$  and  $a > 0$ , we have

$$\begin{aligned}
1266 \quad & F_{\tilde{T}_n + \check{\Delta}_n}(u + a) - F_{\tilde{T}_n + \check{\Delta}_n}(u) \\
1267 \quad & \leq \underbrace{\left| F_{\tilde{T}_n + \check{\Delta}_n}(u + a) - F_{\tilde{T}_n + \check{\Delta}_n}(u + a) \right|}_{\text{Bounded by (8.55)}} + \underbrace{\left| F_{\tilde{T}_n + \check{\Delta}_n}(u) - F_{\tilde{T}_n + \check{\Delta}_n}(u) \right|}_{\text{Bounded by (8.55)}} \\
1268 \quad & + \underbrace{\left| F_{\tilde{T}_n + \check{\Delta}_n}(u + a) - G_n(u + a) \right|}_{\text{Bounded by (8.56)}} + \underbrace{\left| G_n(u + a) - G_n(u) \right|}_{\sup_{u,n} |G'_n(u)| < \infty} \\
1269 \quad & + \underbrace{\left| F_{\tilde{T}_n + \check{\Delta}_n}(u) - G_n(u) \right|}_{\text{Bounded by (8.56)}} \\
1270 \quad & \leq C \cdot a + O(\rho_n^{-1/2} \cdot n^{-1})
\end{aligned}$$

1271 Then applying Lemma 8.2 and noticing that  $\mathcal{M}(\rho_n, n; R)$  dominates all of  $\rho_n^{-1/2} \cdot n^{-1}$ ,  
 1272  $(\rho_n \cdot n)^{-1}$  and  $n^{-1} \log n$  completes the proof of (8.54).

1273 Next, we focus on proving (8.56). In this proof, we shall set  $\gamma = n$  in Esseen's smoothing  
 1274 lemma and break the integration range into three parts:  $|t| \in (0, n^\epsilon)$ ,  $(n^\epsilon, n^{1/2})$  and  $(n^{1/2}, n)$

1275 LEMMA 8.3. *We have the following bounds:*

(a). *For any fixed  $\epsilon > 0$ , we have*

$$\int_{n^\epsilon}^n \left| \frac{\text{Ch.f.}^1(G_n; t)}{t} \right| dt = O(n^{-1})$$

1276

(b). *For a small enough constant  $c_\rho > 0$ , if  $\rho_n \leq c_\rho (\log n)^{-1}$ , we have*

$$\int_{C_1 n^{1/2}}^n \left| \frac{\mathbb{E} \left[ e^{it(\tilde{T}_n + \tilde{\Delta}_n)} \right]}{t} \right| dt = O(n^{-1})$$

1277 *for an arbitrary constant  $C_1 > 0$ .*

(c). *For a small enough constant  $C_1 > 0$  and arbitrary fixed  $\epsilon > 0$ , we have*

$$\int_{n^\epsilon}^{C_1 n^{1/2}} \left| \frac{\mathbb{E} \left[ e^{it(\tilde{T}_n + \tilde{\Delta}_n)} \right]}{t} \right| dt = O(n^{-1} \log n).$$

(d). *For fixed  $\epsilon > 0$  chosen such that  $\epsilon \leq 1/7$ , then we have*

$$\int_0^{n^\epsilon} \left| \frac{\mathbb{E} \left[ e^{it(\tilde{T}_n + \tilde{\Delta}_n)} \right] - \text{Ch.f.}(G_n; t)}{t} \right| dt = O((\rho_n \cdot n)^{-1} + n^{-1} \log n).$$

1278 PROOF OF LEMMA 8.3. First of all, we notice that between two parts  $\tilde{T}_n$  and  $\tilde{\Delta}_n$ , the  
 1279 former is completely determined by  $W$ , and the latter follows  $N(0, (\rho_n \cdot n)^{-1} \cdot \sigma_w^2)$ , where  
 1280  $\sigma_w^2 \asymp 1$  is a U-statistic of  $X_1, \dots, X_n$ . We have

$$\begin{aligned} 1281 \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{it\tilde{\Delta}_n} \right] &= \mathbb{E} \left[ \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{it\tilde{\Delta}_n} \mid W \right] \right] = \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot \mathbb{E} \left[ e^{it\tilde{\Delta}_n} \mid W \right] \right] \\ 1282 &= \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right] \end{aligned}$$

1283 Then we prove each of the bounds in the lemma.

1284 (a). Notice that for each of  $k = -1, 0, 1, 2, 3, \dots$ , we always have  $t^k e^{-t^2/2} \leq C_k e^{-t^2/3}$  when  
 1285  $t > 1$  for universal constants  $C_k > 0$  that only depend on  $k$ . From the classical literature  
 1286 on Hermite polynomials, we recall that function  $\text{Ch.f.}(G_n; t)$  takes the form of  $e^{-t^2/2}$   
 1287 multiplies a polynomial of  $t$ . Therefore, for  $k = -1, 0, 1, 2, 3 \dots$

$$1288 \int_{n^\epsilon}^n |\text{Ch.f.}(G_n; t)/t| dt \leq (C_{-1} + \dots + C_{d_g-1}) \int_{n^\epsilon}^\infty e^{-t^2/3} dt = O(n^{-1})$$

1289 where  $d_g := \text{degree of Ch.f.}(G_n; t)$  is a fixed finite number.

---

<sup>1</sup>Ch.f.: characteristic function. For the Edgeworth expansion function  $G_n$  that is not necessarily a valid CDF, its Ch.f. is defined to be its Fourier transform.

1290 (b). For  $|t| \geq n^{1/2}$ , we have

$$\begin{aligned}
1291 \quad & \left| \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right] \right| \leq \mathbb{E} \left[ \left| e^{it\tilde{T}_n} \right| \cdot \left| e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right| \right] \\
1292 \quad & = \mathbb{E} \left[ e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right] \leq \mathbb{E} \left[ e^{-(\rho_n \cdot n)^{-1} \mathbb{E}[\sigma_w^2] / 4 \cdot t^2} \right] + \mathbb{P} \left( \sigma_w^2 < \mathbb{E}[\sigma_w^2] / 4 \right) \\
1293 \quad (8.57) \quad & \leq e^{-C_1 \cdot \rho_n^{-1}} + e^{-C_2 n} = C n^{-C_1 \cdot c_\rho^{-1}}
\end{aligned}$$

since  $\rho_n^{-1} = c_\rho^{-1} \log n$ , and notice that  $\mathbb{P}(\sigma_w^2 < \mathbb{E}[\sigma_w^2]/4)$  diminishes exponentially fast because  $\sigma_w^2$  is a U-statistic (as will be proved in the proof of part (c) below) dominated by its linear part and concentration inequalities such as Bernstein's. Then choosing  $c_\rho = (4C_1)^{-1}$  finishes the proof of Lemma 8.3-(b) since

$$\int_{C_1 n^{1/2}}^n t^{-1} dt = O(\log n)$$

1294 (c). For this part of the proof, we show that  $\sigma_w^2$  can be written as the sum of U-statistics thus  
1295 Hoeffding's decomposition to U-statistics conveniently applies to it<sup>2</sup>. Then we combine  
1296 this argument with the argument used in [21]. Recall that  $\hat{\Theta}_{ij} \asymp \rho_n^{-1} \cdot n^{1/2}$ , and it is a U-  
1297 statistic with the index set  $\{1, \dots, n\} \setminus \{i, j\}$ , thus the Hoeffding's decomposition implies:

$$(8.58) \quad \hat{\Theta}_{ij} \cdot \rho_n \cdot n^{-1/2} = \theta_{ij} + \frac{C}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i, j}} \check{g}_1(X_k; X_i, X_j) + \tilde{O}_p(n^{-1} \cdot \log n)$$

1298 where  $\theta_{ij} := \mathbb{E}[\hat{\Theta}_{ij} | X_i, X_j] \cdot \rho_n \cdot n^{-1/2}$ , and we used [97] to obtain a probabilistic upper  
1299 bound of the higher order terms in Hoeffding's decomposition. Then we have

$$\begin{aligned}
1300 \quad & \sigma_w^2 = \rho_n \cdot n \cdot \text{Var} \left( \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{\Theta}_{ij} \eta_{ij} | W \right) = \frac{\rho_n \cdot n}{\binom{n}{2}^2} \sum_{1 \leq i < j \leq n} \hat{\Theta}_{ij}^2 W_{ij} (1 - W_{ij}) \\
1301 \quad & = \frac{\rho_n \cdot n}{\binom{n}{2}^2} \cdot \rho_n^{-2} \cdot n \cdot \sum_{1 \leq i < j \leq n} \left\{ \theta_{ij} + \frac{C}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i, j}} \check{g}_1(X_k; X_i, X_j) + \tilde{O}_p(n^{-1} \cdot \log n) \right\}^2 \cdot W_{ij} (1 - W_{ij}) \\
1302 \quad & = \frac{\rho_n^{-1} n^2}{\binom{n}{2}^2} \sum_{1 \leq i < j \leq n} \left\{ \theta_{ij}^2 + \frac{2C\theta_{ij}}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i, j}} \check{g}_1(X_k; X_i, X_j) + \tilde{O}_p(n^{-1} \cdot \log n) \right\} W_{ij} (1 - W_{ij}) \\
1303 \quad & = \frac{\rho_n^{-1} \cdot n^2 \sum_{1 \leq i < j \leq n} \theta_{ij}^2 W_{ij} (1 - W_{ij})}{\binom{n}{2}^2} \\
(8.59) \quad & + \frac{\rho_n^{-1} \cdot n^2 \cdot 2C}{(n-2) \cdot \binom{n}{2}^2} \sum_{\substack{1 \leq i < j \leq n \\ 1 \leq k \leq n \\ k \neq i, j}} \check{g}_1(X_k; X_i, X_j) W_{ij} (1 - W_{ij}) + \tilde{O}_p(n^{-1} \cdot \log n)
\end{aligned}$$

<sup>2</sup>Notice that in this part of the proof, we cannot simply bound the  $\sigma_w$  term away because it is dependent on any individual term in the expansion of  $\tilde{T}_n$ .

1305 where we used the fact  $|(n-2)^{-1} \sum_{1 \leq k \leq n; k \neq i, j} \check{g}_1(X_k; X_i, X_j)| = \tilde{O}_p((n^{-1} \log n)^{1/2})$   
 1306 by Bernstein inequality.

Clearly, the first term in (8.59) is a U-statistic of degree 2, where the individual term is at the order

$$\frac{\rho_n^{-1} \cdot n^2 \cdot \theta_{ij}^2 \cdot W_{ij}(1 - W_{ij})}{\binom{n}{2}} \asymp \frac{\rho_n^{-1} \cdot n^2 \cdot 1 \cdot \rho_n}{n^2} \asymp 1$$

1307 Now we focus on the second term and re-express it as a U-statistic. We have

$$\begin{aligned} 1308 \quad & \sum_{\substack{1 \leq i < j \leq n \\ 1 \leq k \leq n \\ k \neq i, j}} \theta_{ij} \check{g}_1(X_k; X_i, X_j) W_{ij}(1 - W_{ij}) = \frac{1}{2} \sum_{\substack{1 \leq \{i, j, k\} \leq n \\ i \neq j, j \neq k, k \neq i}} \theta_{ij} \check{g}_1(X_k; X_i, X_j) W_{ij}(1 - W_{ij}) \\ 1309 \quad & = \frac{1}{2} \sum_{\substack{1 \leq \{i, j, k\} \leq n \\ i \neq j, j \neq k, k \neq i}} \left[ \frac{1}{3} \left\{ \theta_{ij} \check{g}_1(X_k; X_i, X_j) W_{ij}(1 - W_{ij}) \right. \right. \\ 1310 \quad & \left. \left. + \theta_{ki} \check{g}_1(X_j; X_k, X_i) W_{ki}(1 - W_{ki}) + \theta_{jk} \check{g}_1(X_i; X_j, X_k) W_{jk}(1 - W_{jk}) \right\} \right] \end{aligned}$$

(8.60)

$$1311 \quad =: \sum_{1 \leq i < j < k \leq n} \check{H}(X_i, X_j, X_k)$$

1312 where we denote

$$\begin{aligned} 1313 \quad & \check{H}(X_i, X_j, X_k) := \theta_{ij} \check{g}_1(X_k; X_i, X_j) W_{ij}(1 - W_{ij}) \\ 1314 \quad & + \theta_{ki} \check{g}_1(X_j; X_k, X_i) W_{ki}(1 - W_{ki}) + \theta_{jk} \check{g}_1(X_i; X_j, X_k) W_{jk}(1 - W_{jk}) \end{aligned}$$

Clearly,  $\check{H}(X_i, X_j, X_k)$  is symmetric in  $X_i, X_j, X_k$ , and the individual term

$$\frac{\rho_n^{-1} \cdot n^2 \cdot 2C \cdot \binom{n}{3}}{(n-2) \cdot \binom{n}{2}^2} \cdot \check{H}(X_i, X_j, X_k) \asymp \frac{\rho_n^{-1} \cdot n^2 \cdot n^3}{n^5} \cdot \rho_n \asymp 1$$

1315 So the second term on the RHS of (8.59) is a U-statistic of degree 3. Therefore,  $\sigma_w^2$  can be  
 1316 re-expressed as Hoeffding's decomposition for U-statistics as follows

$$(8.61) \quad \sigma_w^2 = \mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i) + \tilde{O}_p(n^{-1} \cdot \log n)$$

1317 where we again applied [97] to derive the probabilistic upper bound for the higher order  
 1318 terms in Hoeffding's decomposition.

1319 Now, we are ready to upper bound the characteristic function for  $n^\epsilon \leq |t| \leq n^{1/2}$

$$\begin{aligned} 1320 \quad & \left| \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right] \right| \\ 1321 \quad & \leq \left| \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} t^2 / 2 \cdot \left\{ \mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i) + \tilde{O}_p(n^{-1} \cdot \log n) \right\}} \right] \right| \end{aligned}$$

(8.62)

$$1322 \quad = \left| \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} t^2 / 2 \cdot \left\{ \mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i) \right\}} \cdot \left( 1 + \tilde{O}_p(\rho_n^{-1} \cdot n^{-2} \log n \cdot t^2) \right) \right] \right|$$

1323 where in the last line, we used the fact that  $|e^z - 1| = O(|z|)$  for all universally bounded  
 1324  $z \in \mathbb{C}$  (here  $|t| = O(n^{1/2})$ ) and by assumption  $\rho_n \cdot \log n = O(1)$ ). Then since

$$1325 \quad (8.63) \quad \int_{n^\epsilon}^{n^{1/2}} \frac{\rho_n^{-1} \cdot n^{-2} \cdot \log n \cdot t^2}{t} dt \asymp (\rho_n \cdot n)^{-1} \cdot \log n$$

1326 we know that this  $\tilde{O}_p(\rho_n^{-1} \cdot n^{-2} \cdot \log n \cdot t^2)$  term can be ignored in (8.62). Continuing  
 1327 (8.62), we have

$$1328 \quad \text{RHS of (8.62)} \leq e^{-(\rho_n \cdot n)^{-1} t^2 / 2 \cdot \mathbb{E}[\sigma_w^2]} \cdot \left| \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-\rho_n^{-1} \cdot n^{-2} \cdot \sum_{i=1}^n g_{\sigma,1}(X_i) \cdot t^2} \right] \right|$$

1329 We are going to show that  $\tilde{T}_n$  can be expressed as a U-statistic of degree 2 plus an  
 1330  $\tilde{O}_p(n^{-1} \log^{3/2} n)$  remainder term, which can be ignored. Indeed,

$$1331 \quad \begin{aligned} \tilde{T}_n &= U_n^\# + \Delta_n - \frac{1}{2} \cdot U_n^\# \cdot \delta_n \\ &= \frac{1}{\sqrt{n}\xi_1} \sum_{i=1}^n g_1(X_i) + \frac{r-1}{\sqrt{n}(n-1)\xi_1} \sum_{1 \leq i < j \leq n} g_2(X_i, X_j) \\ &\quad + \frac{1}{n^{3/2}\xi_1} \sum_{i=1}^n g_1(X_i) \sum_{j=1}^n \frac{g_1^2(X_j) - \xi_1^2}{\xi_1^2} + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n). \end{aligned}$$

1334 Since  $n^{-3/2} \sum_{i=1}^n g_1(X_i) (g_1(X_i)^2 - \xi_1^2) / \xi_1^3 = \tilde{O}_p(n^{-1} \cdot \log^{1/2} n)$ , we can write

$$1335 \quad \begin{aligned} \tilde{T}_n &= \frac{1}{\sqrt{n}\xi_1} \sum_{i=1}^n g_1(X_i) + \frac{r-1}{\sqrt{n}(n-1)\xi_1} \sum_{1 \leq i < j \leq n} g_2(X_i, X_j) \\ 1336 \quad &+ \frac{1}{n^{3/2}\xi_1} \sum_{1 \leq i < j \leq n} \frac{g_1(X_i) (g_1^2(X_j) - \xi_1^2) + g_1(X_j) (g_1^2(X_i) - \xi_1^2)}{\xi_1^2} + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n) \\ 1337 \quad &=: \frac{1}{\sqrt{n}\xi_1} \sum_{i=1}^n g_1(X_i) + \frac{r-1}{\sqrt{n}(n-1)} \sum_{1 \leq i < j \leq n} \tilde{g}_2(X_i, X_j) + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n) \end{aligned}$$

which therefore is expressed as a U-statistic of degree 2 plus an  $\tilde{O}_p(n^{-1} \cdot \log^{3/2} n)$  term,  
 where  $\mathbb{E}[\tilde{g}_2(X_i, X_j)] = 0$  and  $\mathbb{E}[\tilde{g}_2^2(X_i, X_j)] = O(1)$ . To prove the claimed bound, we  
 can choose a positive integer  $m$  (depending on  $t$ ) and write

$$\sum_{1 \leq i < j \leq n} \tilde{g}_2(X_i, X_j) = \sum_{i=1}^m \sum_{j=i+1}^n \tilde{g}_2(X_i, X_j) + \sum_{i=m+1}^{n-1} \sum_{j=i+1}^n \tilde{g}_2(X_i, X_j)$$

1338 Then the arguments of [21, eq. (2.17)-(2.20)] can be applied here. Notice that this part  
 1339 of the proof of [21] does not require non-lattice assumption, but all it requires on the  
 1340 behavior of  $|\mathbb{E}[e^{itg_1(X_i)/(\sqrt{n}\xi_1)}]|$  is its closeness to 1 for  $t/\sqrt{n} \approx 0$ . Indeed, for  $n\rho_n \gg 1$   
 1341 and  $t \leq c_1 n^{1/2}$  with small  $c_1 > 0$ ,

$$1342 \quad \begin{aligned} &|\mathbb{E} e^{itg_1(X_i)/(\sqrt{n}\xi_1) - \rho_n^{-1} n^{-2} t^2 / 2g_{\sigma,1}(X_i)}| \\ 1343 \quad &\leq \left| \mathbb{E} \left( 1 + \frac{1}{2} \left( \frac{itg_1(X_i)}{\sqrt{n}\xi_1} - \frac{t^2 g_{\sigma,1}(X_i)}{2\rho_n n^2} \right)^2 \right) \right| + O \left( \mathbb{E} \left| \frac{itg_1(X_i)}{\sqrt{n}\xi_1} - \frac{t^2 g_{\sigma,1}(X_i)}{2\rho_n n^2} \right|^3 \right) \\ 1344 \quad &\leq 1 - \frac{t^2}{3n} \leq \exp \left\{ -\frac{t^2}{3n} \right\}. \end{aligned}$$

1345 The proof of Lemma 8.3-(c) is therefore completed after applying the arguments of [21,  
 1346 eq. (2.17)-(2.20)].

1347 (d). Finally, in this part, we calculate the expansion of  $\mathbb{E}\left[e^{it\tilde{T}_n}\right]$  and derive the Edgeworth  
 1348 expansion for  $|t| \leq n^\epsilon$  for a small enough fixed  $\epsilon$ . The main portion of the proof for  
 1349 this part, i.e., our calculations in (8.69), (8.70), (8.73) and (8.74) that we are going to  
 1350 present, follow the roadmap in classical literature on Edgeworth expansion for noiseless  
 1351 U-statistics, laid out by [21, 71, 90, 96]. Our  $\tilde{T}_n$  is different from their studentization/  
 1352 standardization forms by using a different rescaler, so this part is not a direct corollary  
 1353 of their results. Despite the resulting differences is non-essential, we nonetheless present  
 1354 the full calculation steps for completeness and for the convenience of the readers.

1355 To start, we have

$$\begin{aligned}
 1356 & \mathbb{E}\left[e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2}\right] \\
 1357 & (8.64) \quad = \mathbb{E}\left[e^{it\tilde{T}_n} \cdot \left\{1 - \frac{\sigma_w^2 t^2}{2\rho_n \cdot n} + \frac{\sigma_w^4 t^4}{8\rho_n^2 \cdot n^2} + O\left(\frac{\sigma_w^6 t^6}{\rho_n^3 \cdot n^3}\right)\right\}\right]
 \end{aligned}$$

1358 as long as  $n\rho_n = \omega(n^{2\epsilon})$ . We first bound the remainder, we have  $\int_0^{n^\epsilon} (\sigma_w^6 t^6) (\rho_n^3 n^3) \cdot$   
 1359  $t^{-1} dt \asymp n^{6\epsilon} \cdot (\rho_n \cdot n)^{-3}$ . Since the assumption of Theorem 3.1 implies that  $\rho_n = \omega(n^{-1/2})$   
 1360 in any case, so setting  $\epsilon \leq 1/13$  yields  $n^{6\epsilon} \cdot (\rho_n \cdot n)^{-3} = O(n^{-1})$ . We have

$$\begin{aligned}
 1361 & e^{it\tilde{T}_n} = e^{it(U_n^\# + \Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n)} \\
 1362 & = e^{itU_n^\#} \left\{1 + \left(\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n\right) it - \frac{1}{2} \cdot \left(\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n\right)^2 t^2\right\} \\
 1363 & (8.65) \quad + \tilde{O}_p\left(\left|\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n\right|^3 t^3\right)
 \end{aligned}$$

1364 To bound the remainder term, notice that  $|1 - \sigma_w^2 t^2 / (\rho_n \cdot n)| \leq 1$  for  $|t| \leq n^\epsilon$ , where we  
 1365 recall that Theorem 3.1 we are proving here always assumes  $\rho_n = \omega(n^{-1/2})$  in all cases.  
 1366 Then, setting  $\epsilon \leq 1/7$  together with the fact  $U_n^\# = \tilde{O}_p(\log^{1/2} n)$ ,  $\Delta_n = \tilde{O}_p(n^{-1/2} \log n)$ ,  
 1367  $\delta_n = \tilde{O}_p(n^{-1/2} \log^{1/2} n)$ , by Bernstein's inequality and [97], we have

$$\begin{aligned}
 1368 & \int_0^{n^\epsilon} \left|\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n\right|^3 t^3 \cdot \frac{1}{t} dt = \tilde{O}_p\left(n^{-3/2} \cdot n^{3\epsilon} \log^{3/2} n\right) \\
 1369 & (8.66) \quad = \tilde{O}_p(n^{-15/14} \log^{3/2} n) = \tilde{O}_p(n^{-1})
 \end{aligned}$$

1370 and this remainder term can also be ignored. Now we deal with the main part of the terms.

1371 Set  $\varphi_n(t) := \mathbb{E}\left[e^{it \cdot \frac{g_1(X_1)}{\sqrt{n \cdot \xi_1}}}\right]$ . Then by Section VI, Lemma 4 of [108], we have

$$1372 \quad (8.67) \quad \varphi_n^n(t) = e^{-t^2/2} \left(1 - n^{-1/2} \cdot \frac{i\mathbb{E}[g_1^3(X_1)]t^3}{6\xi_1^3}\right) + O\left(n^{-1} \cdot \log n \cdot P_0(t)e^{-t^2/4}\right)$$

$$1373 \quad \varphi_n^{n-k}(t) = \varphi_n^n(t) + O\left(n^{-1} \cdot \log n \cdot P_k(t)e^{-t^2/4}\right)$$

1374 for any fixed  $k = 0, 1, 2, 3$ , where  $P_0(t), \dots, P_k(t)$  are fixed polynomials of  $t$  and each  
 1375 of them can be divided by  $t$ . Here, we first focus on  $\mathbb{E}[e^{it\tilde{T}_n}]$ , and then handle  $\mathbb{E}[e^{it\tilde{T}_n} \cdot$   
 1376  $\sigma_w^2 t^2 / (\rho_n \cdot n)]$ . For  $\mathbb{E}[e^{it\tilde{T}_n}]$ , by ignoring the small term in (8.65), we have

$$1377 \quad (8.68) \quad \mathbb{E}\left[e^{it\tilde{T}_n}\right] = \mathbb{E}\left[e^{itU_n^\#} \left\{1 + it \left(\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n\right) - \frac{t^2}{2} \left(\Delta_n - \frac{1}{2}U_n^\# \cdot \delta_n\right)^2\right\}\right]$$

1378 Now we inspect each term on the RHS of (8.68). For  $\mathbb{E}[e^{itU_n^\#}]$  we use (8.67) and obtain  
 1379  $\mathbb{E}[itU_n^\#] = \varphi_n^n(t)$ . For the next term, recall that  $\mathbb{E}[g_2(X_1, X_2)] = 0$  and  $\mathbb{E}[g_1^k(X_1)g_2(X_1, X_2)] =$   
 1380  $0$  for all  $k \in \mathbb{N}$ . We have

$$\begin{aligned}
 1381 \quad \mathbb{E}\left[e^{itU_n^\#} \cdot it\Delta_n\right] &= \mathbb{E}\left[e^{itU_n^\#} \cdot it \cdot \frac{r-1}{\sqrt{n}(n-1)} \sum_{1 \leq i < j \leq n} \frac{g_2(X_i, X_j)}{\xi_1}\right] \\
 1382 &= \frac{it(r-1)}{\sqrt{n}(n-1)} \cdot \binom{n}{2} \cdot \varphi_n^{n-2}(t) \cdot \mathbb{E}\left[e^{\frac{it(g_1(X_1)+g_1(X_2))}{\sqrt{n}\xi_1}} \cdot \frac{g_2(X_1, X_2)}{\xi_1}\right] \\
 1383 &= \frac{it(r-1)\sqrt{n}}{2} \cdot \varphi_n^{n-2}(t) \cdot \mathbb{E}\left[\frac{g_2(X_1, X_2)}{\xi_1} + \frac{it(g_1(X_1) + g_1(X_2))g_2(X_1, X_2)}{\sqrt{n} \cdot \xi_1^2}\right. \\
 1384 &\quad \left. - \frac{t^2\{g_1^2(X_1) + 2g_1(X_1)g_1(X_2) + g_1^2(X_2)\} \cdot g_2(X_1, X_2)}{2n\xi_1^3} + \right] + O\left(n^{-1} \cdot e^{-t^2/4} \cdot \text{Poly}(t)\right) \\
 1385 &= \frac{it(r-1)\sqrt{n}}{2} \cdot \varphi_n^{n-2}(t) \cdot \mathbb{E}\left[\frac{g_2(X_1, X_2)}{\xi_1} + \frac{2itg_1(X_1)g_2(X_1, X_2)}{\sqrt{n} \cdot \xi_1^2}\right. \\
 1386 &\quad \left. - \frac{t^2\{g_1^2(X_1) + g_1(X_1)g_1(X_2)\} \cdot g_2(X_1, X_2)}{n\xi_1^3} + \right] + O\left(n^{-1} \cdot e^{-t^2/4} \cdot \text{Poly}(t)\right) \\
 1387 &= \frac{-it^3(r-1)}{2\sqrt{n} \cdot \xi_1^3} \cdot \varphi_n^{n-2}(t) \cdot \mathbb{E}[g_1(X_1)g_1(X_2) \cdot g_2(X_1, X_2)] + O\left(n^{-1} \cdot e^{-t^2/4} \cdot \text{Poly}(t)\right) \\
 (8.69) & \\
 1388 &= e^{-t^2/2} \cdot \frac{-it^3(r-1)}{2\sqrt{n} \cdot \xi_1^3} \cdot \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] + O\left(n^{-1} \cdot e^{-t^2/4} \cdot \text{Poly}(t)\right)
 \end{aligned}$$

1389 We use the approximation to  $\delta_n$  given by Lemma 3.1-(d). When we use it here, we may  
 1390 ignore any  $\tilde{O}_p(n^{-1} \log n)$  remainder term, which is justified by Lemma 8.2 in the real  
 1391 domain, not the frequency domain that characteristic function works with. We thus have

$$\begin{aligned}
 1392 \quad \mathbb{E}\left[e^{itU_n^\#} \cdot it \left(-\frac{1}{2}U_n^\# \cdot \delta_n\right)\right] &= -\frac{1}{2}it \cdot \mathbb{E}\left[e^{it\frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n}\xi_1}} \cdot \left\{\frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}\right\}\right. \\
 (8.70) & \\
 1393 &\quad \left. \cdot \left(\frac{\sum_{j=1}^n \{g_1^2(X_j) - \xi_1^2\}}{n\xi_1^2} + \frac{2(r-1)\sum_{i=1}^n \sum_{j \neq i} g_1(X_i)g_2(X_i, X_j)}{n(n-1)\xi_1^2}\right) + \tilde{O}_p(n^{-1} \cdot \log^{3/2} n)\right]
 \end{aligned}$$

1394 We consider the expression into two parts by the two terms inside the parenthesis on the  
 1395 RHS of the equation, and inspect them respectively. Ignoring the  $\tilde{O}_p(n^{-1} \cdot \text{Polylog}(n))$   
 1396 remainder, for the first part, we have

$$\begin{aligned}
 1397 \quad &-\frac{1}{2}it \cdot \mathbb{E}\left[e^{it\frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n}\xi_1}} \cdot \left\{\frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}\right\} \cdot \left(\frac{\sum_{j=1}^n \{g_1^2(X_j) - \xi_1^2\}}{n\xi_1^2}\right)\right] \\
 (8.71) & \\
 1398 &= -\frac{1}{2}it \cdot \mathbb{E}\left[e^{it\frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n}\xi_1}} \cdot \left\{\sum_{i=1}^n \frac{g_1(X_i)(g_1^2(X_i) - \xi_1^2)}{n\sqrt{n} \cdot \xi_1^3} + \sum_{\substack{i, j \in \{1, \dots, n\} \\ i \neq j}} \frac{g_1(X_i)(g_1^2(X_j) - \xi_1^2)}{n\sqrt{n} \cdot \xi_1^3}\right\}\right]
 \end{aligned}$$

Further breaking the RHS down and handle the two summations in the fancy bracket separately, we have

$$\begin{aligned}
 & -\frac{1}{2}\text{it} \cdot \mathbb{E} \left[ e^{\text{it} \frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}} \cdot \left\{ \sum_{i=1}^n \frac{g_1(X_i) (g_1^2(X_i) - \xi_1^2)}{n\sqrt{n} \cdot \xi_1^3} \right\} \right] \\
 & = -\frac{1}{2}\text{it} \cdot \varphi_n^{n-1}(t) \cdot n \cdot \mathbb{E} \left[ \left\{ 1 + \frac{\text{it} \cdot g_1(X_1)}{\sqrt{n} \cdot \xi_1} \right\} \cdot \left\{ \frac{g_1(X_1) (g_1^2(X_1) - \xi_1^2)}{n\sqrt{n} \cdot \xi_1^3} \right\} \right] \\
 & \quad + O\left(n^{-1} \cdot e^{-t^2/4t^2} \cdot \text{Poly}(t)\right) \\
 (8.72) \quad & = -\frac{1}{2} \cdot \frac{\text{it} \varphi_n^{n-1}(t)}{\sqrt{n} \cdot \xi_1^3} \cdot \mathbb{E} [g_1^3(X_1)] + O\left(n^{-1} \cdot e^{-t^2/4t^2} \cdot \text{Poly}(t)\right)
 \end{aligned}$$

and

$$\begin{aligned}
 & -\frac{1}{2}\text{it} \cdot \mathbb{E} \left[ e^{\text{it} \frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}} \cdot \left\{ \sum_{\substack{i,j \in \{1, \dots, n\} \\ i \neq j}} \frac{g_1(X_i) (g_1^2(X_j) - \xi_1^2)}{n\sqrt{n} \cdot \xi_1^3} \right\} \right] \\
 & = -\frac{1}{2}\text{it} \cdot \varphi_n^{n-2}(t) \cdot n(n-1) \cdot \mathbb{E} \left[ \left\{ 1 + \text{it} \cdot \frac{g_1(X_1)}{\sqrt{n} \cdot \xi_1} - \frac{t^2 g_1^2(X_1)}{2n\xi_1^2} \right\} \right. \\
 & \quad \left. \left\{ 1 + \text{it} \cdot \frac{g_1(X_2)}{\sqrt{n} \cdot \xi_1} - \frac{t^2 g_1^2(X_2)}{2n\xi_1^2} \right\} \cdot \left\{ \frac{g_1(X_1) (g_1^2(X_2) - \xi_1^2)}{n\sqrt{n} \cdot \xi_1^3} \right\} \right] + O\left(n^{-1} \cdot e^{-t^2/4t^2} \cdot \text{Poly}(t)\right) \\
 & = \frac{1}{2}\text{it}^3 \cdot \varphi_n^{n-2}(t) \cdot n(n-1) \cdot \mathbb{E} \left[ \frac{g_1^2(X_1) g_1(X_2) \{g_1^2(X_2) - \xi_1^2\}}{n^2 \sqrt{n} \cdot \xi_1^5} \right] + O\left(n^{-1} \cdot e^{-t^2/4t^2} \cdot \text{Poly}(t)\right) \\
 (8.73) \quad & = \frac{1}{2} \frac{\text{it}^3 \varphi_n^{n-2}(t)}{\sqrt{n} \cdot \xi_1^3} \cdot \mathbb{E} [g_1^3(X_1)] + O\left(n^{-1} \cdot e^{-t^2/4t^2} \cdot \text{Poly}(t)\right)
 \end{aligned}$$

Now we calculate Part 2 of the RHS of (8.70). We have

$$\begin{aligned}
 & -\frac{1}{2}\text{it} \cdot \mathbb{E} \left[ e^{\text{it} \frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}} \cdot \left\{ \frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1} \right\} \cdot \left( \frac{2(r-1) \sum_{i=1}^n \sum_{j \neq i} g_1(X_i) g_2(X_i, X_j)}{n(n-1)\xi_1^2} \right) \right] \\
 & = -\frac{(r-1)\text{it}}{\xi_1^2} \cdot \mathbb{E} \left[ e^{\text{it} \frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}} \cdot \left\{ \frac{g_1(X_1) + g_1(X_2)}{\sqrt{n} \cdot \xi_1} \cdot g_1(X_1) g_2(X_1, X_2) \right\} \right] \\
 & \quad - \frac{(r-1)\text{it}}{\xi_1^2} (n-2) \cdot \mathbb{E} \left[ e^{\text{it} \frac{\sum_{i=1}^n g_1(X_i)}{\sqrt{n} \cdot \xi_1}} \cdot \left\{ \frac{g_1(X_3)}{\sqrt{n} \cdot \xi_1} \cdot g_1(X_1) g_2(X_1, X_2) \right\} \right] \\
 & = -\frac{(r-1)\text{it}}{\sqrt{n} \cdot \xi_1^3} \cdot \varphi_n^{n-2}(t) \cdot \mathbb{E} [g_1(X_1) g_1(X_2) g_2(X_1, X_2)] + O\left(n^{-1} \cdot e^{-t^2/4t^2} \cdot \text{Poly}(t)\right) \\
 & \quad - \frac{(r-1)\text{it}}{\sqrt{n} \cdot \xi_1^3} (n-2) \cdot \varphi_n^{n-3}(t) \cdot \mathbb{E} \left[ e^{\text{it} \frac{g_1(X_1)}{\sqrt{n} \cdot \xi_1}} \cdot \left\{ 1 + \frac{\text{it} g_1(X_2)}{\sqrt{n} \cdot \xi_1} - \frac{t^2 g_1^2(X_2)}{2n\xi_1^2} \right\} \right. \\
 & \quad \left. \cdot \left\{ 1 + \frac{\text{it} g_1(X_3)}{\sqrt{n} \cdot \xi_1} - \frac{t^2 g_1^2(X_3)}{2n\xi_1^2} \right\} \cdot g_1(X_1) g_2(X_1, X_2) g_1(X_3) \right]
 \end{aligned}$$

$$\begin{aligned}
1418 \quad &= -\frac{(r-1)\mathfrak{i}t}{\sqrt{n} \cdot \xi_1^3} \cdot \varphi_n^{n-2}(t) \cdot \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] + O\left(n^{-1} \cdot e^{-t^2/4} t^2 \cdot \text{Poly}(t)\right) \\
1419 \quad &- \frac{(r-1)\mathfrak{i}t}{\sqrt{n} \cdot \xi_1^3} (n-2) \cdot \varphi_n^{n-3}(t) \cdot \mathbb{E}\left[\frac{-t^2}{n\xi_1^2} \cdot g_1(X_1)g_1(X_2)g_2(X_1, X_2)g_1^2(X_3)\right] \\
(8.74) \quad & \\
1420 \quad &= \frac{(r-1)\mathfrak{i}(t^3-t)}{\sqrt{n} \cdot \xi_1^3} \cdot e^{-t^2/2} \cdot \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] + O\left(n^{-1} \cdot e^{-t^2/4} t^2 \cdot \text{Poly}(t)\right)
\end{aligned}$$

1421 Collecting terms (8.69), (8.73) and (8.74), we have

$$\begin{aligned}
1422 \quad &\mathbb{E}\left[e^{\mathfrak{i}t(U_n^\# + \tilde{\Delta}_n + \Delta_n - \frac{1}{2}U_n^\# \delta_n)}\right] \\
1423 \quad &= e^{-t^2/2} \cdot \left\{ 1 - \left( \frac{\mathbb{E}[g_1^3(X_1)]}{2} + (r-1)\mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right) \cdot \frac{\mathfrak{i}t}{\sqrt{n} \cdot \xi_1^3} \right. \\
1424 \quad &\quad \left. + \left( \frac{\mathbb{E}[g_1^3(X_1)]}{3} + \frac{(r-1)}{2}\mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right) \cdot \frac{\mathfrak{i}t^3}{\sqrt{n} \cdot \xi_1^3} \right\} \\
(8.75) \quad &+ O\left(n^{-1} \log n \cdot e^{-t^2/4} \cdot \text{Poly}(t)\right)
\end{aligned}$$

1426 The remainder term is clearly ignorable if plugged into the Esseen's smoothing lemma.  
1427 It only remains to deal with the  $\sigma_w^2 t^2 / (\rho_n \cdot n)$  term and the  $\sigma_w^4 t^4 / (\rho_n^2 \cdot n^2)$  term in (8.64).  
1428 By (8.61), we have

$$\begin{aligned}
1429 \quad &\mathbb{E}\left[e^{\mathfrak{i}t\tilde{T}_n} \cdot \frac{\sigma_w^2 t^2}{\rho_n \cdot n}\right] \\
1430 \quad &= \left[ e^{\mathfrak{i}t\tilde{T}_n} \left( \mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i) + \tilde{O}_p(n^{-1} \cdot \log n) \right) \right] \cdot \frac{t^2}{\rho_n \cdot n} \\
1431 \quad &= \mathbb{E}\left[e^{\mathfrak{i}t\tilde{T}_n}\right] \cdot \frac{\mathbb{E}[\sigma_w^2]t^2}{\rho_n \cdot n} + \mathbb{E}\left[e^{\mathfrak{i}t\tilde{T}_n} \cdot g_{\sigma;1}(X_1)\right] \cdot \frac{t^2}{\rho_n \cdot n} + O\left(\frac{t^2 \log n}{\rho_n \cdot n^2}\right)
\end{aligned}$$

1432 Now we discuss the three terms on the RHS. Term 1:

$$\begin{aligned}
1433 \quad &\int_0^{n^\epsilon} \left| \mathbb{E}\left[e^{\mathfrak{i}t\tilde{T}_n}\right] \cdot \frac{\mathbb{E}[\sigma_w^2]t^2}{\rho_n \cdot n} \cdot \frac{1}{t} \right| dt = \int_0^{n^\epsilon} O\left(e^{-t^2/4} \cdot \text{Poly}(t)\right) \cdot (\rho_n \cdot n)^{-1} \\
1434 \quad &= O\left((\rho_n \cdot n)^{-1}\right)
\end{aligned}$$

Term 2: by mimicking the derivations in our (8.72) and also referring to (2.11) in [21], we see that

$$\left| \mathbb{E}\left[e^{\mathfrak{i}t\tilde{T}_n} \cdot g_{\sigma;1}(X_1)\right] \right| = O\left(e^{-t^2/4} \cdot \text{Poly}(t)\right)$$

1435 Therefore, it can be bounded in exactly the same way as term 1.

1436 For term 3, we have

$$\int_0^{n^\epsilon} \frac{t^2}{\rho_n \cdot n} \cdot \frac{1}{t} dt = (\rho_n \cdot n)^{-1} \cdot n^{2\epsilon-1} \leq (\rho_n \cdot n)^{-1}$$

1438 where recall that  $\epsilon < 1/2$ . The  $\sigma_w^4 t^4 / (\rho_n^2 \cdot n^2)$  term can be bounded exactly similarly and  
1439 we omit the proof here. This finishes the proof of Lemma 8.3-(d).

1440 Now we return to the proof of Theorem 3.1. Plugging the results of Lemma 8.3 back into  
 1441 Lemma 8.1 completes the proof of Theorem 3.1 with the assumption  $\rho_n = O((\log n)^{-1})$ .

1442 If Cramer's condition holds instead of the upper bound on  $\rho_n$ , then the derivation steps  
 1443 in (2.21)–(2.22) in [21] can be reproduced, where their  $t_N$  can be understood as  $n^{r_0}$  for any  
 1444 fixed  $r_0 \in (0, 1)$ . It would suffice for our purpose to use any  $r_0 \in (1/2, 1)$ . Notice that their “ $r$ ”  
 1445 has different meaning than ours. This extends the integrative range that our Lemma 8.3-(c)  
 1446 holds valid from the original range  $(n^\epsilon, C_1 \cdot n^{1/2})$  to  $(n^\epsilon, n^{r_0})$ , and we only need to prove  
 1447 Lemma 8.3-(b) on the integrative range  $(n^{r_0}, n)$  instead of  $(C_1 \cdot n^{1/2}, n)$ . Then our proof of  
 1448 Lemma 8.3-(b) can be revised into

$$\begin{aligned}
 1449 \quad & \left| \mathbb{E} \left[ e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right] \right| \leq \mathbb{E} \left[ \left| e^{it\tilde{T}_n} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right| \right] \\
 1450 \quad & = \mathbb{E} \left[ e^{-n^{-1} \sigma_w^2 t^2 / 2} \right] \leq \mathbb{E} \left[ e^{-n^{2r_0-1} \cdot \mathbb{E}[\sigma_w^2] / 4} \right] + \mathbb{P} \left( \sigma_w^2 < \mathbb{E}[\sigma_w^2] / 4 \right) \\
 1451 \quad (8.76) \quad & \leq e^{-C_1 \cdot n^{2r_0-1}} + e^{-C_2 n} < n^{-2}
 \end{aligned}$$

1452 where in the second line we replaced  $\rho_n$  by 1 to majorize.

1453

□

1454 **4.4. Proof of Theorem 3.2.** It is easy to verify that

$$(8.77) \quad \tilde{O}_p(p_n) \tilde{O}_p(q_n) = \tilde{O}_p(p_n q_n), \quad \text{and} \quad \tilde{O}_p(p_n) + \tilde{O}_p(q_n) = \tilde{O}_p(p_n + q_n)$$

1455 We also easily have  $O(p_n) \tilde{O}_p(q_n) = \tilde{O}_p(p_n q_n)$  since  $O(\cdot)$  implies  $\tilde{O}_p(\cdot)$ , but it is not guar-  
 1456 anteed that  $O_p(p_n) \tilde{O}_p(q_n) = \tilde{O}_p(p_n q_n)$  if the distribution of  $O_p(p_n)$  is heavy tailed. The  
 1457 presence of edge-wise observational errors introduces extra technical complications to the  
 1458 proof of Theorem 3.2 beyond the analysis for empirical Edgeworth expansions for noise-  
 1459 less U-statistics such as [71, 96] and [110]. We shall carefully address this. By the proofs of  
 1460 Lemma 3.1-(c) and (d), and recall that  $\hat{\xi}_1^2 = n \hat{S}^2 / r$  and  $\xi_1^2 = n \sigma_n^2 / r$ , we have

$$\frac{(\hat{\xi}_1 + \xi_1)(\hat{\xi}_1 - \xi_1)}{\rho_n^{2s}} \asymp \frac{\hat{\xi}_1^2 - \xi_1^2}{n \sigma_n^2} \asymp \delta_n + \hat{\delta}_n = \tilde{O}_p(n^{-1/2} \log^{1/2} n)$$

1461 Then noticing that  $\hat{\xi}_1 / \xi_1 = 1 + \tilde{O}_p(1)$  and thus  $\hat{\xi}_1 \asymp \xi_1 \asymp \rho_n^s$  with probability at least  $1 -$   
 1462  $O(n^{-1})$ , we have  $\hat{\xi}_1 - \xi_1 = \tilde{O}_p(\rho_n^s \cdot n^{-1/2} \log^{1/2} n)$ . Therefore

$$\hat{\xi}_1^3 - \xi_1^3 = \tilde{O}_p(\rho_n^{3s} \cdot n^{-1/2} \log^{1/2} n)$$

implying that

$$\left| \frac{1}{\sqrt{n} \xi_1^3} - \frac{1}{\sqrt{n} \hat{\xi}_1^3} \right| = \tilde{O}_p \left( \frac{\log^{1/2} n}{\rho_n^{3s} \cdot n} \right).$$

1463 Recall that  $\|F_{\hat{T}_n}(x) - G_n(x)\|_\infty = O(\mathcal{M}(\rho_n, n; R))$ , where

$$\begin{aligned}
 1464 \quad G_n(x) &= \Phi(x) + \frac{\varphi(x)}{\sqrt{n} \cdot \xi_1^3} \cdot \left\{ \left( \frac{x^2}{3} + \frac{1}{6} \right) \mathbb{E}[g_1^3(X_1)] \right. \\
 1465 \quad &\quad \left. + \frac{r-1}{2} (x^2 + 1) \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\}.
 \end{aligned}$$

1466 As a result, in order to prove  $\|\hat{G}_n(x) - G_n(x)\|_\infty = \tilde{O}_p(\mathcal{M}(\rho_n, n; R))$ , it suffices to show  
 1467 that

$$1468 \quad \max \left\{ \left| \hat{\mathbb{E}}g_1^3(X_1) - \mathbb{E}g_1^3(X_1) \right| \right.$$

$$\begin{aligned}
1469 \quad & , \left| \widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] - \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\} \\
1470 \quad & = \begin{cases} \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1/2} \log^{1/2} n) & \text{if } R \text{ is acyclic} \\ \tilde{O}_p(\rho_n^{3s-2/r} \cdot n^{-1/2} \log^{1/2} n) & \text{if } R \text{ is cyclic} \end{cases}
\end{aligned}$$

where we used the fact that  $\sup_{x \in \mathbb{R}} |x|^3 \varphi(x) = O(1)$ . We will show that the empirical moments  $\widehat{\mathbb{E}}[g_1^3(X_1)]$  and  $\widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$  converge to  $\mathbb{E}[g_1^3(X_1)]$  and  $\mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$ , respectively, at rates no slower than  $\tilde{O}_p(\rho_n^{3s-0.5} \cdot n^{-1/2} \log^{1/2} n)$  for both acyclic and cyclic cases under respective network sparsity conditions. The convergence of  $\widehat{\mathbb{E}}[g_1^3(X_1)]$  to  $\mathbb{E}[g_1^3(X_1)]$  can be established using (8.43). Recall the definitions of  $\widehat{a}_i$  and  $a_i$  from (8.26) and (8.27),

$$\widehat{\mathbb{E}}[g_1^3(X_1)] = \frac{1}{n} \sum_{i=1}^n (\widehat{a}_i - \widehat{U}_n)^3 \quad \text{and} \quad \mathbb{E}[g_1^3(X_1)] = \mathbb{E} \left[ (\mathbb{E}[h(X_1, \dots, X_r) | X_1] - \mu_n)^3 \right].$$

1471 Observe that

$$\begin{aligned}
1472 \quad & \left| \widehat{\mathbb{E}}[g_1^3(X_1)] - \mathbb{E}[g_1^3(X_1)] \right| \leq \left| \sum_{i=1}^n (\widehat{a}_i - \widehat{U}_n)^3 - \sum_{i=1}^n (a_i - \mu_n)^3 \right| / n \\
1473 \quad & + \left| \sum_{i=1}^n (a_i - \mu_n)^3 / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_1] - \mu_n)^3 \right|
\end{aligned}$$

(8.78)

$$1474 \quad = \left| \sum_{i=1}^n (a_i - \mu_n)^3 / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_1] - \mu_n)^3 \right| + \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n)$$

1475 where the last inequality is due to the facts  $a_i \asymp \mu_n \asymp \rho_n^s$ ,  $|\widehat{a}_i - a_i| = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$

1476 and  $|\widehat{U}_n - \mu_n| = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$  due to the proof of Lemma 3.1 (a), (b) and (c).

1477 Moreover, we have

$$\begin{aligned}
1478 \quad & \left| \sum_{i=1}^n (a_i - \mu_n)^3 / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_1] - \mu_n)^3 \right| \\
1479 \quad & \leq \left| \sum_{i=1}^n a_i^3 / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_i])^3 \right| \\
1480 \quad & + \rho_n^s \cdot O \left( \left| \sum_{i=1}^n a_i^2 / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_i])^2 \right| \right) \\
1481 \quad (8.79) \quad & + \rho_n^{2s} \cdot O \left( \left| \sum_{i=1}^n a_i / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_i]) \right| \right)
\end{aligned}$$

Recall the definition of  $a_i$  and notice that it is a U-statistic of order  $r-1$  conditioned on  $X_i$ . By the standard concentration inequality of U-statistic [97], we have

$$|a_i - \mathbb{E}[h(X_1, \dots, X_r) | X_i]| = \tilde{O}_p(\rho_n^s \cdot n^{-1/2} \log^{1/2} n).$$

By decomposing  $a_i = (a_i - \mathbb{E}[h(X_1, \dots, X_r) | X_i]) + \mathbb{E}[h(X_1, \dots, X_r) | X_i]$ , we have

$$\rho_n^{2s} \cdot O \left( \left| \sum_{i=1}^n a_i / n - \mathbb{E}(\mathbb{E}[h(X_1, \dots, X_r) | X_i]) \right| \right) = \tilde{O}_p(\rho_n^{3s} \cdot n^{-1/2} \log^{1/2} n)$$

where we used the facts  $\{\mathbb{E}[h(X_1, \dots, X_r)|X_i]\}_{i=1}^n$  are i.i.d. random variables so that

$$\left| n^{-1} \sum_{i=1}^n \mathbb{E}[h(X_1, \dots, X_r)|X_i] - \mathbb{E}[h(X_1, \dots, X_r)] \right| = \tilde{O}_p(\rho_n^{3s} n^{-1/2} \log^{1/2} n).$$

1482 By a similar strategy, we can prove that the bound  $\tilde{O}_p(\rho_n^{3s} \cdot n^{-1/2} \log^{1/2} n)$  also holds for the  
 1483 other two terms in RHS of (8.79). Together with (8.78), we conclude that

$$(8.80) \quad \left| \widehat{\mathbb{E}}g_1^3(X_1) - \mathbb{E}g_1^3(X_1) \right| = \tilde{O}_p(\rho_n^{3s-0.5} \cdot n^{-1/2} \log^{1/2} n).$$

1484 The proof of the convergence of  $\widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$ , however, needs separate care.  
 1485 Recall that

$$1486 \quad \widehat{g}_1(X_i) := \frac{1}{\binom{n-1}{r-1}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(A_{i, i_1, \dots, i_{r-1}}) - \widehat{U}_n = \widehat{a}_i - \widehat{U}_n$$

$$1487 \quad \widehat{g}_2(X_i, X_j) := \frac{1}{\binom{n-2}{r-2}} \sum_{\substack{1 \leq i_1 < \dots < i_{r-2} \leq n \\ i_1, \dots, i_{r-2} \neq i, j}} h(A_{i, j, i_1, \dots, i_{r-2}}) - \widehat{U}_n - \widehat{g}_1(X_i) - \widehat{g}_1(X_j)$$

1488 Unlike that  $\widehat{g}_1(X_i)$  converges to the corresponding  $g_1(X_i)$ , the randomness in  $h(A_{i, j, i_1, \dots, i_{r-2}})$   
 1489 introduced by the edge  $A_{ij}$  is not suppressed by an average over  $\{i_1, \dots, i_{r-2} : i_1, \dots, i_{r-2} \neq$   
 1490  $i, j\}$ . Therefore, the convergence of  $\widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$  has to be discussed as  
 1491 a whole. We first show that given  $W$ ,  $\widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$  converges to its  
 1492 ‘‘population-sample’’ version replacing  $A$  by  $W$  in its definition, then show the convergence  
 1493 of that version to the eventual expectation form. Observe that

$$1494 \quad \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \widehat{g}_1(X_i)\widehat{g}_1(X_j)\widehat{g}_2(X_i, X_j) - \mathbb{E}g_1(X_1)g_1(X_2)g_2(X_1, X_2)$$

$$1495 \quad = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} [\widehat{g}_1(X_i)\widehat{g}_1(X_j)\widehat{g}_2(X_i, X_j) - g_1(X_i)g_1(X_j)g_2(X_i, X_j)]$$

$$1496 \quad + \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} g_1(X_i)g_1(X_j)g_2(X_i, X_j) - \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)].$$

1497 It is easy to bound the second term. By the definition of  $g_1(X_i), g_2(X_i, X_j)$ , we notice that  
 1498 clearly  $\binom{n}{2}^{-1} \sum_{1 \leq i < j \leq n} g_1(X_i)g_1(X_j)g_2(X_i, X_j)$  is a degree-two U-statistic. By the stan-  
 1499 dard concentration inequality of U-statistic [97],

$$1500 \quad \left| \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} g_1(X_i)g_1(X_j)g_2(X_i, X_j) - \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right|$$

$$1501 \quad = \tilde{O}_p(\rho_n^{3s} n^{-1/2} \log^{1/2} n)$$

1502 where we used the fact  $g_1(X_i)g_1(X_j)g_2(X_i, X_j) = O(\rho_n^{3s})$  a.s. Therefore, it suffices to upper  
 1503 bound

$$(8.81) \quad \mathfrak{R}_1 := \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} [\widehat{g}_1(X_i)\widehat{g}_1(X_j)\widehat{g}_2(X_i, X_j) - g_1(X_i)g_1(X_j)g_2(X_i, X_j)].$$

1504 The convergence of  $\widehat{g}_1(X_i)$  to  $g_1(X_i)$  is straightforward. Indeed,

$$1505 \quad \widehat{g}_1(X_i) - g_1(X_i) = \widehat{a}_i - \mathbb{E}[h(X_1, \dots, X_r)|X_i] + (\mu_n - \widehat{U}_n).$$

1506 Recall from Lemma 3.1(a), (b) and (c),  $|\hat{U}_n - \mu_n| = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$ . We then  
 1507 prove the first term on RHS of above equation. Clearly,

$$1508 \quad |\hat{a}_i - \mathbb{E}[h(X_1, \dots, X_r)|X_i]| \leq |\hat{a}_i - a_i| + |a_i - \mathbb{E}[h(X_1, \dots, X_r)|X_i]|$$

$$1509 \quad = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$$

1510 where the last inequality is due to the bounds of  $|\hat{a}_i - a_i|$  and  $|a_i - \mathbb{E}[h(X_1, \dots, X_r)|X_i]|$  as  
 1511 shown above. Therefore, conditioned on  $X_i$ , we have  $|\hat{g}_1(X_i) - g_1(X_i)| = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$ .  
 1512 Now, we re-express  $\mathfrak{K}_1$  from (8.81) as

$$1513 \quad \mathfrak{K}_1 = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) [\hat{g}_2(X_i, X_j) - g_2(X_i, X_j)]$$

$$1514 \quad + \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} [\hat{g}_1(X_i) \hat{g}_1(X_j) g_2(X_i, X_j) - g_1(X_i) g_1(X_j) g_2(X_i, X_j)]$$

$$1515 \quad = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) [\hat{g}_2(X_i, X_j) - g_2(X_i, X_j)] + \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n),$$

1516 where we used the fact  $|g_1(X_i)| = O(\rho_n^s)$ , a.s. It suffices to bound the first term on RHS.  
 1517 Define

$$1518 \quad (8.82) \quad \hat{a}_{ij} := \frac{1}{\binom{n-2}{r-2}} \sum_{\substack{1 \leq i_1 < i_2 < \dots < i_{r-2} \leq n \\ i_1, \dots, i_{r-2} \neq i, j}} h(A_{i,j,i_1,i_2,\dots,i_{r-2}})$$

$$1519 \quad a_{ij} := \frac{1}{\binom{n-2}{r-2}} \sum_{\substack{1 \leq i_1 < i_2 < \dots < i_{r-2} \leq n \\ i_1, \dots, i_{r-2} \neq i, j}} h(W_{i,j,i_1,i_2,\dots,i_{r-2}}).$$

1520 Then we can re-express the  $\hat{g}_2(X_i, X_j) - g_2(X_i, X_j)$  factor as follows

$$1521 \quad \hat{g}_2(X_i, X_j) - g_2(X_i, X_j) = (\hat{a}_{ij} - a_{ij}) + (a_{ij} - \mathbb{E}[h(X_1, \dots, X_r)|X_i, X_j])$$

$$1522 \quad - (\hat{U}_n - \mu_n) - (\hat{g}_1(X_i) - g_1(X_i)) - (\hat{g}_1(X_j) - g_1(X_j)).$$

1523 Similarly to our earlier derivations, using the concentration of U-statistics, we have  $(a_{ij} -$   
 1524  $\mathbb{E}[h(X_1, \dots, X_r)|X_i, X_j]) = \tilde{O}_p(\rho_n^s n^{-1/2} \log^{1/2} n)$ . Since  $\hat{U}_n - \mu_n = \tilde{O}_p(\rho_n^{s-1/2} n^{-1/2} \log^{1/2} n)$   
 1525 and  $\hat{g}_1(X_i) - g_1(X_i) = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$ , we have

$$1526 \quad \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) [\hat{g}_2(X_i, X_j) - g_2(X_i, X_j)]$$

$$1527 \quad = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) (\hat{a}_{ij} - a_{ij}) + \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n).$$

1528 Therefore, we have

$$1529 \quad \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) \hat{g}_2(X_i, X_j) - \mathbb{E}g_1(X_1)g_1(X_2)g_2(X_1, X_2)$$

$$1530 \quad = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) (\hat{a}_{ij} - a_{ij}) + \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n).$$

1531 Recall the definitions of  $\hat{a}_i$  and  $a_i$  from (8.26) and (8.27). We write

$$\begin{aligned}
 1532 \quad & \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) (\hat{a}_{ij} - a_{ij}) = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i \hat{a}_j (\hat{a}_{ij} - a_{ij}) \\
 1533 \quad & - \frac{2}{n} \sum_{1 \leq i \leq n} \hat{U}_n \hat{a}_i (\hat{a}_i - a_i) + \hat{U}_n^2 (\hat{U}_n - U_n) \\
 1534 \quad & = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i \hat{a}_j (\hat{a}_{ij} - a_{ij}) + \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n)
 \end{aligned}$$

1535 where the last equation is due to  $a_i \asymp U_n \asymp \rho_n^s$  a.s.,  $|\hat{a}_i - a_i| = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)$ ,  
 1536  $|\hat{U}_n - U_n| = \tilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1} \log^{1/2} n)$  due to Lemma 3.1 (b). Therefore,

$$\begin{aligned}
 1537 \quad & \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) \hat{g}_2(X_i, X_j) - \mathbb{E}[g_1(X_1) g_1(X_2) g_2(X_1, X_2)] \\
 1538 \quad (8.83) \quad & = \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i \hat{a}_j (\hat{a}_{ij} - a_{ij}) + \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n).
 \end{aligned}$$

1539 It remains to bound the first term on RHS. We rewrite it as

$$\begin{aligned}
 1540 \quad & \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i \hat{a}_j (\hat{a}_{ij} - a_{ij}) = \frac{1}{n(n-1)} \sum_{1 \leq i \neq j \leq n} \hat{a}_i \hat{a}_j (\hat{a}_{ij} - a_{ij}) \\
 1541 \quad (8.84) \quad & = \frac{1}{n} \sum_{i=1}^n \hat{a}_i \cdot \left( \frac{1}{n-1} \sum_{j \neq i} \hat{a}_j (\hat{a}_{ij} - a_{ij}) \right).
 \end{aligned}$$

1542 We then establish the upper bound for  $\sum_{j \neq i} \hat{a}_j (\hat{a}_{ij} - a_{ij}) / (n-1)$  for each fixed  $i$ . We have

$$\begin{aligned}
 1543 \quad & \frac{1}{n-1} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} (\hat{a}_j - a_j) (\hat{a}_{ij} - a_{ij}) \\
 1544 \quad & = \frac{1}{(n-1)^2} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} \sum_{\substack{1 \leq i' \leq n \\ i' \neq j}} (\hat{a}_{i'j} - a_{i'j}) (\hat{a}_{ij} - a_{ij}) \\
 1545 \quad (8.85) \quad & = \frac{1}{(n-1)^2} \left\{ \sum_{\substack{1 \leq j \leq n \\ j \neq i}} (\hat{a}_{ij} - a_{ij})^2 + \sum_{\substack{1 \leq \{i', j\} \leq n \\ i' \neq i \\ j \neq i, i'}} (\hat{a}_{i'j} - a_{i'j}) (\hat{a}_{ij} - a_{ij}) \right\}
 \end{aligned}$$

1546 Similar to the derivation of (8.43) by expanding  $\sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ \{i, j\} \subset \{i_1, \dots, i_r\}}} h(A_{i_1, \dots, i_r})$ , we have

$$\begin{aligned}
 1547 \quad & \hat{a}_{ij} = \hat{\Theta}_{ij} \eta_{ij} + \frac{1}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i, j}} \left( \hat{\Theta}_{i, k; i, j} \eta_{ik} + \hat{\Theta}_{j, k; i, j} \eta_{jk} \right) \\
 1548 \quad (8.86) \quad & + a_{ij} + \tilde{O}_p(\rho_n^{s-1} \cdot n^{-1} \log n)
 \end{aligned}$$

where

$$\max \{ |\hat{\Theta}_{ij}|, |\hat{\Theta}_{i, k; i, j}|, |\hat{\Theta}_{j, k; i, j}| \} \leq \rho_n^{s-1}, a.s.$$

1549 We note that, similarly as the derivation of (8.43), the bound (8.86) holds under the sparsity  
 1550 condition  $\rho_n = \omega(n^{-1})$  for acyclic  $R$  and  $\rho_n = \omega(n^{-2/r})$  for cyclic  $R$ .

1551 Now we discuss the two terms on the RHS of (8.85). For term 1 on the RHS of (8.85), we  
1552 have

$$\begin{aligned}
1553 \quad & \frac{1}{(n-1)^2} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} (\hat{a}_{ij} - a_{ij})^2 \\
1554 \quad & = \frac{1}{(n-1)^2} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} \left\{ \dot{\Theta}_{i,j} \eta_{ij} + \frac{1}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i,j}} \left( \dot{\Theta}_{i,k;i,j} \eta_{ik} + \dot{\Theta}_{j,k;i,j} \eta_{jk} \right) \right\}^2 \\
1555 \quad & \quad + \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-2} \log n) \\
1556 \quad & \asymp n^{-2} \left\{ \sum_{\substack{1 \leq j \leq n \\ j \neq i}} \dot{\Theta}_{i,j}^2 \eta_{ij}^2 + \frac{2}{n-2} \sum_{\substack{1 \leq \{j,k\} \leq n \\ j \neq i \\ k \neq i,j}} \dot{\Theta}_{i,j} \eta_{ij} \left( \dot{\Theta}_{i,k;i,j} \eta_{ik} + \dot{\Theta}_{j,k;i,j} \eta_{jk} \right) \right. \\
1557 \quad (8.87) \quad & \left. + \frac{1}{(n-2)^2} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} \left( \sum_{\substack{1 \leq k \leq n \\ k \neq i,j}} \left( \dot{\Theta}_{i,k;i,j} \eta_{ik} + \dot{\Theta}_{j,k;i,j} \eta_{jk} \right) \right)^2 \right\} + \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-2} \log n)
\end{aligned}$$

1558 Now we bound each term on the RHS of (8.87). Inspecting the expectation of term 1 on  
1559 the RHS of (8.87) and using Bernstein inequality, we know it is  $\tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1} + \rho_n^{2s-3/2} \cdot$   
1560  $n^{-3/2} \log^{1/2} n)$ . Term 2 on the RHS of (8.87) is mean zero so we can focus on the concen-  
1561 tration. Its  $\eta_{ij} \eta_{ik}$  part can be bounded by inspecting the concentration averaging over  $j$  and  
1562 over  $k$ , respectively, and see that this part is bounded as  $\tilde{O}_p(\rho_n^{2s-2} \cdot n^{-1} (\rho_n n^{-1/2} \log^{1/2} n)^2)$ ,  
1563 and this upper bound is dominated by the bound of term 1, thus it is ignorable. Using Theo-  
1564 rem 8.1, the  $\eta_{ij} \eta_{jk}$  part of term 2 can be bounded as follows

$$1565 \quad n^{-3} \sum_{\substack{1 \leq \{j,k\} \leq n \\ j \neq i \\ k \neq i,j}} \dot{\Theta}_{i,j} \dot{\Theta}_{j,k;i,j} \eta_{ij} \eta_{jk} = \tilde{O}_p \left( n^{-3} \cdot \rho_n^{2s-2} \cdot \max \left\{ \underbrace{\sqrt{\rho_n^2 \cdot n^2 \log n}}_{\text{"Variance"}}, \underbrace{\rho_n \cdot n \log n}_{\text{"}\Xi_1\text{"}} \right\} \right)$$

1566 and is thus ignorable. Now noticing that each  $\eta$  is bounded by 1, using Bernstein's inequality,  
1567 term 3 on the RHS of (8.87) is  $\tilde{O}_p(n^{-4} \cdot \rho_n^{2s-1} \cdot n^2) = \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-2})$  and thus ignorable.  
1568 Therefore, term 1 on the RHS of (8.85) is  $\tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1} \log n + \rho_n^{2s-3/2} \cdot n^{-3/2} \log^{1/2} n)$ .

1569 Now we bound term 2 on the RHS of (8.85). By a similar treatment, we have

$$\begin{aligned}
1570 \quad & \frac{1}{(n-1)^2} \sum_{\substack{1 \leq \{i',j\} \leq n \\ i' \neq i \\ j \neq i,i'}} (\hat{a}_{i'j} - a_{i'j})(\hat{a}_{ij} - a_{ij}) \\
1571 \quad & = \frac{1}{(n-1)^2} \sum_{\substack{1 \leq \{i',j\} \leq n \\ i' \neq i \\ j \neq i,i'}} \left\{ \dot{\Theta}_{i',j} \eta_{i'j} + \frac{1}{n-2} \sum_{\substack{1 \leq k_1 \leq n \\ k_1 \neq i',j}} \left( \dot{\Theta}_{i',k_1;i',j} \eta_{i'k_1} + \dot{\Theta}_{j,k_1;i',j} \eta_{jk_1} \right) \right\}
\end{aligned}$$

$$\begin{aligned}
 & \cdot \left\{ \dot{\Theta}_{ij} \eta_{ij} + \frac{1}{n-2} \sum_{\substack{1 \leq k_2 \leq n \\ k_2 \neq i, j}} \left( \dot{\Theta}_{i, k_2; i, j} \eta_{ik_2} + \dot{\Theta}_{j, k_2; i, j} \eta_{jk_2} \right) \right\} + \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1} \log n) \\
 & \asymp n^{-2} \left[ \sum_{\substack{1 \leq \{i', j\} \leq n \\ i' \neq i \\ j \neq i, i'}} \dot{\Theta}_{i', j} \dot{\Theta}_{i, j} \eta_{i'j} \eta_{ij} + \frac{1}{n-2} \sum_{\substack{1 \leq \{i', j, k_2\} \leq n \\ i' \neq i \\ j \neq i, i' \\ k_2 \neq i, j}} \left( \dot{\Theta}_{i', j} \dot{\Theta}_{i, k_2; i, j} \eta_{i'j} \eta_{ik_2} + \dot{\Theta}_{i', j} \dot{\Theta}_{j, k_2; i, j} \eta_{i'j} \eta_{jk_2} \right) \right. \\
 & + \frac{1}{n-2} \sum_{\substack{1 \leq \{i', j, k_1\} \leq n \\ i' \neq i \\ j \neq i, i' \\ k_1 \neq i, j}} \left( \dot{\Theta}_{i, j} \dot{\Theta}_{i', k_1; i', j} \eta_{ij} \eta_{i'k_1} + \dot{\Theta}_{i, j} \dot{\Theta}_{j, k_1; i', j} \eta_{ij} \eta_{jk_1} \right) \\
 & \left. + \tilde{O}_p((\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n)^2) + \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1} \log n) \right]
 \end{aligned} \tag{8.88}$$

Now we bound the RHS of (8.88). Again, by Theorem 8.1, the first term is bounded by

$$\sum_{\substack{1 \leq \{i', j\} \leq n \\ i' \neq i \\ j \neq i, i'}} \dot{\Theta}_{i', j} \eta_{i'j} = \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1} \log n). \tag{8.89}$$

Terms 2 and 3 on the RHS of (8.88) can be bounded exactly similarly. Here we only present the bounding of term 2. We have

$$\frac{1}{(n-2)^3} \sum_{\substack{1 \leq \{i', j, k_2\} \leq n \\ i' \neq i \\ j \neq i, i' \\ k_2 \neq i, j}} \dot{\Theta}_{i', j} \dot{\Theta}_{i, k_2; i, j} \eta_{i'j} \eta_{ik_2} = \frac{1}{(n-2)^3} \sum_{1 \leq i' \leq n} \left( \sum_{j \neq i, i'} \dot{\Theta}_{i', j} \eta_{i'j} \sum_{k_2 \neq i, j} \dot{\Theta}_{i, k_2; i, j} \eta_{ik_2} \right) \tag{8.90}$$

$$= n^{-3} \rho_n^{2s-2} \tilde{O}_p((\rho_n^{1/2} n^{1/2} \log^{1/2} n)^2) = \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-2} \log n)$$

and using Theorem 8.1, we have

$$\begin{aligned}
 & \frac{1}{(n-2)^3} \sum_{\substack{1 \leq \{i', j, k_2\} \leq n \\ i' \neq i \\ j \neq i, i' \\ k_2 \neq i, j}} \dot{\Theta}_{i', j} \dot{\Theta}_{j, k_2; i, j} \eta_{i'j} \eta_{jk_2} \\
 & = n^{-3} \rho_n^{2s-2} \cdot \tilde{O}_p(\max\{\rho_n \cdot n^{3/2} \log^{1/2} n, \rho_n \cdot n \log n\}) = \tilde{O}_p(\rho_n^{2s-1} \cdot n^{-3/2} \log n)
 \end{aligned} \tag{8.91}$$

Collecting all results, we see that term 2 on the RHS of (8.85) is  $\tilde{O}_p(\rho_n^{2s-1} \cdot n^{-1} \log n + \rho_n^{2s-3/2} \cdot n^{-3/2} \log n)$ . We thus conclude that

$$\left| \frac{1}{n-1} \sum_{\substack{1 \leq j \leq n \\ j \neq i}} (\hat{a}_{ij} - a_{ij})(\hat{a}_j - a_j) \right| = \tilde{O}_p(\rho_n^{2s-1} n^{-1} \log n)$$

1585 under the given sparsity condition  $\rho_n = \omega(n^{-1/2})$ , which holds for both acyclic and cyclic  
1586  $R$ .

1587 Now we return to the main proof and continue (8.84). We have

$$\begin{aligned}
1588 \quad & \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i \hat{a}_j (\hat{a}_{ij} - a_{ij}) \\
1589 \quad &= \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i a_j (\hat{a}_{ij} - a_{ij}) + \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{a}_i (\hat{a}_j - a_j) (\hat{a}_{ij} - a_{ij}) \\
1590 \quad &= \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} a_i a_j (\hat{a}_{ij} - a_{ij}) + \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} (\hat{a}_i - a_i) a_j (\hat{a}_{ij} - a_{ij}) \\
1591 \quad &\quad + \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n + \rho_n^{3s-3/2} \cdot n^{-3/2} \log n) \\
1592 \quad &= \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n) + \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n + \rho_n^{3s-3/2} \cdot n^{-3/2} \log n) \\
1593 \quad (8.92) \quad &= \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n + \rho_n^{3s-3/2} \cdot n^{-3/2} \log n)
\end{aligned}$$

1594 where the second to last line is due to

$$\begin{aligned}
1595 \quad & \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} a_i a_j (\hat{a}_{ij} - a_{ij}) \\
1596 \quad &= \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} a_i a_j \left\{ \dot{\Theta}_{i,j} \eta_{ij} + \frac{1}{n-2} \sum_{\substack{1 \leq k \leq n \\ k \neq i,j}} \left( \dot{\Theta}_{i,k;i,j} \eta_{ik} + \dot{\Theta}_{j,k;i,j} \eta_{jk} \right) \right\} \\
1597 \quad &\quad + \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n) \\
1598 \quad &= \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} a_i a_j \dot{\Theta}_{i,j} \eta_{ij} + \frac{1}{\binom{n}{2}} \sum_{\substack{1 \leq \{i,j,k\} \leq n \\ i \neq j; j \neq k; k \neq i}} a_i a_j \frac{\dot{\Theta}_{i,k;i,j} \eta_{ik}}{n-2} + \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n) \\
1599 \quad (Bernstein) \quad &= \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n) + \frac{1}{\binom{n}{2}} \sum_{\substack{1 \leq \{i,k\} \leq n \\ i \neq k}} a_i \left( \sum_{\substack{1 \leq j \leq n \\ j \neq i,k}} a_j \frac{\dot{\Theta}_{i,k;i,j}}{n-2} \right) \eta_{ik} \\
1600 \quad (8.93) \quad &= \tilde{O}_p(\rho_n^{3s-1} \cdot n^{-1} \log n)
\end{aligned}$$

1601 Now we may conclude that

$$\begin{aligned}
1602 \quad & \left| \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_i) \hat{g}_1(X_j) \hat{g}_2(X_i, X_j) - \mathbb{E}[g_1(X_1) g_1(X_2) g_2(X_1, X_2)] \right| \\
1603 \quad (8.94) \quad &= \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n)
\end{aligned}$$

1604 This completes the proof of Theorem 3.2.

1605 **8.5. Proof of Theorem 3.3.** We will inherit the notation of  $\hat{a}_i$  from (8.27) in the proof of  
1606 Lemma 3.1. It suffices to show (3.15), which would then imply the closeness between  $F_{\hat{T}_n}$

1607 and  $F_{\hat{T}_n; \text{bootstrap}}$  by repeating our arguments for proving (8.54) and (8.55) using Lemma 8.2.

1608 Observe that

$$\begin{aligned}
 1609 \quad & \binom{n}{r} \cdot \hat{U}_n = \sum_{1 \leq i_1 < \dots < i_r \leq n} h(A_{i_1, \dots, i_r}) \\
 1610 \quad & \text{(For any } i) = \sum_{\substack{1 \leq i_1 < \dots < i_{r-1} \leq n \\ i_1, \dots, i_{r-1} \neq i}} h(A_{i, i_1, \dots, i_{r-1}}) + \sum_{\substack{1 \leq i_1 < \dots < i_r \leq n \\ i_1, \dots, i_r \neq i}} h(A_{i_1, \dots, i_r}) \\
 1611 \quad & = \binom{n-1}{r-1} \cdot \hat{a}_i + \binom{n-1}{r} \cdot \hat{U}_n^{(-i)}
 \end{aligned}$$

1612 Simplifying both sides, we have

$$1612 \quad (8.95) \quad \hat{U}_n^{(-i)} - \hat{U}_n = -\frac{r}{n-r} (\hat{a}_i - \hat{U}_n)$$

1613 Therefore,

$$\begin{aligned}
 1614 \quad & n \left( \hat{S}_n^2 - \hat{S}_{n; \text{jackknife}}^2 \right) \\
 1615 \quad & = \frac{r^2}{n} \sum_{i=1}^n (\hat{a}_i - \hat{U}_n)^2 - (n-1) \sum_{i=1}^n \left( \hat{U}_n^{(-i)} - \hat{U}_n \right)^2 \\
 1616 \quad & = \frac{1}{n} \sum_{i=1}^n \left[ r^2 (\hat{a}_i - \hat{U}_n)^2 - n(n-1) \cdot \frac{r^2}{(n-r)^2} (\hat{a}_i - \hat{U}_n)^2 \right] \\
 1617 \quad (8.96) \quad & = \frac{1}{n} \sum_{i=1}^n r^2 \left\{ 1 - \frac{n(n-1)}{(n-r)^2} \right\} (\hat{a}_i - \hat{U}_n)^2 = O(\hat{S}_n^2)
 \end{aligned}$$

where in the last line, recall that  $\hat{S}_n^2 := r^2 \sum_{i=1}^n (\hat{a}_i - \hat{U}_n)^2 / n^2$ . Therefore,

$$\hat{S}_n^2 - \hat{S}_{n; \text{jackknife}}^2 = O(\hat{S}_n^2/n) \implies |\hat{S}_n - \hat{S}_{n; \text{jackknife}}| = O(\hat{S}_n/n).$$

1618 This proves (3.15) and thus completes the proof of Theorem 3.3.

1619 **PROOF OF THEOREM 3.4.** It suffices to prove the Berry-Esseen bound for the normal ap-  
 1620 proximation. By definition,  $\|G_n(u) - \Phi(u)\|_\infty = O(n^{-1/2})$  and by the proof of Theorem 3.2,  
 1621 we know that  $\|\hat{G}_n(u) - G_n(u)\|_\infty = \tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \cdot n^{-1/2} = \tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \wedge o_p(1)$ .  
 1622 The proof is partitioned into two parts, for  $O(\mathcal{M}(\rho_n, n; R))$  and  $o(1)$  bounds, respectively.

1623 **Part I: proof of the  $O(\mathcal{M}(\rho_n, n; R))$  bound when  $\rho_n n > \log^{1/2} n$  (acyclic) or  $\rho_n^{r/2} n > \log^{1/2} n$  (cyclic).**

1624 We begin by recalling the decomposition of  $\hat{T}_n$  and inspect whether each component de-  
 1625 pends on  $\rho_n$  or not. Using Lemma 3.1 for the sparse regime, we have

$$\begin{aligned}
 1626 \quad & \hat{T}_n = \left\{ \underbrace{U_n^\# + \Delta_n}_{\text{No } \rho_n} + \underbrace{\hat{\Delta}_n}_{\text{Depend on } \rho_n} + \underbrace{\tilde{O}_p(n^{-1} \log^{3/2} n)}_{\text{No } \rho_n} \right\} \cdot \left( 1 + \underbrace{\delta_n}_{\text{No } \rho_n} + \underbrace{\hat{\delta}_n}_{\text{Depend on } \rho_n} \right) \\
 1627 \quad & = \left\{ U_n^\# + \tilde{O}_p(n^{-1/2} \log^{1/2} n) + \hat{\Delta}_n + \tilde{O}_p(n^{-1} \log^{3/2} n) \right\} \\
 1628 \quad & \cdot \left\{ 1 + \tilde{O}_p(n^{-1/2} \log^{1/2} n) + \tilde{O}_p(\mathcal{M}(\rho_n, n; R)) \right\}
 \end{aligned}$$

$$(8.97) \quad = U_n^\# + \widehat{\Delta}_n (1 + \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))) + \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$$

To see that last equality of (8.97), it is not difficult to prove that  $U_n^\# \cdot \widehat{\delta}_n$  is also  $\widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$  using the method for proving that  $\widehat{\delta}_n = \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$ , but due to its involvement we omit the proof here.

Now we use (8.97) for this proof. First, we discuss the term  $\widehat{\Delta}_n \cdot \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$ . By an ordinary Bernstein's inequality, we have

$$\begin{aligned} \mathbb{P}(|\check{\Delta}_n| > u) &\leq 2 \exp \left\{ -\frac{C_1 u^2 n^4}{C_2 n^2 \cdot \rho_n \cdot \rho_n^{-2} \cdot n + C_3 \rho_n^{-1} \cdot n^{1/2} \cdot u \cdot n^2} \right\} \\ &\leq 2 \exp \left\{ -C_4 (\rho_n \cdot n) \cdot u^2 \right\} \end{aligned}$$

Therefore,  $\check{\Delta}_n = \widetilde{O}_p((\rho_n \cdot n)^{-1/2} \log^{1/2} n)$ . Therefore, we have

$$\begin{aligned} &\widehat{\Delta}_n \cdot \widetilde{O}_p(\mathcal{M}(\rho_n, n; R)) \\ &= (\check{\Delta}_n + \check{R}_n) \cdot \widetilde{O}_p(\mathcal{M}(\rho_n, n; R)) \\ &= \left\{ \widetilde{O}_p((\rho_n \cdot n)^{-1/2} \log^{1/2} n) + \widetilde{O}_p(\mathcal{M}(\rho_n, n; R)) \right\} \cdot \widetilde{O}_p(\mathcal{M}(\rho_n, n; R)) \end{aligned}$$

Therefore, the term  $\widehat{\Delta}_n \cdot \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$  is ignorable compared to  $\widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$ . Thus, recalling  $\check{R}_n = \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$ , we have

$$\widehat{T}_n = U_n^\# + \check{\Delta}_n + \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$$

Now it only remains to show that

$$\|F_{U_n^\# + \check{\Delta}_n + \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))}(u) - \Phi(u)\|_\infty = O(\mathcal{M}(\rho_n, n; R))$$

Similar to the proof of Theorem 3.1, we are going to break this down into three steps. Recall the definition of  $\check{\Delta}_n$  from the proof of Theorem 3.1, we shall prove

$$(8.98) \quad \left\| F_{U_n^\# + \check{\Delta}_n + \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))}(u) - F_{U_n^\# + \check{\Delta}_n}(u) \right\|_\infty = \widetilde{O}_p(\mathcal{M}(\rho_n, n; R))$$

$$(8.99) \quad \left\| F_{U_n^\# + \check{\Delta}_n}(u) - F_{U_n^\#}(u) \right\|_\infty = O(\rho_n^{-1/2} \cdot n^{-1})$$

$$(8.100) \quad \left\| F_{U_n^\#}(u) - \Phi(u) \right\|_\infty = O((\rho_n \cdot n)^{-1} \log^{1/2} n)$$

We start from proving (8.100). Notice that this part of the proof only requires that  $\rho_n n > \log^{1/2} n$  regardless of the shape of the motif, since the asymptotic orders  $U_n^\# \asymp 1$  and  $\widehat{\Delta}_n \asymp (\rho_n \cdot n)^{-1/2}$  do not depend on the motif. The stronger condition  $\rho_n^{r/2} n > \log^{1/2} n$  is still necessary to deduce (8.98) from (8.99) and (8.100) using Lemma 8.2; a second reason is that the error bound  $\mathcal{M}(\rho_n, n; R)$  for cyclic motifs would not diminish to zero if  $\rho_n^{r/2} n \leq \log^{1/2} n$ . We are going to apply the Esseen's smoothing lemma on the interval  $t \in [-\rho_n \cdot n \log^{-1/2} n, \rho_n \cdot n \log^{-1/2} n]$ . The integral we shall need to bound is

$$(8.101) \quad \int_{-\rho_n \cdot n \log^{-1/2} n}^{\rho_n \cdot n \log^{-1/2} n} \left| \frac{\mathbb{E}[e^{it(U_n^\# + \check{\Delta}_n)}] - e^{-t^2/2}}{t} \right| dt$$

The following intermediate result in the proof of Lemma 8.3-(c) remains valid:

$$\begin{aligned} \mathbb{E} \left[ e^{it(U_n^\# + \check{\Delta}_n)} \right] &= \mathbb{E} \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1} \sigma_w^2 t^2 / 2} \right] \\ &= \mathbb{E} \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1} t^2 / 2} \{ \mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i) \} \cdot (1 + \widetilde{O}_p(\rho_n^{-1} n^{-2} \log n \cdot t^2)) \right] \end{aligned}$$

1656 For  $t \leq \rho_n \cdot n \log^{-1/2} n \ll n^{1/2}$ , the remainder's contribution to the integral (8.101) is

$$(8.102) \quad \int_0^{\rho_n \cdot n \log^{-1/2} n} \rho_n^{-1} n^{-2} \log n \cdot t^2 / t dt = O(\rho_n) \ll n^{-1/2}$$

1657 Therefore, for the rest of the proof in this part, we can directly ignore the remainder term's  
 1658 contribution according to (8.102). Now we bound the main part. Suppose  $C_0 > 0$  is a very  
 1659 large constant. We discuss two cases

- **Case 1:**  $\rho_n \cdot n \log^{-1/2} n \geq \{C_0 \log(\rho_n \cdot n)\}^{1/2}$ . In this case, we break the integral in (8.101) into two parts:

$$\int_0^{\{C_0 \log(\rho_n \cdot n)\}^{1/2}} + \int_{\{C_0 \log(\rho_n \cdot n)\}^{1/2}}^{\rho_n \cdot n \log^{-1/2} n}$$

1660 By (8.102), we can ignore the remainder. Similar to the intermediate step in the proof of  
 1661 Lemma 8.3-(d), using Section VI, Lemma 4 of [108], we have

$$\begin{aligned} 1662 & \mathbb{E} \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1}(t^2/2)\{\mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i)\}} \right] \\ 1663 & = e^{-(\rho_n \cdot n)^{-1}(t^2/2)\mathbb{E}[\sigma_w^2]} \cdot \mathbb{E} \left[ e^{it \sum_{i=1}^n g_1(X_i)/(\sqrt{n}\xi_1) - (\rho_n \cdot n)^{-1}t^2/(2n) \cdot \sum_{i=1}^n g_{\sigma;1}(X_i)} \right] \\ 1664 & = e^{-(\rho_n \cdot n)^{-1}(t^2/2)\mathbb{E}[\sigma_w^2]} \cdot \mathbb{E} \left[ e^{it \sum_{i=1}^n g_1(X_i)/(\sqrt{n}\xi_1)} \cdot \left( 1 + \tilde{O}_p \left( \frac{t^2 \log^{1/2} n}{\rho_n n^{3/2}} \right) \right) \right] \\ 1665 & = e^{-(\rho_n \cdot n)^{-1}(t^2/2)\mathbb{E}[\sigma_w^2]} \cdot \mathbb{E} \left[ e^{it \sum_{i=1}^n g_1(X_i)/(\sqrt{n}\xi_1)} \right] + O\left( \frac{t^2 \log^{1/2} n}{\rho_n n^{3/2}} \right) \\ 1666 & = e^{-(\rho_n \cdot n)^{-1}(t^2/2)\mathbb{E}[\sigma_w^2]} \cdot \mathbb{E} \left\{ e^{-t^2/2} + O(n^{-1/2}t^3 e^{-t^2/2}) \right\} + O\left( \frac{t^2 \log^{1/2} n}{\rho_n n^{3/2}} \right) \end{aligned}$$

1667 Therefore,

$$\begin{aligned} 1668 & \left| \mathbb{E} \left[ e^{itU_n^\#} \cdot e^{-C(\rho_n \cdot n)^{-1}t^2/2\{\mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i)\}} \right] - e^{-t^2/2} \right| \\ 1669 & \leq e^{-t^2/2} \left| e^{-C(\rho_n \cdot n)^{-1}t^2} - 1 \right| + O(n^{-1/2}t^3 e^{-t^2/2}) + O\left( \frac{t^2 \log^{1/2} n}{\rho_n n^{3/2}} \right) \\ 1670 & \leq e^{-t^2/2} \cdot O((\rho_n \cdot n)^{-1}t^2) + O(n^{-1/2}t^3 e^{-t^2/2}) + O\left( \frac{t^2 \log^{1/2} n}{\rho_n n^{3/2}} \right) \end{aligned}$$

1671 Consequently,

$$(8.103) \quad \int_0^{\{C_0 \log(\rho_n \cdot n)\}^{1/2}} \left| \frac{\mathbb{E}[e^{it(U_n^\# + \tilde{\Delta}_n)}] - e^{-t^2/2}}{t} \right| dt = O((\rho_n \cdot n)^{-1})$$

1672 where we recall (8.102) to simplify notation. For the second part of the integral, we can  
 1673 reproduce the steps in the proof of Theorem 3.1 and obtain

$$\begin{aligned} 1674 & \mathbb{E} \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1}(t^2/2)\{\mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i)\}} \right] \\ 1675 & = e^{-(\rho_n \cdot n)^{-1}(t^2/2)\mathbb{E}[\sigma_w^2]} \cdot \mathbb{E} \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1}t^2/2\{\frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i)\}} \right] \\ 1676 & \leq (1 - C_1 \cdot t^2/n)^n \leq e^{-C_2 t^2} \leq e^{-C_2 \cdot C_0 \log(\rho_n \cdot n)} \leq (\rho_n \cdot n)^{-2}. \end{aligned}$$

Therefore we have

$$\begin{aligned} & \int_{\{C_0 \log(\rho_n n)\}^{1/2}}^{(\rho_n n) \log^{-1/2} n} \left| \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1}(t^2/2)\{\mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i)\}} \right] \right| \cdot t^{-1} dt \\ & \leq \int_{\{C_0 \log(\rho_n n)\}^{1/2}}^{(\rho_n n) \log^{-1/2} n} \left| \left[ e^{itU_n^\#} \cdot e^{-(\rho_n \cdot n)^{-1}(t^2/2)\{\mathbb{E}[\sigma_w^2] + \frac{1}{n} \sum_{i=1}^n g_{\sigma;1}(X_i)\}} \right] \right| dt \\ & \leq \int_{\{C_0 \log(\rho_n n)\}^{1/2}}^{(\rho_n n) \log^{-1/2} n} (\rho_n \cdot n)^{-2} dt \leq (\rho_n \cdot n)^{-1} \end{aligned}$$

Moreover, we choose  $C_0 \geq 4$  so that

$$\begin{aligned} & \int_{\{C_0 \log(\rho_n n)\}^{1/2}}^{(\rho_n n) \log^{-1/2} n} \frac{e^{-t^2/2}}{t} dt \leq \int_{\{C_0 \log(\rho_n n)\}^{1/2}}^{(\rho_n n) \log^{-1/2} n} e^{-t^2/2} dt \\ & \leq (\rho_n n) \log^{-1/2}(n) \cdot e^{-(C_0/2) \cdot \log(\rho_n n)} \leq \frac{\rho_n n}{(\rho_n n)^{C_0/2}} \leq \frac{1}{\rho_n n}. \end{aligned}$$

Therefore, we have

$$(8.104) \quad \int_{\{C_0 \log(\rho_n n)\}^{1/2}}^{(\rho_n n) \log^{-1/2} n} \left| \frac{\mathbb{E}[e^{it(U_n^\# + \tilde{\Delta}_n)}] - e^{-t^2/2}}{t} \right| dt = O_p((\rho_n \cdot n)^{-1})$$

Combining (8.103) and (8.104) proves (8.101).

- **Case 2:**  $\rho_n \cdot n \log^{-1/2} n < \{C_0 \log(\rho_n \cdot n)\}^{1/2}$ . The proof in this case is even easier, since (8.103) remains valid and implies (8.101).

Plugging (8.101) back into the Esseen's smoothing lemma proves (8.100). Notice that the  $\log^{1/2} n$  factor in the eventual error bound comes from the second term on the RHS of (8.53).

Next, reproducing the proof (8.55), we prove (8.99) by combining (8.100) and Lemma 3.1-(b).

Finally, the proof of (8.98) is done by combining (8.55) and Lemma 8.2. The proof of this part is exactly similar to the proof of (8.54). This completes the proof of the  $O(\mathcal{M}(\rho_n, n; R))$  bound.

**Part II: proof of the  $o(1)$  bound when  $1 < \rho_n n \leq \log^{1/2} n$  (acyclic) or  $1 < \rho_n^{r/2} n \leq \log^{1/2} n$  (cyclic).**

The error bounds we derived in Part I of this proof focused on establishing finite sample error rates, and consequently need to bound the tail probability at the price of a  $\log^{1/2} n$  factor multiplied on the error bound. Taking the acyclic motif setting as an example, to counter the log factor in the error bound, we also need to assume  $\rho_n = \omega(n^{-1} \log^{1/2} n)$  rather than  $\rho_n = \omega(n^{-1})$ . For  $\rho_n : n^{-1} < \rho_n \leq n^{-1} \log^{1/2} n$ , despite establishing an explicit finite-sample error bound is still possible, the formula and derivation are rather complicated. For cleanness of presentation, in this paper, we slightly lower the goal and only aim at deriving uniform consistency. Consequently, the proof can be done by slightly varying the proof of the first part of Theorem 3.4. In this proof, we do not need to show an explicit error rate, so we do not need " $\tilde{O}_p$ " any more, and " $o_p$ " would suffice for our purpose. For the convenience of narration, we define

$$\tilde{\mathcal{M}}(\rho_n, n; R) := \begin{cases} (\rho_n \cdot n)^{-1/2} + n^{-1} \cdot \log^{3/2} n, & \text{For acyclic } R \\ (\rho_n^{r/2} \cdot n)^{-1/2} + n^{-1} \cdot \log^{3/2} n, & \text{For cyclic } R \end{cases}$$

We first present a variant of Lemma 3.1.

1697 LEMMA 8.4. *Under the conditions of Theorem 3.4, we have the following results:*

- 1698 (a) *Identical to Lemma 3.1-(a).*  
 1700 (b) *We have*

$$\widehat{\Delta}_n := \frac{\widehat{U}_n - U_n}{\sigma_n} = \check{\Delta}_n + \check{R}_n$$

1699 where  $\check{\Delta}_n$  and  $\check{R}_n$  satisfy

1700 (8.105) 
$$\check{R}_n = o_p(\widetilde{\mathcal{M}}(\rho_n, n; R))$$

1701 and the original (3.12) in Lemma 3.1-(b) holds for  $\check{\Delta}_n$ , where the definition and asymptotic order of  $\sigma_w$  is identical to that in Lemma 3.1,

- 1703 (c)  $\widehat{\delta}_n = o_p(\widetilde{\mathcal{M}}(\rho_n, n; R))$ ,  
 1704 (d) *Identical to Lemma 3.1-(d).*

1705 PROOF OF LEMMA 8.4. The proof of this lemma can be obtained by slightly varying the  
 1706 proof of Lemma 3.1.

- 1707 (a) (No additional proof needed.)  
 1708 (b) The only change we need to make to the proof of Lemma 3.1-(b) to make it a valid proof  
 1709 here is to replace (8.18) by the following concentration inequality:

(8.106)

$$\begin{aligned} 1710 & \mathbb{P} \left( \check{R}_n := \frac{\check{R}}{\binom{n}{r} \cdot \sigma_n} \geq C \cdot \widetilde{\mathcal{M}}(\rho_n, n; R) \right) \\ 1711 & \leq \begin{cases} \max \left\{ \exp \left( -\frac{((\rho_n \cdot n)^{-1} \cdot (\rho_n \cdot n)^{1/2})^2}{(\rho_n \cdot n)^{-2}} \right), \exp \left( -\frac{(\rho_n \cdot n)^{-1} \cdot (\rho_n \cdot n)^{1/2}}{(\rho_n \cdot n)^{-1} \cdot n^{-1/2}} \right) \right\}, & \text{for acyclic } R; \\ \max \left\{ \exp \left( -\frac{((\rho_n^{-r/2} \cdot n^{-1}) \cdot (\rho_n^{r/2} \cdot n)^{1/2})^2}{(\rho_n^{-r/2} \cdot n^{-1})^2} \right), \exp \left( -\frac{(\rho_n^{-r/2} \cdot n^{-1}) \cdot (\rho_n^{r/2} \cdot n)^{1/2}}{\rho_n^{-3} n^{-5/2}} \right) \right\}, & \text{for cyclic } R; \end{cases} \\ 1712 & = o(1) \end{aligned}$$

1713 The proof of this part is completed.

- (c) We only need to change how we use Theorem 3 of Schudy and Sviridenko [116] in (8.42), into the following way

$$\sum_{\substack{\text{All possible } (v,p): \\ v \geq 2, p \geq 3}} \widehat{\Delta}^{(i;v,p)} = o_p(\rho_n^s \cdot n^{r-1} \cdot \widetilde{\mathcal{M}}(\rho_n, n; R))$$

1714 and for the rest of the proof of Lemma 3.1-(c), replace every remainder term in the format  
 1715 of “ $\widetilde{O}_p(\cdots \times \mathcal{M}(\rho_n, n; R))$ ” by “ $o_p(\cdots \times \widetilde{\mathcal{M}}(\rho_n, n; R))$ ”. This completes the proof.

- 1716 (d) (No additional proof needed.)

1717 □

Now we return to the proof of the second part of Theorem 3.4. The proof is completed by slightly varying (8.97) in the proof of the first part of this theorem by

$$\widehat{T}_n = U_n^\# + \widehat{\Delta}_n(1 + o_p(\widetilde{\mathcal{M}}(\rho_n, n; R))) + o_p(\widetilde{\mathcal{M}}(\rho_n, n; R))$$

1718 Then recall the definition of “ $o_p$ ” and apply Lemma 2 of Maesono [96] (setting  $T = \widehat{T}_n$ ,  
 1719  $\widetilde{T} = U_n^\#$  and  $\alpha = \widetilde{\mathcal{M}}(\rho_n, n; R)$  and  $H(x) = \Phi(x)$ ). We have

$$\begin{aligned} 1720 & \|F_{\widehat{T}_n}(u) - \Phi(u)\|_\infty \leq \|F_{U_n^\#}(u) - \Phi(u)\|_\infty + \mathbb{P} \left\{ |\widehat{T}_n - U_n^\#| \geq \widetilde{\mathcal{M}}(\rho_n, n; R) \right\} \\ 1721 & + O(\widetilde{\mathcal{M}}(\rho_n, n; R)) = o(1) + o(1) + o(1) \rightarrow 0 \end{aligned}$$

1722 This completes the proof of the second part and thus the proof of the error bound of the  
1723 population version Edgeworth expansion in Theorem 3.4.

1724 Next, we prove the error bound for the empirical version Edgeworth expansion in Theorem  
1725 3.4. Similar to the proof of the population version, we discuss two cases.

1726 Part I: the proof of the  $\tilde{O}_p(\mathcal{M}(\rho_n, n; R))$  bound when  $\rho_n n > \log^{1/2} n$  (acyclic) or  $\rho_n^{r/2} n >$   
1727  $\log^{1/2} n$  (cyclic) is easily done by citing the following intermediate results from the proof of  
1728 Theorem 3.2.

$$1729 \quad (8.107) \quad \hat{\xi}_1^3 - \xi_1^3 = \tilde{O}_p(\rho_n^{3s} \cdot n^{-1/2} \log^{1/2} n),$$

$$1730 \quad (8.108) \quad |\hat{\mathbb{E}}[g_1^3(X_1)] - \mathbb{E}[g_1^3(X_1)]| = \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n),$$

$$1731 \quad \left| \binom{n}{2}^{-1} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_1) \hat{g}_1(X_2) \hat{g}_2(X_1, X_2) - \mathbb{E}[g_1(X_1) g_1(X_2) g_2(X_1, X_2)] \right|$$

$$1732 \quad (8.109) \quad = \tilde{O}_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \log^{1/2} n).$$

1733 The proof of this part is then instantly done by combining these results with the statement  
1734 about the population Edgeworth expansion in the sparse case that we just proved above.

1735 Part II: proof of the  $o_p(1)$  bound when  $1 < \rho_n n \leq \log^{1/2} n$  (acyclic) or  $1 < \rho_n^{r/2} n \leq$   
1736  $\log^{1/2} n$ . To prove for this regime, we only need to slightly vary the proof of Theorem 3.2.  
1737 Set a series  $\varrho_n \rightarrow \infty$  as follows:

$$1738 \quad (8.110) \quad \varrho_n := \begin{cases} \rho_n \cdot n, & \text{for acyclic } R, \\ \rho_n^{r/2} \cdot n, & \text{for cyclic } R. \end{cases}$$

1738 By replacing the  $\log^{1/2} n$  factor in all the “ $u$ ” values that we set in Theorem 8.1 by  $\varrho_n$ , where  
1739 we apply it in the proof of Lemma 3.1-(c),(d) and in the proof of Theorem 3.2, we establish  
1740 the following analogous intermediate results:

$$1741 \quad (8.111) \quad \hat{\xi}_1^3 - \xi_1^3 = O_p(\rho_n^{3s} \cdot n^{-1/2} \cdot \varrho_n) = o_p(\rho_n^{3s})$$

$$1742 \quad (8.112) \quad |\hat{\mathbb{E}}[g_1^3(X_1)] - \mathbb{E}[g_1^3(X_1)]| = O_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \cdot \varrho_n) = o_p(\rho_n^{3s-1/2})$$

$$1743 \quad \left| \binom{n}{2}^{-1} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_1) \hat{g}_1(X_2) \hat{g}_2(X_1, X_2) - \mathbb{E}[g_1(X_1) g_1(X_2) g_2(X_1, X_2)] \right|$$

$$1744 \quad (8.113) \quad = O_p(\rho_n^{3s-1/2} \cdot n^{-1/2} \cdot \varrho_n) = o_p(\rho_n^{3s-1/2})$$

1745 This implies  $\|G_n(u) - \hat{G}_n(u)\|_\infty = o_p((\rho_n n)^{-1/2}) = o_p(1)$  (thus immediately completes  
1746 the proof of Part II) by simply reproducing the rest of the proof of Theorem 3.2. The proof  
1747 of the entire Theorem 3.4 is now complete. □

1748

8.6. *Proof of Theorem 4.1.* We will mainly prove for the node sub-sampling network  
bootstrap scheme [17], and the corresponding conclusion for the re-sampling scheme can  
be obtained easily by slightly varying the proof for sub-sampling. Conditioned on  $A$ , since  
the sub-sampling objects in network models are nodes rather than latent variables  $X_j$ 's<sup>3</sup>,

<sup>3</sup>In other words,  $X_j$ 's in the bootstrap procedure are deemed fixed.

we change the notation for simplicity. Define  $\mathcal{V}_* = \{1 \leq v_1 < v_2 < \dots < v_{n^*} \leq n\}$  to be uniformly sampled from all size- $n^*$  subsets of  $[n]$ . That is,

$$\mathbb{P}(\mathcal{V}_* = \{i_1, \dots, i_{n^*}\}) = \frac{1}{\binom{n}{n^*}} \quad \forall 1 \leq i_1 < \dots < i_{n^*} \leq n.$$

Define the bootstrap expectation  $\mathbb{E}^*$  to be taken with respect to the randomness of  $\mathcal{V}_*$ . The sub-sampling bootstrap sample network moment  $\hat{U}_{n^*}^b$  calculated from the sub-network  $A_{\mathcal{V}_*, \mathcal{V}_*}$  calculated according to [17] is

$$\hat{U}_{n^*}^b = \frac{1}{\binom{n^*}{r}} \sum_{i_1 < \dots < i_r \subset \mathcal{V}_*} h(A_{i_1, i_2, \dots, i_r}).$$

1749 To emphasize that the randomness in this bootstrap setting is solely due to  $\mathcal{V}_*$  and simplify  
1750 notation, we define  $\hat{g}_1^b(v_1)$ , taking the argument  $v_1$  rather than  $X_{v_1}$ , as follows

(8.114)

$$1751 \quad \hat{g}_1^b(v_1) := \frac{n-1}{n-n^*} \left\{ \frac{1}{\binom{n^*-1}{r-1}} \mathbb{E}^* \left[ \sum_{i_1, \dots, i_{r-1} \subset \mathcal{V}_* \setminus v_1} h(A_{v_1, i_1, \dots, i_{r-1}}) \middle| v_1 \right] - \hat{U}_n \right\}$$

$$1752 \quad \hat{g}_2^b(v_1, v_2) := \frac{n-3}{n-n^*-1} \left( \frac{n-2}{n-n^*} \left\{ \mathbb{E}^* \left[ \frac{1}{\binom{n^*-2}{r-2}} \sum_{i_1, \dots, i_{r-2} \subset \mathcal{V}_* \setminus \{v_1, v_2\}} h(A_{v_1, v_2, i_1, \dots, i_{r-2}}) \middle| v_1, v_2 \right] \right. \right.$$

(8.115)

$$1753 \quad \left. - \hat{U}_n \right\} - \hat{g}_1^b(v_1) - \hat{g}_1^b(v_2) \Big)$$

1754 where the finite population correction term  $(n-1)/(n-n^*)$  comes from [23, (1.2)]. where  
1755 again the finite population correction term  $(n-3)/(n-n^*-1)$  is due to [23, (1.3)]. Recall  
1756 that  $\hat{S}_{n^*}^*$  is a jackknife estimator of  $\text{Var}^*(\hat{U}_{n^*}^b | A)$  and that the bootstrap test statistic as

$$(8.116) \quad \hat{T}_{n^*}^* = \frac{\hat{U}_{n^*}^b - \hat{U}_n}{\hat{S}_{n^*}^*}$$

1757 By our proof of Theorem 3.3, the difference between a jackknife estimator and an estimator  
1758 based on  $\xi_1^*$  is ignorable, and we are free to choose either. Here we use the jackknife estimator  
1759 in order to better connect with Bloznelis [23]. To start, we check that  $\mathbb{E}^*[\hat{U}_{n^*}^b] = \hat{U}_n$  where  
1760 the expectation is taken with respect to the randomness of  $\mathcal{V}_*$ , so that (8.116) is an valid  
1761 studentization of the U-statistic. To see this, notice that

$$1762 \quad \mathbb{E}^*[\hat{U}_{n^*}^b] = \frac{1}{\binom{n}{n^*}} \sum_{\mathcal{V}_* \subset [n]} \hat{U}_{n^*}^b = \frac{1}{\binom{n}{n^*}} \frac{1}{\binom{n^*}{r}} \sum_{\mathcal{V}_* \subset [n]} \sum_{i_1 < \dots < i_r \subset \mathcal{V}_*} h(A_{i_1, i_2, \dots, i_r}).$$

1763 On the RHS, each summand  $h(A_{i_1, \dots, i_r})$  appears  $\binom{n-r}{n^*-r}$  times. Therefore,

$$1764 \quad \sum_{\mathcal{V}_* \subset [n]} \sum_{i_1 < \dots < i_r \subset \mathcal{V}_*} h(A_{i_1, i_2, \dots, i_r}) = \binom{n-r}{n^*-r} \sum_{1 \leq i_1 < \dots < i_r \leq n} h(A_{i_1, \dots, i_r})$$

$$1765 \quad = \binom{n-r}{n^*-r} \binom{n}{r} \hat{U}_n.$$

1766 As a result,

$$\begin{aligned}
 1767 \quad \mathbb{E}^*[\widehat{U}_{n^*}^b] &= \frac{1}{\binom{n}{n^*}} \frac{1}{\binom{n^*}{r}} \sum_{\mathcal{V}_* \subset [n]} \sum_{i_1 < \dots < i_r \in \mathcal{V}_*} h(A_{i_1, i_2, \dots, i_r}) \\
 1768 \quad &= \frac{1}{\binom{n}{n^*}} \frac{1}{\binom{n^*}{r}} \binom{n-r}{n^*-r} \binom{n}{r} \cdot \widehat{U}_n = \widehat{U}_n
 \end{aligned}$$

1769 To investigate the distribution of  $\widehat{T}_{n^*}^*$  under the finite-population sampling obeying  $\mathcal{V}_*$ , we  
 1770 define the bootstrap Edgeworth expansion by

$$\begin{aligned}
 1771 \quad G_{n^*}^*(x) &:= \Phi(x) + \frac{\varphi(x)}{\sqrt{n^*(1-n^*/n)} \cdot (\xi_1^*)^3} \cdot \left\{ \frac{2x^2+1}{6} \cdot \mathbb{E}^* \left\{ \widehat{g}_1^b(v_1) \right\}^3 \right. \\
 1772 \quad (8.117) \quad &\left. + \frac{r-1}{2} \cdot (x^2+1) \mathbb{E}^* [\widehat{g}_1^b(v_1) \widehat{g}_1^b(v_2) \widehat{g}_2^b(v_1, v_2)] \right\}
 \end{aligned}$$

1773 where recall the definitions of  $\widehat{g}_1^b(\cdot), \widehat{g}_2^b(\cdot, \cdot)$  from (8.114) and (8.115), respectively. Here,  
 1774  $(\xi_1^*)^2 := \text{Var}^*(\widehat{g}_1^b(v_1)|A) = \mathbb{E}^*[(\widehat{g}_1^b(v_1))^2]$ .

1775 Next, we are going to apply Theorem 1 of [23]. The Cramer's condition (1.11) in Theorem  
 1776 1 in [23] is different from the conventional version, and we need to check that it indeed  
 1777 holds in our setting. Specifically, in our setting, it suffices to prove that there exists a positive  
 1778 sequence  $\{t_n\} \rightarrow \infty$  and a universal constant  $M_1 : 0 < M_1 < 1$ , such that

$$\mathbb{P} \left( \sup_{t \in (0, t_n)} \left| n^{-1} \sum_{j=1}^n e^{it_n \widehat{g}_1(X_j)/\widehat{\xi}_1} \right| \leq M_1 < 1 \right) \xrightarrow{p} 1$$

1779 because our eventual bounds are  $O_p$  bounds, and in the proof we can choose to discuss only  
 1780 events that will happen with high probability. Recall from the proof of Theorem 3.2 that we  
 1781 have shown the following facts

$$\begin{aligned}
 1782 \quad |\widehat{g}_1(X_i) - g_1(X_i)| &= \widetilde{O}_p(\rho_n^{s-1/2} \cdot n^{-1/2} \log^{1/2} n) \\
 1783 \quad |\widehat{\xi}_1 - \xi_1| &= \widetilde{O}_p(\rho_n^s \cdot n^{-1/2} \log^{1/2} n)
 \end{aligned}$$

and the simple fact that  $\xi_1 = \rho_n^s$  a.s. Therefore, we have

$$|\widehat{g}_1(X_j)/\widehat{\xi}_1 - g_1(X_j)/\xi_1| = \widetilde{O}_p(\rho_n^{-1/2} \cdot n^{-1/2} \log^{1/2} n)$$

1784 Recall that our assumption implies  $\rho_n n \rightarrow \infty$  throughout this paper (regardless of  $R$  shapes,  
 1785 all assumptions we made imply this). Choosing  $t_n = (\rho_n \cdot n)^{1/4}$ , we have

$$\begin{aligned}
 1786 \quad &\sup_{t \in (0, t_n)} \left| \frac{1}{n} \sum_{j=1}^n e^{itg_1(X_j)/\xi_1} - \frac{1}{n} \sum_{j=1}^n e^{it\widehat{g}_1(X_j)/\widehat{\xi}_1} \right| \\
 1787 \quad &\leq \sup_{t \in (0, t_n)} t \cdot \max_{1 \leq j \leq n} \left| g_1(X_j)/\xi_1 - \widehat{g}_1(X_j)/\widehat{\xi}_1 \right| \cdot e^{t \cdot |g_1(X_j)/\xi_1 - \widehat{g}_1(X_j)/\widehat{\xi}_1|} \\
 1788 \quad &(\text{w.p. } 1 - Cn^{-1}) \leq \sup_{t \in (0, t_n)} t(\rho_n \cdot n)^{-1/2} \log^{1/2} n \cdot e^{t(\rho_n \cdot n)^{-1/2} \log^{1/2} n} \leq t_n(\rho_n \cdot n)^{-1/2}
 \end{aligned}$$

1789 under the specified sparsity conditions.

1790 It suffices to bound  $\sup_{t \in (0, t_n)} \left| n^{-1} \sum_{j=1}^n e^{itg_1(X_j)/\xi_1} \right|$ . For every given  $t \in \mathcal{T}_n := \{k/n : k \in$   
 1791  $\mathbb{N}, k/n \leq t_n\}$ , by Bernstein's inequality, we have

$$1792 \quad \mathbb{P} \left( \left| n^{-1} \sum_{j=1}^n e^{itg_1(X_j)/\xi_1} - \mathbb{E} \left[ e^{itg_1(X_1)/\xi_1} \right] \right| > \epsilon \right) \leq 2e^{-Cn\epsilon^2}$$

1793 Therefore, setting  $M_1 := \limsup_{t \rightarrow \infty} |\mathbb{E} [e^{itg_1(X_1)/\xi_1}]|$ , by the Cramer's condition we as-  
 1794 sumed in Theorem 4.1, we have  $M_1 \in (0, 1)$  and  $(1 + M_1)/2 \in (0, 1)$ . Therefore

$$1795 \quad \mathbb{P} \left( \sup_{t \in \mathcal{T}_n} \left| n^{-1} \sum_{j=1}^n e^{itg_1(X_j)/\xi_1} \right| > (1 + M_1)/2 \right) \leq |\mathcal{T}_n| \cdot 2e^{-C_3n(M_1/2)^2} \leq e^{-C_4n}$$

1796 for some universal constants  $C_3, C_4 > 0$ . Now noticing that for any  $t \in (0, t_n)$ , let  $t'$  be the  
 1797 best approximation to  $t$  in  $\mathcal{T}_n$ , we have

$$1798 \quad \sup_{t \in (0, t)} \left| \frac{1}{n} \sum_{j=1}^n e^{itg_1(X_j)/\xi_1} - \frac{1}{n} \sum_{j=1}^n e^{it'g_1(X_j)/\xi_1} \right|$$

$$1799 \quad (\text{w.h.p.}) \leq |t - t'|(\rho_n \cdot n)^{-1/2} \cdot e^{|t-t'|(\rho_n \cdot n)^{-1/2}} \leq t_n \cdot (\rho_n \cdot n)^{-1/2} \rightarrow 0$$

1800 The verification that our ordinary Cramer's condition implies the sample version in [23] is  
 1801 thus finished.

1802 By Theorem 1 of [23], the sampling distribution of  $\hat{T}_{n^*}^*$  by node sub-sampling enjoys the  
 1803 following uniform bound

$$(8.118) \quad \left\| F_{\hat{T}_{n^*}^*}(u) - G_{n^*}^*(u) \right\|_{\infty} = o_p((n^*)^{-1/2})$$

1804 It then suffices to establish the connection between  $G_{n^*}^*(u)$  and  $\hat{G}_{n^*(1-n^*/n)}(u)$ . The  
 1805 proof strategy is to show that (8.117) can be transcribed, with  $\mathbb{E}^*$  replaced by  $\hat{\mathbb{E}}$ 's and  
 1806  $\hat{g}_1^b(v_1), \hat{g}_2^b(v_1, v_2)$  replaced with  $\hat{g}_1(X_1), \hat{g}_2(X_1, X_2)$ , respectively. Then the comparison of  
 1807 the Edgeworth coefficients in  $G_{n^*}^*(u)$  and  $\hat{G}_{n^*(1-n^*/n)}(u)$  would complete the proof. To  
 1808 proceed, now we focus on analyzing the core quantities  $\hat{g}_1^b(v_1)$  and  $\hat{g}_2^b(v_1, v_2)$ . For  $\hat{g}_1^b(v_1)$ ,  
 1809 since conditioning on  $v_1 \in \mathcal{V}_*$ , the rest indexes  $\{v_2, \dots, v_{n^*}\}$  are uniformly sampled from  
 1810  $\{\{i_1, \dots, i_{n^*-1}\} \subset [n] \setminus v_1\}$ , we have

$$1811 \quad \frac{1}{\binom{n^*-1}{r-1}} \cdot \mathbb{E}^* \left[ \sum_{i_1, \dots, i_{r-1} \subset \mathcal{V}_* \setminus v_1} h(A_{v_1, i_1, \dots, i_{r-1}}) \middle| v_1 \right]$$

$$1812 \quad = \frac{1}{\binom{n^*-1}{r-1}} \frac{1}{\binom{n-1}{n^*-1}} \sum_{\mathcal{V}_* \subset [n]: v_1 \in \mathcal{V}_*} \sum_{i_1, \dots, i_{r-1} \in \mathcal{V}_* \setminus v_1} h(A_{v_1, i_1, \dots, i_{r-1}})$$

$$1813 \quad (\text{By (8.26)}) = \frac{1}{\binom{n^*-1}{r-1}} \frac{1}{\binom{n-1}{n^*-1}} \binom{n-r}{n^*-r} \binom{n-1}{r-1} \cdot \hat{a}_{v_1} = \hat{a}_{v_1}.$$

1814 where in the second equality, we noticed that each  $h(A_{v_1, i_1, \dots, i_{r-1}})$  appears  $\binom{n-r}{n^*-r}$  times in  
 1815 the first line. Therefore,

$$(8.119) \quad \hat{g}_1^b(v_1) = \frac{n-1}{n-n^*} [\hat{a}_{v_1} - \hat{U}_n] = \frac{n-1}{n-n^*} \cdot \hat{g}_1(X_{v_1})$$

1816 where  $\hat{g}_1(X_{v_1})$  appeared (in “ $\hat{\mathbb{E}}$ ” terms) in Theorem 3.2. Then we have

$$1817 \quad \mathbb{E}^* \left[ \{\hat{g}_1^b(v_1)\}^3 \right] = \frac{1}{n} \sum_{i=1}^n \left( \frac{n-1}{n-n^*} \right)^3 (\hat{a}_i - \hat{U}_n)^3 = \left( \frac{n-1}{n-n^*} \right)^3 \hat{\mathbb{E}}[g_1^3(X_1)]$$

$$1818 \quad (\xi_1^*)^2 = \text{Var}^*(\hat{g}_1^b(v_1)|A) = \mathbb{E}^*[(\hat{g}_1^b(v_1))^2] = \frac{(n-1)^2}{(n-n^*)^2} \cdot \hat{\xi}_1^2$$

1819 where the definitions of  $\hat{\xi}_1$  and  $\hat{\mathbb{E}}g_1^3(X_1)$  can also be recalled by reviewing Theorem 3.2.  
 1820 Now we turn to analyzing  $\mathbb{E}^*[\{\hat{g}_1^b(v_1)\hat{g}_1^b(v_2)\hat{g}_2^b(v_1, v_2)\}]$ . The main part of the definition of  
 1821  $\hat{g}_2^b(v_1, v_2)$  can be re-expressed as follows

$$1822 \quad \mathbb{E}^* \left[ \frac{1}{\binom{n^*-2}{r-2}} \sum_{i_1, \dots, i_{r-2} \subset \mathcal{V}_* \setminus \{v_1, v_2\}} h(A_{v_1, v_2, i_1, \dots, i_{r-2}}) \middle| v_1, v_2 \right]$$

$$1823 \quad = \frac{1}{\binom{n-2}{n^*-2}} \frac{1}{\binom{n^*-2}{r-2}} \sum_{\mathcal{V}_* \subset [n]: v_1, v_2 \in \mathcal{V}_*} \sum_{i_1, \dots, i_{r-2} \subset \mathcal{V}_* \setminus \{v_1, v_2\}} h(A_{v_1, v_2, i_1, \dots, i_{r-2}})$$

$$1824 \quad = \frac{1}{\binom{n-2}{n^*-2}} \frac{1}{\binom{n^*-2}{r-2}} \binom{n-2}{r-2} \binom{n-r}{n^*-r} \hat{a}_{v_1 v_2} = \hat{a}_{v_1 v_2}$$

1825 where we recall the definition of  $\hat{a}_{ij}$  from (8.82). Combining this with (8.119), we have

$$1826 \quad \hat{g}_2^b(v_1, v_2) = \frac{n-3}{(n-n^*-1)} \left[ \frac{n-2}{n-n^*} (\hat{a}_{v_1 v_2} - \hat{U}_n) - \frac{n-1}{n-n^*} (\hat{a}_{v_1} - \hat{U}_n) - \frac{n-1}{n-n^*} (\hat{a}_{v_2} - \hat{U}_n) \right]$$

$$1827 \quad = \frac{(n-3)(n-1)}{(n-n^*-1)(n-n^*)} \left[ (\hat{a}_{v_1 v_2} - \hat{U}_n) - (\hat{a}_{v_1} - \hat{U}_n) - (\hat{a}_{v_2} - \hat{U}_n) \right]$$

$$1828 \quad - \frac{(n-3)}{(n-n^*-1)(n-n^*)} (\hat{a}_{v_1 v_2} - \hat{U}_n).$$

1829 Then we have

$$1830 \quad \mathbb{E}^*[\hat{g}_1^b(v_1)\hat{g}_1^b(v_2)\hat{g}_2^b(v_1, v_2)] = \frac{1}{\binom{n}{2}} \sum_{1 \leq v_1 < v_2 \leq n} \hat{g}_1^b(v_1)\hat{g}_1^b(v_2)\hat{g}_2^b(v_1, v_2)$$

$$1831 \quad = \frac{(n-3)(n-1)^3}{(n-n^*-1)(n-n^*)^3} \hat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)]$$

$$1832 \quad - \frac{(n-3)(n-1)^2}{(n-n^*-1)(n-n^*)^3} \cdot \frac{1}{\binom{n}{2}} \sum_{1 \leq i < j \leq n} \hat{g}_1(X_1)\hat{g}_1(X_2)[\hat{g}_2(X_1, X_2) + \hat{g}_1(X_1) + \hat{g}_1(X_2)]$$

$$1833 \quad = \frac{(n-3)(n-1)^3}{(n-n^*-1)(n-n^*)^3} \hat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] + \tilde{O}_p \left( \frac{(n-3)(n-1)^2}{(n-n^*-1)(n-n^*)^3} \cdot \rho_n^{3s-1} \log^{1/2} n \right)$$

1834 where in the last line, we used that,  $\hat{g}_1(X_1) \stackrel{p}{\approx} \rho_n^s, \hat{g}_2(X_1, X_2) \stackrel{p}{\approx} \rho_n^{s-1}$  with probability at  
 1835 least  $1 - O(n^{-1})$  by the proof of Theorem 3.2. Define  $\alpha_{n^*} = (n-1)/(n-n^*)$ . Now we can  
 1836 rewrite (8.117) as follows

$$1837 \quad G_{n^*}^*(x) = \Phi(x) + \frac{\varphi(x)}{\sqrt{n^*(1-n^*/n)} \cdot \alpha_{n^*}^3 \hat{\xi}_1^3} \left\{ \frac{2x^2+1}{6} \cdot \alpha_{n^*}^3 \hat{\mathbb{E}}[g_1^3(X_1)] \right.$$

$$1838 \quad \left. + \frac{r-1}{2} \cdot \alpha_{n^*}^3 (x^2+1) \hat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\}$$

$$\begin{aligned}
 & - \frac{n-3}{(n-n^*-1)(n-n^*)} \frac{r-1}{2} \cdot \alpha_{n^*}^2 (x^2+1) \cdot \tilde{O}_p(\rho_n^{3s-1} \log^{1/2} n) \Big\} \\
 & = \hat{G}_{n^*(1-n^*/n)}(u) + \tilde{O}_p \left\{ \frac{\log^{1/2} n}{\sqrt{n^*(1-n^*/n)(n-n^*)\rho_n}} \right\}
 \end{aligned}$$

where recall that  $\hat{G}_n(u)$  was defined Theorem 3.2. Finally, we have

$$\left\| G_{n^*}^*(u) - \hat{G}_{n^*(1-n^*/n)}(u) \right\|_\infty = \tilde{O}_p \left\{ \frac{\log^{1/2}(n)}{\sqrt{n^*(1-n^*/n)(n-n^*)\rho_n}} \right\}$$

where the last equation is due to  $\rho_n = \omega(n^{-1/r})$  and  $n - n^* \asymp n$ . Combining this with Theorem 3.1 and Theorem 3.2, by a triangular inequality, we have

$$(8.120) \quad \left\| F_{\hat{T}_n^*}(u) - F_{\hat{T}_{n^*(1-n^*/n)}}(u) \right\|_\infty = o_p((n^*)^{-1/2}).$$

This completes the proof of Theorem 4.1 for sub-sampling, since the uniform convergence rate of the Edgeworth expansion is governed by the worst convergence rate of its coefficient terms.

Now we discuss the re-sampling scheme. Sampling  $\{v_1, \dots, v_{n^*}\}$  with replacement from a finite population  $[n]$  is equivalent to sampling without replacement from a population in which each of  $[n]$  are repeated infinite many times with the same infinite cardinality such that a uniform sampler will still take each of  $[n]$  with equal probabilities. This amounts to set the “ $n$ ” in Bloznelis [23] to “ $n = \infty$ ”<sup>4</sup>. Notice, however, the “ $n$ ” in [23] should not be confused with our network size  $n$  in the expressions of  $\xi_1^*$ ,  $\mathbb{E}^*[\{\hat{g}_1^b(v_1)\}^3]$  and  $\mathbb{E}^*[\hat{g}_1^b(v_1)\hat{g}_1^b(v_2)\hat{g}_2^b(v_1, v_2)]$ . Therefore, the re-sampling bootstrap Edgeworth expansion is the following slight-modification of (8.117):

$$\begin{aligned}
 G_{n^*}^*(x) & := \Phi(x) + \frac{\varphi(x)}{\sqrt{n^*} \cdot (\xi_1^*)^3} \cdot \left\{ \frac{2x^2+1}{6} \cdot \mathbb{E}^* \left\{ \hat{g}_1^b(v_1) \right\}^3 \right. \\
 & \quad \left. + \frac{r-1}{2} \cdot (x^2+1) \mathbb{E}^* [\hat{g}_1^b(v_1)\hat{g}_1^b(v_2)\hat{g}_2^b(v_1, v_2)] \right\}
 \end{aligned}$$

The rest of the proof is exactly similar to that for sub-sampling and thus will be omitted. The proof of Theorem 4.1 is completed.

**PROOF OF THEOREM 4.2.** The key to this proof is to establish the local monotonicity of the function  $G_n(\cdot)$ . The local curvature of  $G_n$  is handier to use than that of  $F_{\hat{T}_n}$ , because the distribution of  $\hat{T}_n$  may not be exactly continuous, and the classical result  $F_Z(Z) \sim \text{Uniform}[0, 1]$  (thus  $\mathbb{P}(F_Z(Z) \leq u) = u$  for any  $u \in [0, 1]$ ) for a continuous random variable  $Z$  does not necessarily apply. On the other hand, by construction,  $G_n$  is always smooth.

Now notice that not only  $G_n(\cdot)$  uniformly converges to the  $N(0, 1)$  CDF  $\Phi(\cdot)$ , but further, these two functions are both smooth and  $\sup_u |G_n'(u) - \Phi'(u)| \rightarrow 0$  (while the CDF  $F_{\hat{T}_n}(\cdot)$  is not necessarily continuous). Therefore, there exists a large enough constant  $n_0$  and small constants  $\epsilon_0 > 0, \delta_0 > 0$ , such that the following two properties hold simultaneously

<sup>4</sup>Here, we clarify that the “ $n$ ” in “ $n = \infty$ ” should be understood as the size of the finite population for bootstrapping, among the notation system of [23], not the “ $n$ ” in most of this paper as the network size.

- 1870 (i). For all  $n \geq n_0$ , we have  $G_n(u) \geq \alpha/2 + \epsilon_0$  for all  $u \geq z_{\alpha/2} + \delta_0$ ; and  $G_n(u) \leq \alpha/2 - \epsilon_0$   
 1871 for all  $u \leq z_{\alpha/2} - \delta_0$   
 1872 (ii). For all  $n \geq n_0$ , on the interval  $u \in [z_{\alpha/2} - \delta_0, z_{\alpha/2} + \delta_0]$ , we have  $0 < C_{\alpha/2} \leq G'_n(u) \leq$   
 1873  $D_{\alpha/2}$  and constants  $C_{\alpha/2}, D_{\alpha/2}$  only depend on  $\alpha$ .

1874 Properties (ii) implies that  $G_n$  is strictly monotone and thus invertible in  $[z_{\alpha/2} - \delta_0, z_{\alpha/2} + \delta_0]$ ,  
 1875 and specifically,  $G_n^{-1}(\alpha/2)$  is well-defined. Then by (i), we have  $G_n^{-1}(\alpha/2) \in [z_{\alpha/2} -$   
 1876  $\delta_0, z_{\alpha/2} + \delta_0]$ , and by (ii), we know that  $G_n^{-1}(u')$  is also Lipschitz on  $u' \in [G_n(z_{\alpha/2} -$   
 1877  $\delta_0), G_n(z_{\alpha/2} + \delta_0)]$ .

1878 Now we are ready to begin the main proof for Type-I error rate. We have

$$\begin{aligned}
 & \text{Type-I error rate} := \mathbb{P}_{H_0} \left( 2 \cdot \min \left\{ \widehat{G}_n(\widehat{T}_n), 1 - \widehat{G}_n(\widehat{T}_n) \right\} < \alpha \right) \\
 & \quad (\alpha \text{ is small}) = \mathbb{E}_{H_0} \left[ \mathbb{1}_{[\widehat{G}_n(\widehat{T}_n) \leq \alpha/2]} + \mathbb{1}_{[\widehat{G}_n(\widehat{T}_n) > 1 - \alpha/2]} \right] \\
 & \quad (\text{Theorems 3.1 + 3.2}) = \mathbb{E}_{H_0} \left[ \mathbb{1}_{[G_n(\widehat{T}_n) \leq \alpha/2]} + \mathbb{1}_{[G_n(\widehat{T}_n) > 1 - \alpha/2]} \right] \\
 & \quad \quad \quad + \mathbb{E}_{H_0} \left[ \mathbb{1}_{[\widehat{G}_n(\widehat{T}_n) \leq \alpha/2]} - \mathbb{1}_{[G_n(\widehat{T}_n) \leq \alpha/2]} \right] \\
 & \quad \quad \quad + \mathbb{E}_{H_0} \left[ \mathbb{1}_{[\widehat{G}_n(\widehat{T}_n) > 1 - \alpha/2]} - \mathbb{1}_{[G_n(\widehat{T}_n) > 1 - \alpha/2]} \right] \\
 & \quad (8.122) \quad = \mathbb{P}_{H_0} \left( G_n(\widehat{T}_n) \leq \alpha/2 \right) + \mathbb{P}_{H_0} \left( G_n(\widehat{T}_n) \geq 1 - \alpha/2 \right) \\
 & \quad \quad \quad + O(\mathcal{M}(\rho_n, n; R)),
 \end{aligned}$$

1886 where the last equality is due to (recall from the proof of Theorem 3.2 that  $\|\widehat{G}_n(x) -$   
 1887  $G_n(x)\|_\infty = \widetilde{O}_p(\rho_n^{-1}n^{-1})$ )

$$\begin{aligned}
 & \mathbb{E}_{H_0} \left| \mathbb{1}_{[\widehat{G}_n(\widehat{T}_n) \leq \alpha/2]} - \mathbb{1}_{[G_n(\widehat{T}_n) \leq \alpha/2]} \right| \\
 & \quad = \mathbb{P}_{H_0} \left( \widehat{G}_n(\widehat{T}_n) \leq \alpha/2, G_n(\widehat{T}_n) > \alpha/2 \right) + \mathbb{P}_{H_0} \left( \widehat{G}_n(\widehat{T}_n) > \alpha/2, G_n(\widehat{T}_n) \leq \alpha/2 \right) \\
 & \quad = \mathbb{P}_{H_0} \left( G_n(\widehat{T}_n) \leq \alpha/2 + O(\rho_n^{-1}n^{-1}), G_n(\widehat{T}_n) > \alpha/2 \right) \\
 & \quad \quad \quad + \mathbb{P}_{H_0} \left( G_n(\widehat{T}_n) > \alpha/2 - O(\rho_n^{-1}n^{-1}), G_n(\widehat{T}_n) \leq \alpha/2 \right) + O(n^{-1}) \\
 & \quad (\text{Invertibility of } G_n(\cdot)) = \mathbb{P}_{H_0} \left( G_n^{-1}(\alpha/2 - O(\rho_n^{-1}n^{-1})) \leq \widehat{T}_n \leq G_n^{-1}(\alpha/2 + O(\rho_n^{-1}n^{-1})) \right) \\
 & \quad \quad \quad + O(n^{-1}) \\
 & \quad (\text{Theorem 3.1}) = G_n(G_n^{-1}(\alpha/2 + O(\rho_n^{-1}n^{-1}))) - G_n(G_n^{-1}(\alpha/2 - O(\rho_n^{-1}n^{-1}))) \\
 & \quad \quad \quad + O(\mathcal{M}(\rho_n, n; R)) = O(\mathcal{M}(\rho_n, n; R)).
 \end{aligned}$$

1896 Now we continue (8.122) and bound  $\mathbb{P} \left( G_n(\widehat{T}_n) \leq \alpha/2 \right)$ . We have

$$\begin{aligned}
 & \mathbb{P} \left( G_n(\widehat{T}_n) \leq \alpha/2 \right) = \mathbb{P} \left( G_n(\widehat{T}_n) \leq \alpha/2, \widehat{T}_n \in [z_{\alpha/2} - \delta_0, z_{\alpha/2} + \delta_0] \right) \\
 & \quad \quad \quad + \mathbb{P} \left( G_n(\widehat{T}_n) \leq \alpha/2, \widehat{T}_n > z_{\alpha/2} + \delta_0 \right) \\
 & \quad \quad \quad + \mathbb{P} \left( G_n(\widehat{T}_n) \leq \alpha/2, \widehat{T}_n < z_{\alpha/2} - \delta_0 \right) \\
 & \quad (\text{Property (i)}) = \mathbb{P} \left( G_n(\widehat{T}_n) \leq \alpha/2, \widehat{T}_n \in [z_{\alpha/2} - \delta_0, z_{\alpha/2} + \delta_0] \right) + \mathbb{P} \left( \widehat{T}_n < z_{\alpha/2} - \delta_0 \right)
 \end{aligned}$$

$$1901 \quad = \mathbb{P}\left(\widehat{T}_n \leq G_n^{-1}(\alpha/2), \widehat{T}_n \in [z_{\alpha/2} - \delta_0, z_{\alpha/2} + \delta_0]\right) + \mathbb{P}\left(\widehat{T}_n < z_{\alpha/2} - \delta_0\right)$$

$$1902 \quad (\text{Since } G_n^{-1}(\alpha/2) \geq z_{\alpha/2} - \delta_0) = \mathbb{P}\left(\widehat{T}_n \leq G_n^{-1}(\alpha/2)\right) = F_{\widehat{T}_n}(G_n^{-1}(\alpha/2))$$

$$1903 \quad = G_n(G_n^{-1}(\alpha/2)) + O(\mathcal{M}(\rho_n, n; R)) = \alpha/2 + O(\mathcal{M}(\rho_n, n; R))$$

1904 The other term  $\mathbb{P}\left(G_n(\widehat{T}_n) \geq 1 - \alpha/2\right)$  can be handled exactly similarly, and the proof of  
1905 part 1 of Theorem 4.2 is completed.

Now we move on to prove part 2 of the theorem.  $|c_n - d_n| = \omega(\rho_n^s \cdot n^{-1/2})$ . When  $H_a$  is true, we have  $\mu_n = d_n$ , and rewrite

$$\widehat{T}_n := \frac{\widehat{U}_n - d_n}{\widehat{S}_n} + \frac{d_n - c_n}{\widehat{S}_n}$$

Since  $\widehat{S}_n = \widetilde{O}_p(\rho_n^s \cdot n^{-1/2})$ , we have

$$\left| \frac{d_n - c_n}{\widehat{S}_n} \right| \xrightarrow{p} \infty, \quad \text{and therefore,} \quad |\widehat{T}_n| \xrightarrow{p} \infty$$

1906 By definition of Type-II error, this finishes the proof of part 2 of Theorem 4.2.

1907

□

1908 **PROOF OF THEOREM 4.3.** We first prove (4.6). By definition, we have

$$1909 \quad \left| F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha}) - \alpha \right| = F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha}) - \alpha \leq F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha}) - F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha} - 0^+)$$

$$1910 \quad \leq \left| F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha}) - G_n(q_{\widehat{T}_n; \alpha}) \right| + \left| G_n(q_{\widehat{T}_n; \alpha}) - G_n(q_{\widehat{T}_n; \alpha} - 0^+) \right|$$

$$1911 \quad + \left| G_n(q_{\widehat{T}_n; \alpha} - 0^+) - F_{\widehat{T}_n}(q_{\widehat{T}_n; \alpha} - 0^+) \right|$$

$$1912 \quad \leq O(\mathcal{M}(\rho_n, n; R)) + 0^+ = O(\mathcal{M}(\rho_n, n; R))$$

1913 where  $0^+$  represents an arbitrarily small positive number that may depend on  $n$ , and in the  
1914 last line we used the fact that  $G_n(x)$  is globally uniformly Lipschitz. This proves (4.6).

1915 Then we prove the horizontal bound (4.7). Define

$$1916 \quad \widetilde{q}_{\widehat{T}_n; \alpha} := z_\alpha - \frac{1}{\sqrt{n} \cdot \xi_1^3} \cdot \left\{ \frac{2z_\alpha^2 + 1}{6} \cdot \mathbb{E}[g_1^3(X_1)] \right.$$

$$1917 \quad \left. + \frac{r-1}{2} \cdot (z_\alpha^2 + 1) \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\}$$

1918 For convenience, let us simply denote the  $n^{-1/2}$  term in the Edgeworth expansion by  $\Gamma(x)$ :

$$1919 \quad \Gamma(x) := \frac{1}{\xi_1^3} \cdot \left\{ \frac{2x^2 + 1}{6} \cdot \mathbb{E}[g_1^3(X_1)] + \frac{r-1}{2} \cdot (x^2 + 1) \mathbb{E}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\}$$

$$1920 \quad \widehat{\Gamma}(x) := \frac{1}{\widehat{\xi}_1^3} \cdot \left\{ \frac{2x^2 + 1}{6} \cdot \widehat{\mathbb{E}}[g_1^3(X_1)] + \frac{r-1}{2} \cdot (x^2 + 1) \widehat{\mathbb{E}}[g_1(X_1)g_1(X_2)g_2(X_1, X_2)] \right\}$$

1921 We have

$$1922 \quad G_n(x) = \Phi(x) + n^{-1/2} \cdot \Gamma(x)\varphi(x)$$

$$1923 \quad \tilde{q}_{\hat{T}_n; \alpha} = z_\alpha - n^{-1/2} \cdot \Gamma(z_\alpha)$$

$$1924 \quad \hat{q}_{\hat{T}_n; \alpha} = z_\alpha - n^{-1/2} \cdot \hat{\Gamma}(z_\alpha)$$

1925 Then the proof of Theorem 3.2 immediately implies that  $|\hat{q}_{\hat{T}_n; \alpha} - \tilde{q}_{\hat{T}_n; \alpha}| = \tilde{O}_p(\rho_n^{-1} n^{-1} \log^{1/2} n)$ .  
 1926 Mimicking the inversion formula in [65], we have

$$1927 \quad G_n \left( x - \frac{1}{\sqrt{n}} \cdot \Gamma(x) \right) = \Phi \left( x - \frac{1}{\sqrt{n}} \cdot \Gamma(x) \right) + \frac{1}{\sqrt{n}} \cdot \Gamma \left( x - \frac{1}{\sqrt{n}} \cdot \Gamma(x) \right) \varphi \left( x - \frac{1}{\sqrt{n}} \cdot \Gamma(x) \right)$$

$$1928 \quad (8.123) \quad = \Phi(x) + O(n^{-1}).$$

1929 Also notice that the remainder bound in (8.123) holds uniformly over all  $x \in \mathbb{R}$ . As [65]  
 1930 pointed out, in a more general setting, the inversion formula (8.123) might not always have  
 1931 a uniform  $O(n^{-1})$  error bound, when the leading term in the Edgeworth expansion contains  
 1932 a jump function component, in which case the uniform error bound of the Cornish-Fisher  
 1933 expansion is just  $O(n^{-1/2})$ . But in our setting,  $\Gamma(x)$  is always continuous, and moreover,  
 1934 Lipschitz, so [65]'s remark would not be a concern.

1935 We continue our proof. By Theorem 3.1 and (8.123), we have

$$1936 \quad (8.124) \quad G_n(\tilde{q}_{\hat{T}_n; \alpha}) = \alpha + O(n^{-1})$$

$$1937 \quad (8.125) \quad G_n(q_{\hat{T}_n; \alpha}) = F_{\hat{T}_n}(q_{\hat{T}_n; \alpha}) + O(\mathcal{M}(\rho_n, n; R))$$

1938 Since for any given  $\alpha$  and large enough  $n$ , properties (i) and (ii) of  $G_n(\cdot)$ , with “ $\alpha/2, \epsilon_0,$   
 1939  $\delta_0, n_0$ ” replaced by “ $\alpha, \epsilon'_0, \delta'_0, n'_0$ ”, around a neighborhood of  $z_\alpha$ . This yields that for large  
 1940 enough  $n$ , both  $q_{\hat{T}_n; \alpha}$  and  $\tilde{q}_{\hat{T}_n; \alpha}$  belong to  $[z_\alpha - \epsilon'_0, z_\alpha + \epsilon'_0]$ . Then using the invertibility and  
 1941 the Lipschitz property of the inverse function of  $G_n(\cdot)$  within this compact neighborhood,  
 1942 we have

$$1943 \quad |\tilde{q}_{\hat{T}_n; \alpha} - q_{\hat{T}_n; \alpha}| \leq |G_n(\tilde{q}_{\hat{T}_n; \alpha}) - G_n(q_{\hat{T}_n; \alpha})|$$

$$1944 \quad (8.126) \quad = O(\mathcal{M}(\rho_n, n; R))$$

1945 Combining this with the error bound on  $|\tilde{q}_{\hat{T}_n; \alpha} - \hat{q}_{\hat{T}_n; \alpha}|$  we obtained earlier finishes the proof  
 1946 of the horizontal error bound (4.7).

Now we prove the vertical error bound (4.8). Here we should be careful that  $\mathbb{P}(\hat{T}_n \leq \hat{q}_{\hat{T}_n; \alpha})$   
 does *not* equal  $F_{\hat{T}_n}(\hat{q}_{\hat{T}_n; \alpha})$ , as the former is non-random and the latter is random. In order to  
 study  $\mathbb{P}(\hat{T}_n \leq \hat{q}_{\hat{T}_n; \alpha})$ , we seek the help from  $\tilde{q}_{\hat{T}_n; \alpha}$  and appeal to the basic definition. By the  
 horizontal error bound, we know that with probability  $1 - O(n^{-1})$ , we have

$$|\tilde{q}_{\hat{T}_n; \alpha} - \hat{q}_{\hat{T}_n; \alpha}| \leq C \cdot \mathcal{M}(\rho_n, n; R)$$

1947 for some constant  $C > 0$ . This yields that under the above event

$$1948 \quad F_{\hat{T}_n}(\tilde{q}_{\hat{T}_n; \alpha} - C \cdot \mathcal{M}(\rho_n, n; R)) = \mathbb{P}(\hat{T}_n \leq \tilde{q}_{\hat{T}_n; \alpha} - C \cdot \mathcal{M}(\rho_n, n; R))$$

$$1949 \quad \leq \mathbb{P}(\hat{T}_n \leq \hat{q}_{\hat{T}_n; \alpha})$$

$$1950 \quad \leq \mathbb{P}(\hat{T}_n \leq \tilde{q}_{\hat{T}_n; \alpha} + C \cdot \mathcal{M}(\rho_n, n; R)) = F_{\hat{T}_n}(\tilde{q}_{\hat{T}_n; \alpha} + C \cdot \mathcal{M}(\rho_n, n; R))$$

1951 Recall that  $G_n(\cdot)$  is globally Lipschitz for large enough  $n$ , we have

$$1952 \quad F_{\hat{T}_n}(\tilde{q}_{\hat{T}_n; \alpha} + C \cdot \mathcal{M}(\rho_n, n; R)) = G_n(\tilde{q}_{\hat{T}_n; \alpha} + C \cdot \mathcal{M}(\rho_n, n; R)) + O(\mathcal{M}(\rho_n, n; R))$$

$$1953 \quad = G_n(\tilde{q}_{\hat{T}_n; \alpha}) + O(\mathcal{M}(\rho_n, n; R)) = \alpha + O(\mathcal{M}(\rho_n, n; R))$$

1954 This proves the vertical error bound (4.8) and concludes the proof of Theorem 4.3.

1955

□

**9. Additional simulation results.**

9.1. *Additional results in Simulation 5.1.* In this section, we show additional simulation results under different network sparsity settings. We tested  $\rho_n \asymp n^{-1/4}, n^{-1/3}$  and  $n^{-1/2}$ . Notice that some of these settings constitute violations of our assumptions  $\rho_n$  assumptions. We adjusted the constant factors in  $\rho_n$  such that all settings start with roughly equal network densities for  $n = 10$ . Results are shown in Figures 6–8 (errors) and Figures 9–11 (time costs), where error bars show standard deviations.

The plots show that the accuracy of all methods depreciate as the network becomes sparser. Recall that our loss function is the error in approximating  $F_{\hat{T}_n}$ , and that  $\hat{T}_n$  is normalized by the denominator  $\hat{S}_n \asymp \rho_n^s \cdot n^{-1/2}$ , it is therefore understandable that sparser networks are more difficult. Apart from that error bounds would depreciate with a smaller  $\rho_n$ , as in our Theorems 3.1 and 3.2; the performances of our method in some scenarios also seemed to be limited by numerical accuracy, possibly in the Monte Carlo evaluations of the true  $F_{\hat{T}_n}$ . But overall, our method remains the best performer and higher-order accurate in scenarios where the sparsity assumptions are satisfied. The time cost plots can be interpreted similarly to that in the main paper text.

9.2. *Additional results in Simulation 5.2.* In this subsection we present the results for more settings, including  $n = 160$  and more sparsity levels. Results are reported in Tables 6–23.

TABLE 6  
*Performance measures of 95% confidence intervals*  
 $n = 80, \rho_n \asymp n^{-1/4}, \text{ graphon: block model}$

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.960(0.196)	0.954(0.209)	0.957(0.203)	0.953(0.212)
	Length = 0.084(0.009)	0.024(0.005)	0.144(0.024)	0.087(0.020)
	LogTime = -8.419(0.135)	-7.450(0.118)	-7.404(0.108)	-6.405(0.774)
Norm. Approx.	0.953(0.212)	0.935(0.247)	0.944(0.230)	0.933(0.251)
	0.084(0.009)	0.024(0.005)	0.144(0.024)	0.087(0.020)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.830(0.376)	0.856(0.351)	0.832(0.374)	0.858(0.349)
	0.058(0.008)	0.019(0.004)	0.106(0.019)	0.069(0.016)
	-2.599(0.028)	-2.137(0.020)	-2.195(0.031)	-0.987(0.015)
Green and Shalizi [61]	0.934(0.249)	0.936(0.245)	0.942(0.234)	0.938(0.241)
	0.082(0.011)	0.027(0.006)	0.145(0.028)	0.089(0.023)
	-1.202(0.019)	0.548(0.051)	0.085(0.052)	0.353(0.012)
Levin and Levina [93]	0.954(0.210)	0.956(0.205)	0.956(0.205)	0.952(0.214)
	0.085(0.011)	0.026(0.006)	0.150(0.028)	0.094(0.023)
	-1.193(0.014)	0.574(0.040)	0.074(0.044)	0.403(0.006)

9.3. *Additional results in Simulation 5.3.* In this simulation, all settings are carried over from Simulation 5.3 except that  $n = 80$ . The results are shown in Figure 12. We observed the anticipated depreciation in the performances of all methods, while our method maintains a consistent advantage over the closest competitors.

The impact of  $\rho_n$  on the computation time is a subtle topic. Since our simulation runs across dense and sparse regimes, for simplicity and wide-applicability of the code, we did not engage sparse matrix computation procedures. Consequently, the time cost for all  $\rho_n$ 's are nearly the same for all methods. Here, we only show the time costs for  $n = 80$  in Figure 13, and the analogous plot for  $n = 160$  looks exactly similar and is thus omitted here. We leave the study of improving computational efficiency for sparse  $\rho_n$ 's to future work.

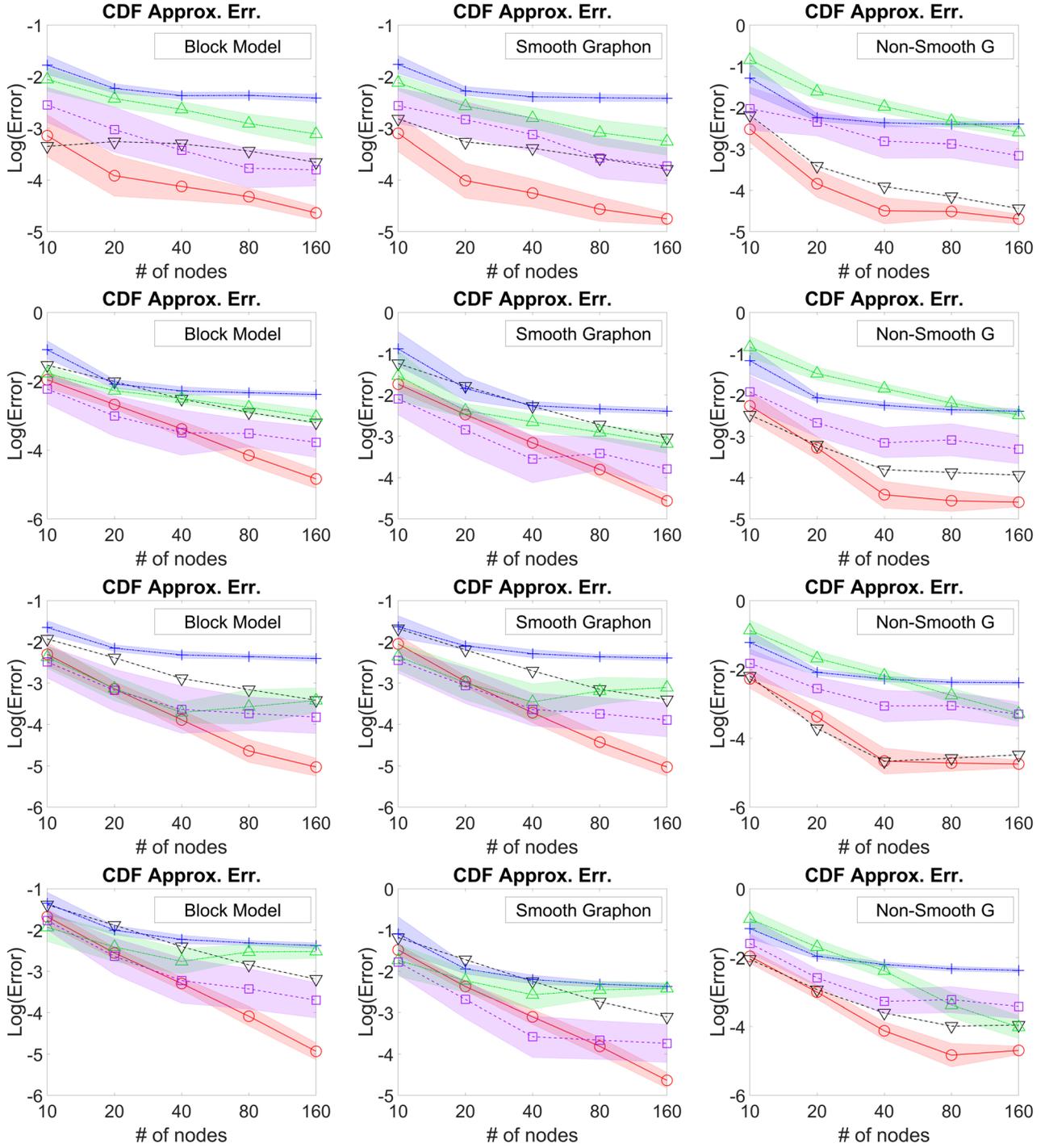


Fig 6: CDF approximation errors,  $\rho_n \asymp n^{-1/4}$ . Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked red circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

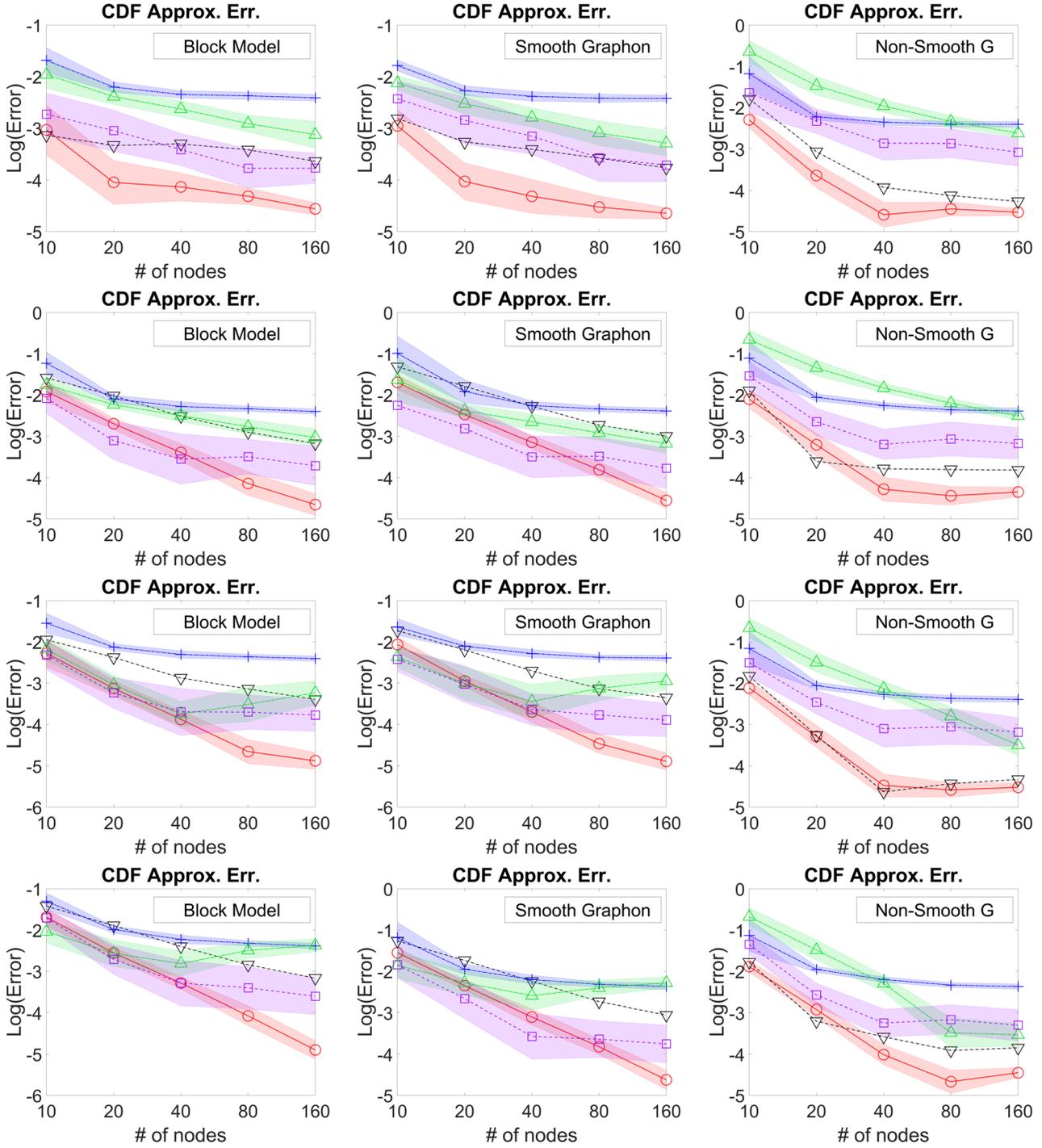


Fig 7: CDF approximation errors,  $\rho_n \asymp n^{-1/3}$ . Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked red circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

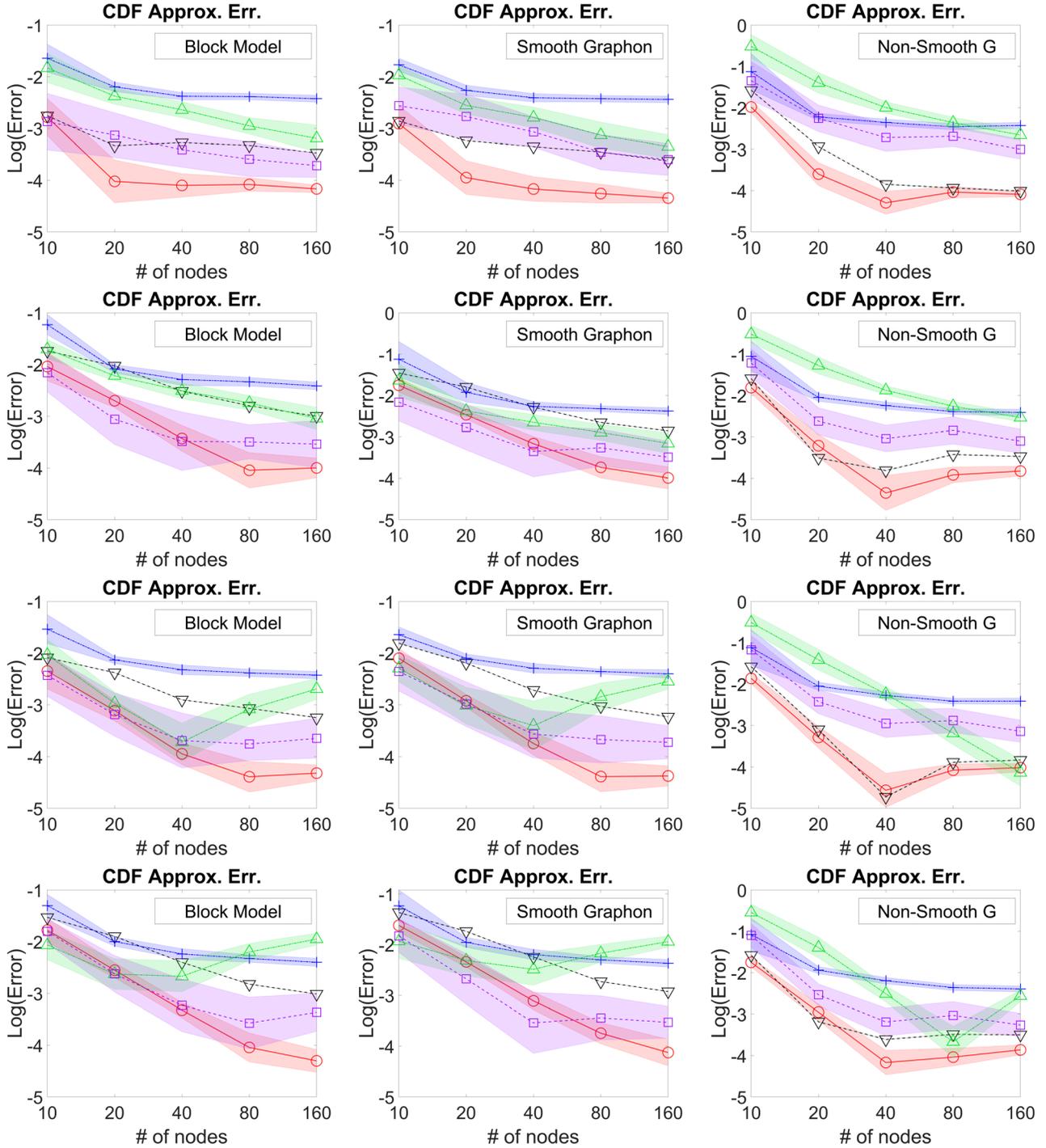


Fig 8: CDF approximation errors,  $\rho_n \asymp n^{-1/2}$ . Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

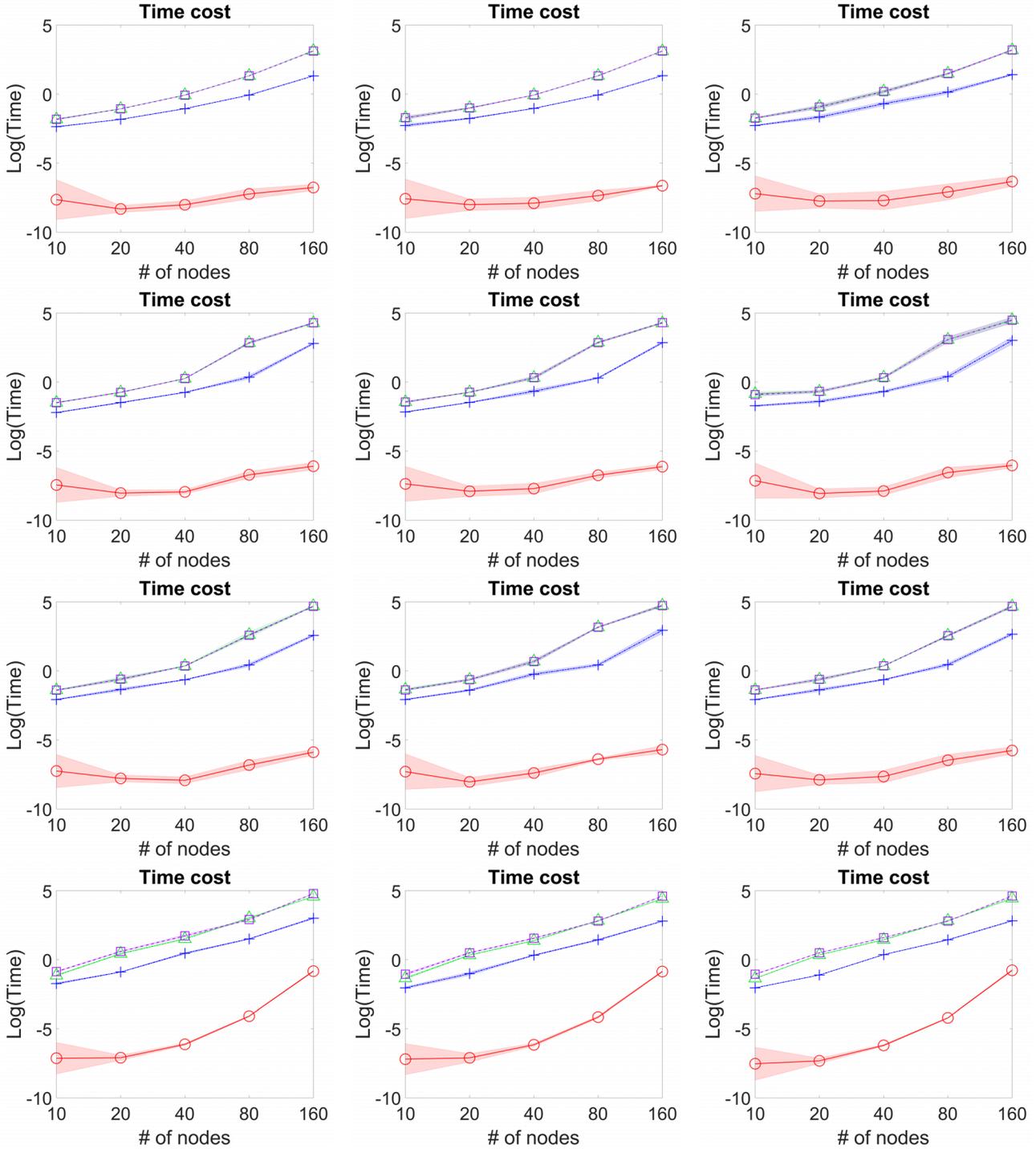


Fig 9: CDF approximation times,  $\rho_n \asymp n^{-1/4}$ . Both axes are  $\log(e)$ -scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked red circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

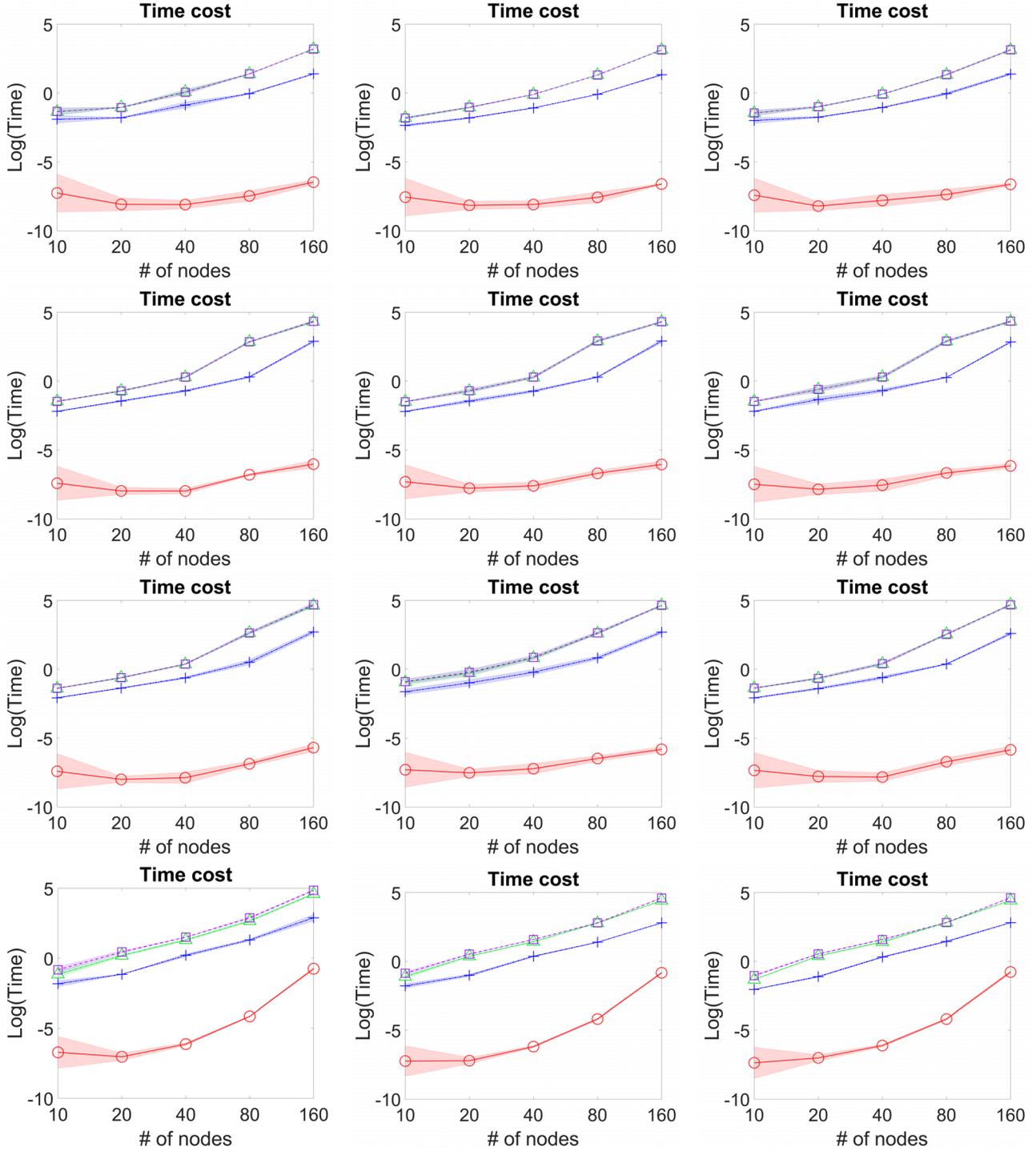


Fig 10: CDF approximation times,  $\rho_n \asymp n^{-1/3}$ . Both axes are  $\log(e)$ -scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked red circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

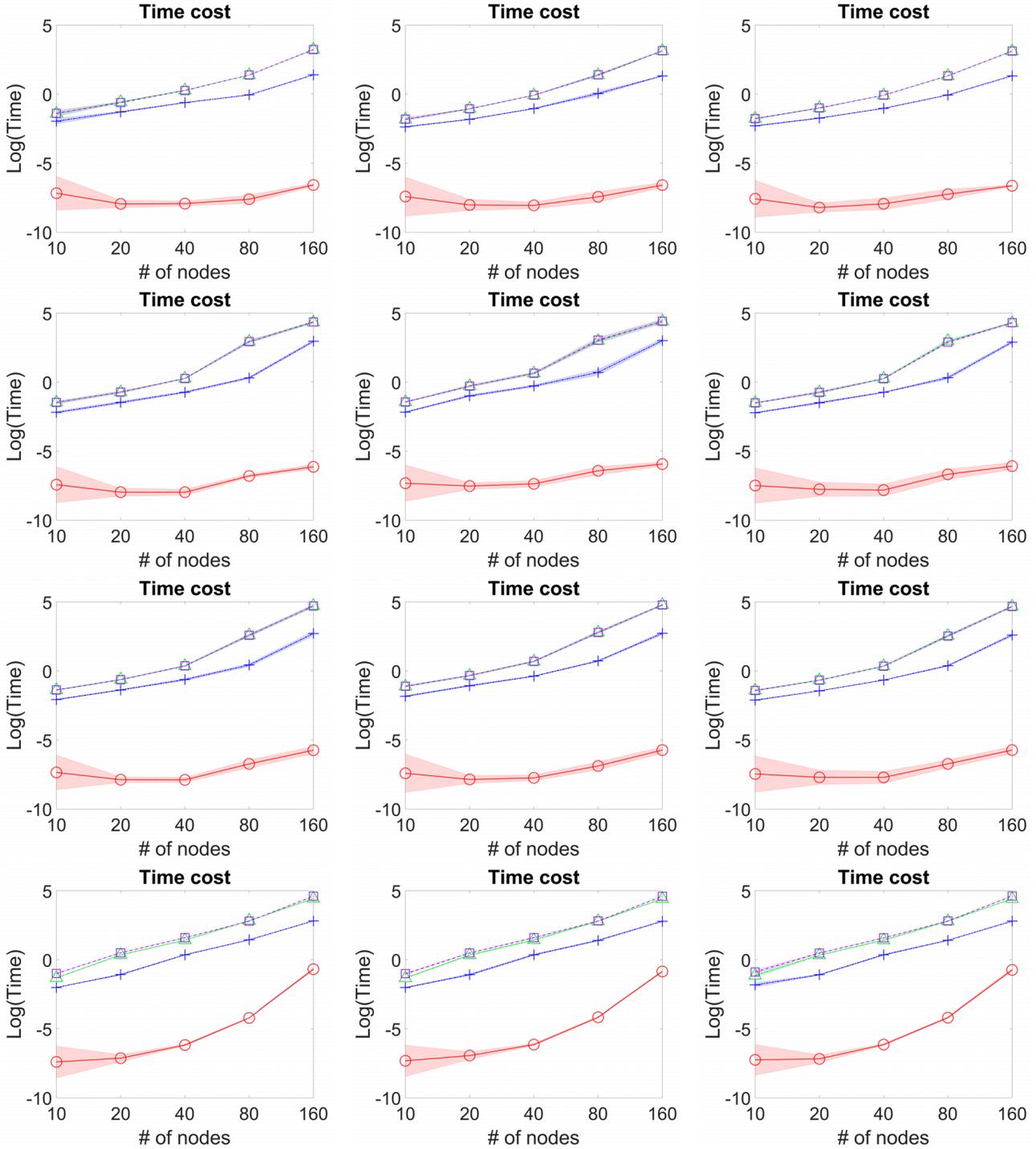


Fig 11: CDF approximation times,  $\rho_n \asymp n^{-1/2}$ . Both axes are  $\log(e)$ -scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked red circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

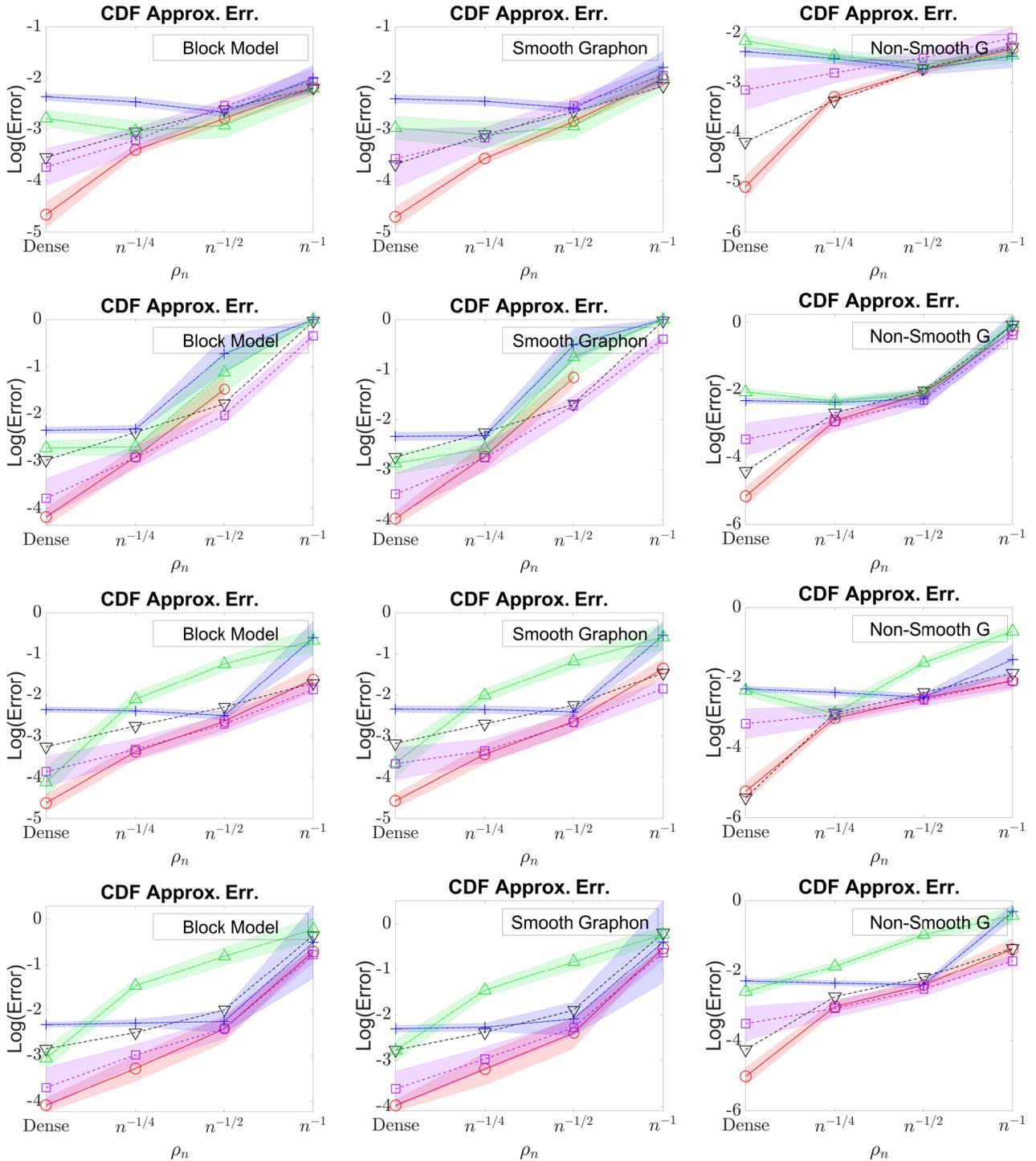


Fig 12: Impact of sparsity on approximation errors,  $n = 80$ . Both axes are  $\log(\epsilon)$ -scaled.

**Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93]. We regarded  $N(0, 1)$  as zero time cost so does not appear in the time cost plot.

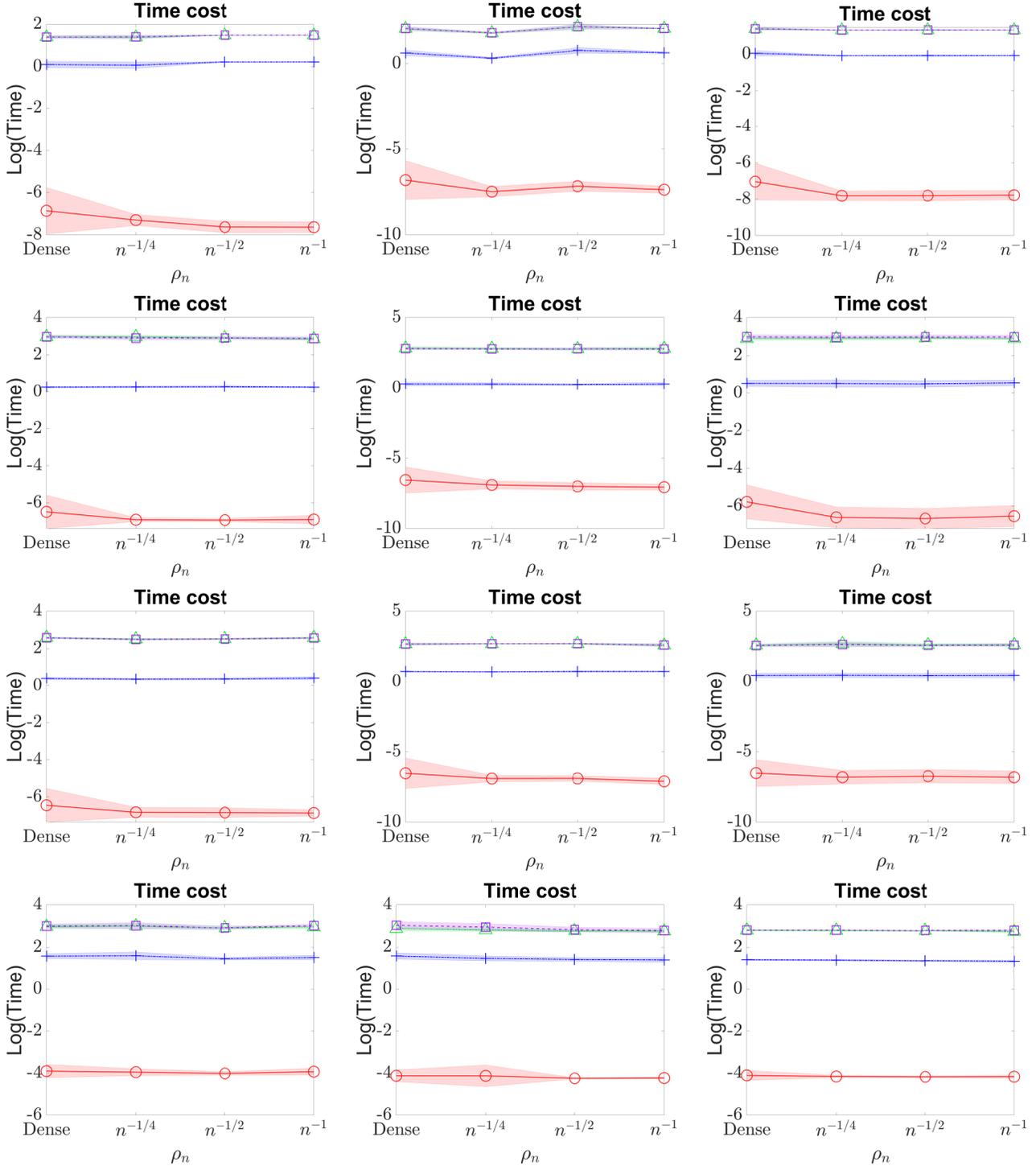


Fig 13: Impact of sparsity on time cost,  $n = 80$ . We used regular (non-sparse) matrix variables in MATLAB. Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93]. We regarded  $N(0, 1)$  as zero time cost so does not appear in the time cost plot.

TABLE 7  
*Performance measures of 95% confidence intervals*  
 $n = 80$ ,  $\rho_n \asymp n^{-1/4}$ , *graphon: smooth graphon*

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.961(0.193)	0.942(0.234)	0.955(0.208)	0.945(0.229)
	Length = 0.078(0.008)	0.013(0.003)	0.101(0.018)	0.050(0.013)
	LogTime = -8.226(0.044)	-7.468(0.116)	-7.439(0.073)	-6.599(0.687)
Norm. Approx.	0.953(0.211)	0.922(0.268)	0.939(0.239)	0.923(0.266)
	0.078(0.008)	0.013(0.003)	0.101(0.018)	0.050(0.013)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.822(0.383)	0.850(0.357)	0.852(0.355)	0.848(0.359)
	0.056(0.007)	0.011(0.003)	0.079(0.015)	0.045(0.012)
	-2.562(0.008)	-2.111(0.099)	-2.198(0.044)	-0.995(0.012)
Green and Shalizi [61]	0.928(0.259)	0.948(0.222)	0.934(0.249)	0.944(0.230)
	0.078(0.010)	0.015(0.004)	0.105(0.021)	0.054(0.015)
	-1.148(0.010)	0.504(0.057)	0.104(0.102)	0.322(0.017)
Levin and Levina [93]	0.942(0.234)	0.960(0.196)	0.954(0.210)	0.962(0.191)
	0.082(0.010)	0.015(0.004)	0.111(0.022)	0.058(0.016)
	-1.146(0.004)	0.514(0.048)	0.056(0.055)	0.387(0.011)

TABLE 8  
*Performance measures of 95% confidence intervals*  
 $n = 80$ ,  $\rho_n \asymp n^{-1/4}$ , *graphon: non-smooth graphon*

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.960(0.196)	0.961(0.193)	0.961(0.193)	0.963(0.189)
	Length = 0.101(0.008)	0.083(0.007)	0.310(0.022)	0.329(0.029)
	LogTime = -7.996(0.058)	-7.652(0.145)	-7.611(0.133)	-6.789(0.618)
Norm. Approx.	0.957(0.202)	0.955(0.208)	0.957(0.204)	0.956(0.206)
	0.101(0.008)	0.083(0.007)	0.310(0.022)	0.329(0.029)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.832(0.374)	0.838(0.369)	0.834(0.372)	0.838(0.369)
	0.070(0.008)	0.059(0.007)	0.216(0.023)	0.233(0.027)
	-2.559(0.048)	-2.151(0.029)	-2.129(0.029)	-1.000(0.042)
Green and Shalizi [61]	0.930(0.255)	0.934(0.249)	0.944(0.230)	0.950(0.218)
	0.097(0.011)	0.083(0.009)	0.301(0.030)	0.323(0.036)
	-1.152(0.027)	0.488(0.054)	0.144(0.041)	0.341(0.033)
Levin and Levina [93]	0.962(0.191)	0.972(0.165)	0.966(0.181)	0.970(0.171)
	0.101(0.011)	0.086(0.010)	0.314(0.031)	0.338(0.038)
	-1.145(0.027)	0.479(0.052)	0.141(0.040)	0.463(0.023)

9.4. *Additional simulation results for degree-corrected stochastic block models.* Here we present the simulation results under a degree-corrected stochastic block model [83]. We generate data from the stochastic block model `BlockModel` that we tested in Section 5, with the following degree correction function

$$\theta(x) := |\cos \pi \cdot (x - 1/2)|.$$

1985 The results are reported in Figures 14 – 17. We observe the clear advantage of our method  
 1986 over benchmarks, as predicted by our theory.

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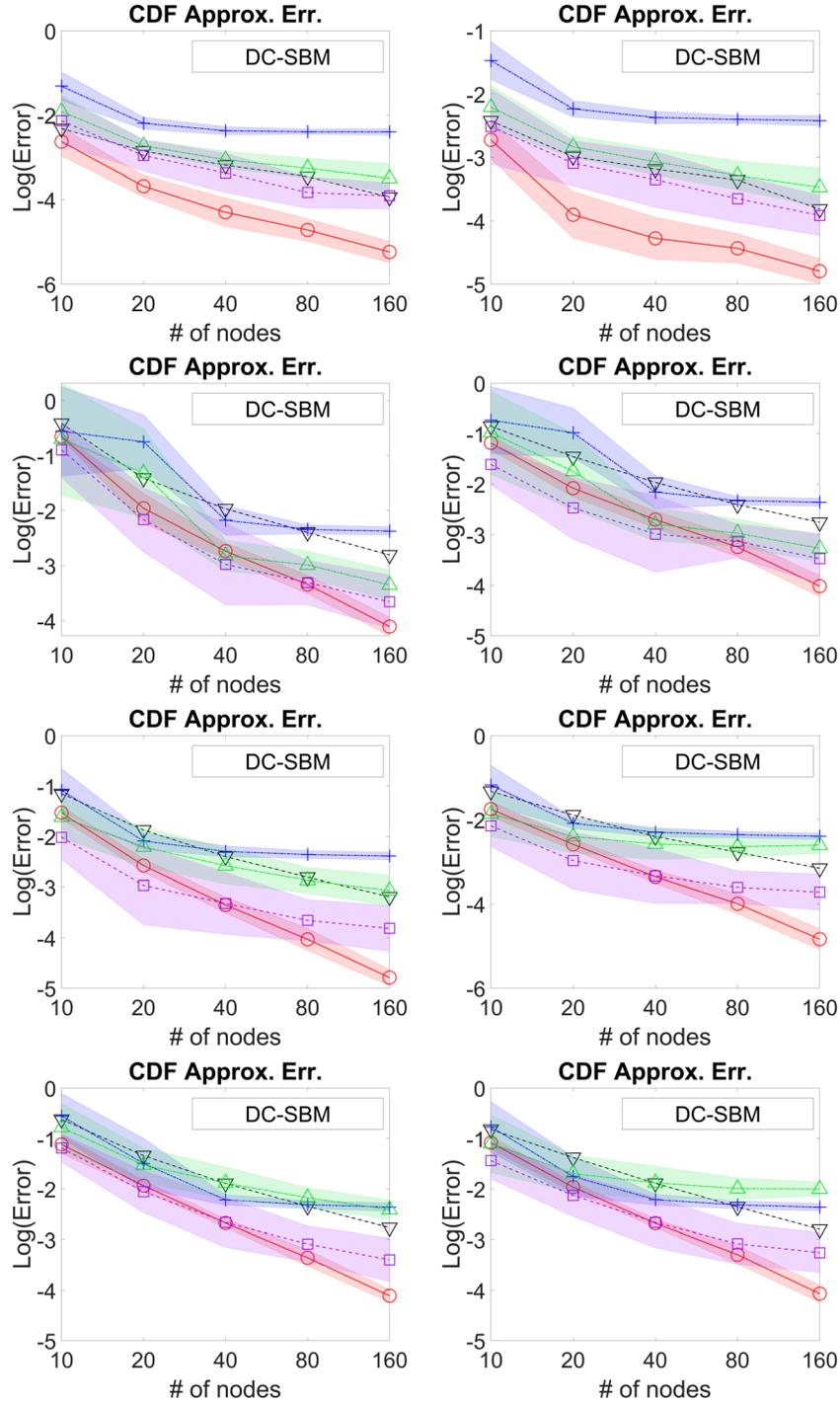


Fig 14: CDF approximation errors for degree-corrected stochastic block model. Sparsity: column 1:  $\rho_n = 1$ ; column 2:  $\rho_n = n^{-1/4}$  Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

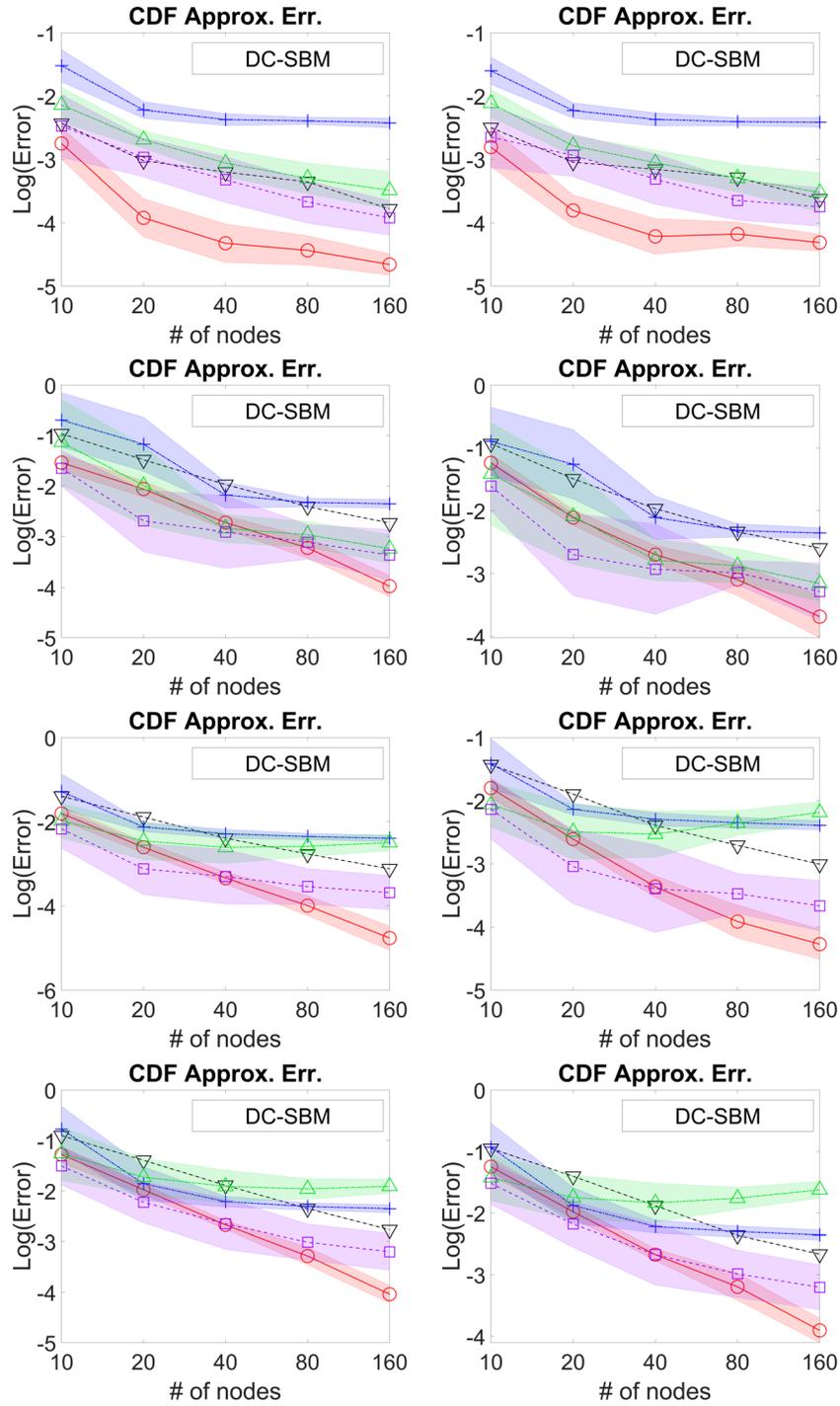


Fig 15: CDF approximation errors for degree-corrected stochastic block model. Sparsity: column 1:  $\rho_n \asymp n^{-1/3}$ ; column 2:  $\rho_n \asymp n^{-1/2}$  Both axes are  $\log(e)$ -scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

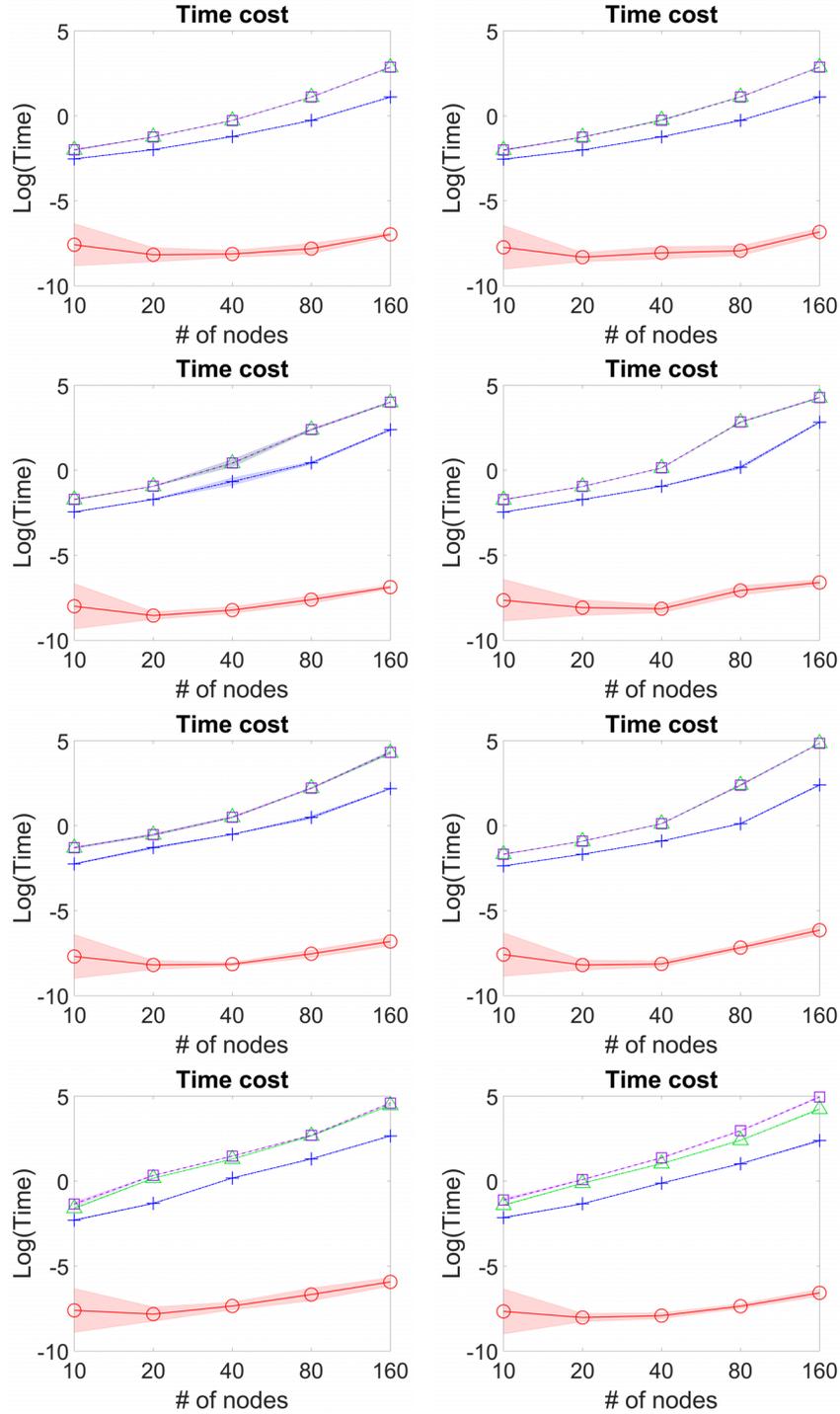


Fig 16: CDF approximation times for degree-corrected stochastic block model. Sparsity: column 1:  $\rho_n \approx 1$ ; column 2:  $\rho_n \approx n^{-1/4}$  Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0,1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\approx n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

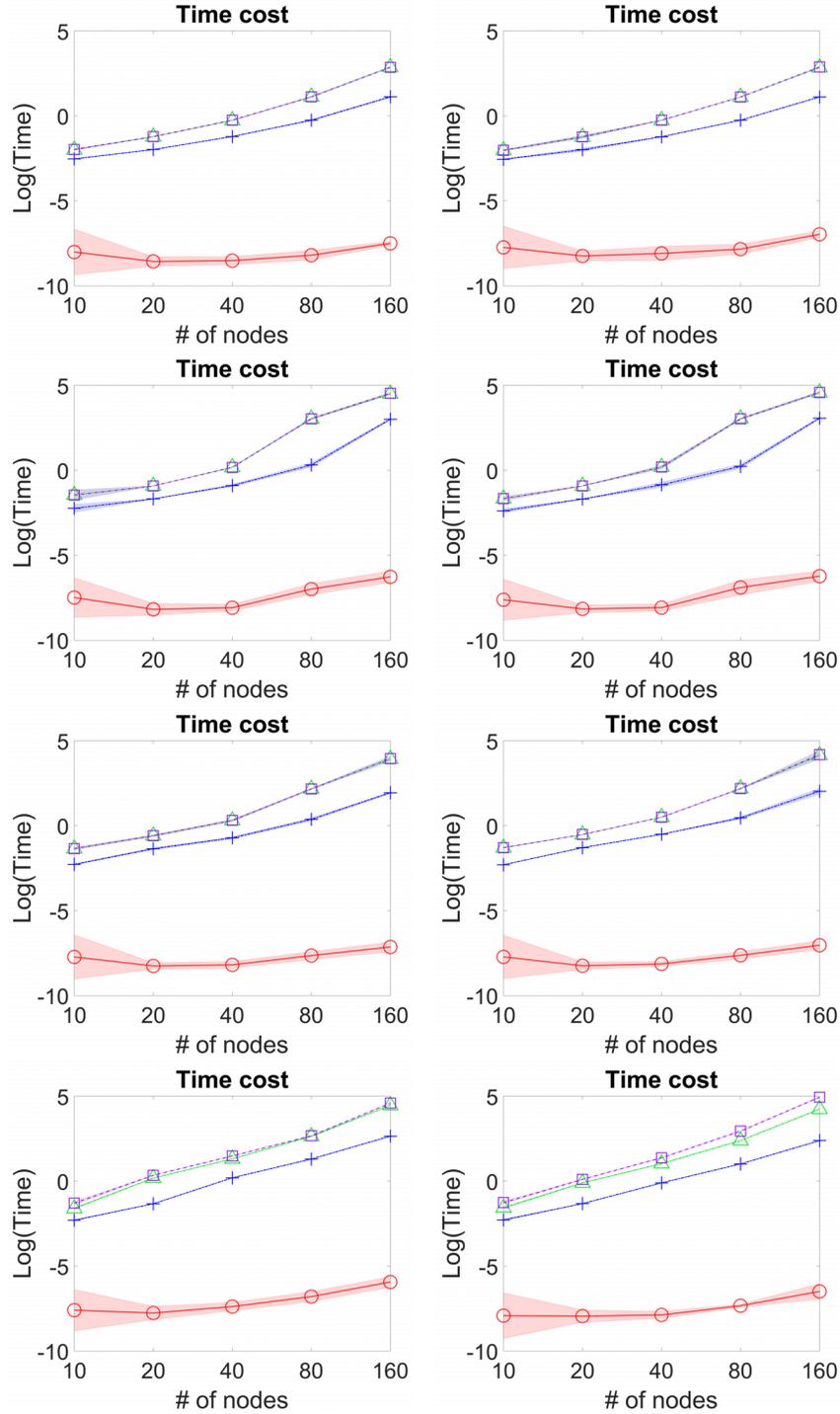


Fig 17: CDF approximation times for degree-corrected stochastic block model. Sparsity: column 1:  $\rho_n \approx n^{-1/3}$ ; column 2:  $\rho_n \approx n^{-1/2}$ . Both axes are log(e)-scaled. **Motifs:** row 1: Edge; row 2: Triangle; row 3: Vshape; row 4: ThreeStar. **Red solid curve marked circle:** our method (empirical Edgeworth); **black dashed curve marked down-triangle:**  $N(0, 1)$  approximation; **green dashed curve marked up-triangle:** re-sampling of  $A$  in [61]; **blue dashed curve marked plus:** [17] sub-sampling  $\asymp n$  nodes; **magenta dashed line with square markers:** ASE plug-in bootstrap in [93].

TABLE 9  
 Performance measures of 95% confidence intervals  
 $n = 80, \rho_n \asymp n^{-1/2}$ , graphon: block model

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.969(0.172)	0.956(0.206)	0.964(0.187)	0.953(0.212)
	Length = 0.046(0.005)	0.003(0.001)	0.037(0.007)	0.011(0.003)
	LogTime = -8.335(0.153)	-7.139(0.113)	-7.212(0.104)	-7.153(0.338)
Norm. Approx.	0.967(0.180)	0.946(0.226)	0.956(0.206)	0.945(0.229)
	0.046(0.005)	0.003(0.001)	0.037(0.007)	0.011(0.003)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.824(0.381)	0.848(0.359)	0.840(0.367)	0.852(0.355)
	0.031(0.005)	0.003(0.001)	0.027(0.006)	0.009(0.003)
	-2.588(0.008)	-2.107(0.084)	-2.123(0.009)	-1.027(0.008)
Green and Shalizi [61]	0.952(0.214)	0.936(0.245)	0.940(0.238)	0.910(0.286)
	0.044(0.007)	0.004(0.001)	0.035(0.008)	0.010(0.003)
	-1.159(0.010)	0.500(0.039)	0.199(0.045)	0.341(0.021)
Levin and Levina [93]	0.972(0.165)	0.966(0.181)	0.966(0.181)	0.962(0.191)
	0.047(0.007)	0.004(0.001)	0.040(0.009)	0.012(0.004)
	-1.148(0.005)	0.521(0.036)	0.220(0.027)	0.444(0.009)

TABLE 10  
 Performance measures of 95% confidence intervals  
 $n = 80, \rho_n \asymp n^{-1/2}$ , graphon: smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.967(0.179)	0.931(0.253)	0.956(0.205)	0.930(0.256)
	Length = 0.042(0.005)	0.002(0.001)	0.026(0.005)	0.006(0.002)
	LogTime = -8.213(0.047)	-7.618(0.111)	-7.152(0.107)	-7.147(0.318)
Norm. Approx.	0.963(0.189)	0.932(0.252)	0.948(0.223)	0.926(0.262)
	0.042(0.005)	0.002(0.001)	0.026(0.005)	0.006(0.002)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.824(0.381)	0.872(0.334)	0.834(0.372)	0.852(0.355)
	0.029(0.004)	Inf(NaN)	0.021(0.005)	0.007(0.002)
	-2.575(0.006)	-2.185(0.026)	-2.114(0.018)	-0.854(0.036)
Green and Shalizi [61]	0.950(0.218)	0.958(0.201)	0.940(0.238)	0.920(0.272)
	0.041(0.006)	0.002(0.001)	0.025(0.006)	0.006(0.002)
	-1.154(0.010)	0.497(0.101)	0.176(0.039)	0.465(0.027)
Levin and Levina [93]	0.956(0.205)	0.974(0.159)	0.960(0.196)	0.968(0.176)
	0.044(0.006)	0.002(0.001)	0.029(0.007)	0.008(0.003)
	-1.150(0.006)	0.487(0.067)	0.181(0.027)	0.440(0.027)

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TABLE 11  
 Performance measures of 95% confidence intervals  
 $n = 80$ ,  $\rho_n \asymp n^{-1/2}$ , graphon: non-smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.974(0.159)	0.974(0.159)	0.974(0.160)	0.973(0.164)
	Length = 0.059(0.005)	0.011(0.002)	0.087(0.009)	0.045(0.006)
	LogTime = -8.196(0.051)	-7.297(0.134)	-7.314(0.140)	-7.011(0.405)
Norm. Approx.	0.973(0.162)	0.969(0.174)	0.973(0.164)	0.970(0.171)
	0.059(0.005)	0.011(0.002)	0.087(0.009)	0.045(0.006)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.850(0.357)	0.852(0.355)	0.852(0.355)	0.860(0.347)
	0.040(0.005)	0.008(0.001)	0.060(0.008)	0.033(0.006)
	-2.543(0.115)	-2.118(0.022)	-2.195(0.021)	-0.965(0.047)
Green and Shalizi [61]	0.944(0.230)	0.942(0.234)	0.950(0.218)	0.946(0.226)
	0.055(0.006)	0.011(0.002)	0.082(0.011)	0.041(0.008)
	-1.135(0.058)	0.590(0.089)	0.134(0.041)	0.383(0.038)
Levin and Levina [93]	0.966(0.181)	0.972(0.165)	0.968(0.176)	0.968(0.176)
	0.059(0.006)	0.012(0.002)	0.090(0.011)	0.048(0.008)
	-1.130(0.059)	0.556(0.029)	0.121(0.036)	0.498(0.024)

TABLE 12  
 Performance measures of 95% confidence intervals  
 $n = 80$ ,  $\rho_n \asymp n^{-1}$ , graphon: block model

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.987(0.115)	0.000(0.000)	0.949(0.219)	0.646(0.478)
	Length = 0.016(0.002)	0.000(0.000)	0.003(0.001)	0.000(0.000)
	LogTime = -8.487(0.176)	-7.230(0.105)	-7.422(0.114)	-7.183(0.322)
Norm. Approx.	0.983(0.129)	0.708(0.455)	0.959(0.199)	0.914(0.280)
	0.016(0.002)	0.000(0.000)	0.003(0.001)	0.000(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.904(0.295)	0.620(0.486)	0.898(0.303)	0.906(0.292)
	0.011(0.002)	Inf(NaN)	Inf(NaN)	Inf(NaN)
	-2.579(0.005)	-2.121(0.050)	-2.137(0.015)	-1.080(0.006)
Green and Shalizi [61]	0.972(0.165)	0.672(0.470)	0.896(0.306)	0.768(0.423)
	0.015(0.002)	Inf(NaN)	0.002(0.001)	0.000(0.000)
	-1.158(0.005)	0.489(0.050)	0.203(0.030)	0.299(0.015)
Levin and Levina [93]	0.984(0.126)	0.706(0.456)	0.996(0.063)	0.996(0.063)
	0.019(0.024)	0.000(0.000)	Inf(NaN)	Inf(NaN)
	-1.146(0.004)	0.509(0.031)	0.223(0.026)	0.434(0.006)

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TABLE 13  
 Performance measures of 95% confidence intervals  
 $n = 80, \rho_n \asymp n^{-1}$ , graphon: smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.984(0.126)	0.000(0.000)	0.909(0.287)	0.503(0.500)
	Length = 0.014(0.002)	0.000(0.000)	0.002(0.001)	0.000(0.000)
	LogTime = -8.205(0.054)	-7.243(0.071)	-7.276(0.103)	-7.098(0.400)
Norm. Approx.	0.981(0.138)	0.426(0.495)	0.948(0.223)	0.875(0.331)
	0.014(0.002)	0.000(0.000)	0.002(0.001)	0.000(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.910(0.288)	0.300(0.461)	0.910(0.288)	0.920(0.273)
	0.009(0.001)	Inf(NaN)	Inf(NaN)	Inf(NaN)
	-1.632(0.005)	-1.182(0.044)	-1.181(0.041)	-0.141(0.009)
Green and Shalizi [61]	0.940(0.239)	0.380(0.488)	0.900(0.302)	0.750(0.435)
	0.013(0.002)	Inf(NaN)	0.001(0.001)	0.000(0.000)
	-0.217(0.010)	1.553(0.030)	1.142(0.042)	1.167(0.016)
Levin and Levina [93]	0.980(0.141)	0.380(0.488)	0.970(0.171)	0.990(0.100)
	41.865(418.438)	Inf(NaN)	Inf(NaN)	Inf(NaN)
	-0.213(0.008)	1.567(0.019)	1.120(0.018)	1.245(0.014)

TABLE 14  
 Performance measures of 95% confidence intervals  
 $n = 80, \rho_n \asymp n^{-1}$ , graphon: non-smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.989(0.103)	0.896(0.305)	0.980(0.139)	0.911(0.285)
	Length = 0.022(0.002)	0.000(0.000)	0.007(0.001)	0.001(0.000)
	LogTime = -8.242(0.096)	-7.355(0.085)	-7.356(0.088)	-7.101(0.343)
Norm. Approx.	0.989(0.106)	0.963(0.189)	0.980(0.142)	0.962(0.192)
	0.022(0.002)	0.000(0.000)	0.007(0.001)	0.001(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.894(0.308)	0.908(0.289)	0.916(0.278)	0.908(0.289)
	0.015(0.002)	Inf(NaN)	0.005(0.001)	0.001(0.000)
	-2.571(0.015)	-2.171(0.012)	-2.128(0.020)	-1.043(0.007)
Green and Shalizi [61]	0.964(0.186)	0.968(0.176)	0.936(0.245)	0.848(0.359)
	0.020(0.002)	Inf(NaN)	0.006(0.001)	0.001(0.000)
	-1.191(0.012)	0.567(0.120)	0.219(0.024)	0.343(0.012)
Levin and Levina [93]	0.986(0.118)	0.984(0.126)	0.992(0.089)	0.992(0.089)
	0.023(0.003)	0.000(0.000)	0.008(0.002)	0.001(0.000)
	-1.183(0.010)	0.529(0.083)	0.236(0.034)	0.450(0.013)

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TABLE 15  
 Performance measures of 95% confidence intervals  
 $n = 160$ ,  $\rho_n \asymp n^{-1/4}$ , graphon: block model

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.957(0.204)	0.954(0.211)	0.954(0.210)	0.951(0.216)
	Length = 0.048(0.004)	0.010(0.001)	0.070(0.008)	0.036(0.006)
	LogTime = -7.068(0.086)	-6.643(0.137)	-6.161(0.337)	-6.132(0.241)
Norm. Approx.	0.954(0.209)	0.943(0.232)	0.949(0.221)	0.943(0.232)
	0.048(0.004)	0.010(0.001)	0.070(0.008)	0.036(0.006)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.828(0.378)	0.834(0.372)	0.828(0.378)	0.836(0.371)
	0.033(0.003)	0.007(0.001)	0.049(0.007)	0.026(0.004)
	-1.198(0.004)	0.547(0.042)	0.138(0.079)	0.328(0.021)
Green and Shalizi [61]	0.934(0.249)	0.940(0.238)	0.940(0.238)	0.940(0.238)
	0.047(0.005)	0.010(0.002)	0.069(0.010)	0.035(0.006)
	0.574(0.006)	2.077(0.047)	2.548(0.041)	2.099(0.005)
Levin and Levina [93]	0.948(0.222)	0.952(0.214)	0.952(0.214)	0.952(0.214)
	0.048(0.005)	0.010(0.002)	0.070(0.010)	0.036(0.006)
	0.582(0.005)	2.096(0.042)	2.541(0.039)	2.268(0.005)

TABLE 16  
 Performance measures of 95% confidence intervals  
 $n = 160$ ,  $\rho_n \asymp n^{-1/4}$ , graphon: smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.958(0.200)	0.949(0.220)	0.954(0.209)	0.951(0.215)
	Length = 0.045(0.003)	0.005(0.001)	0.049(0.006)	0.020(0.004)
	LogTime = -7.305(0.064)	-6.596(0.222)	-6.151(0.328)	-6.117(0.240)
Norm. Approx.	0.954(0.209)	0.941(0.236)	0.946(0.225)	0.941(0.235)
	0.045(0.003)	0.005(0.001)	0.049(0.006)	0.020(0.004)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.850(0.357)	0.856(0.351)	0.842(0.365)	0.860(0.347)
	0.032(0.003)	0.004(0.001)	0.036(0.005)	0.016(0.003)
	-1.151(0.003)	0.512(0.045)	0.126(0.107)	0.314(0.021)
Green and Shalizi [61]	0.948(0.222)	0.946(0.226)	0.944(0.230)	0.950(0.218)
	0.044(0.004)	0.005(0.001)	0.049(0.007)	0.020(0.004)
	0.624(0.013)	2.004(0.050)	2.533(0.040)	2.103(0.010)
Levin and Levina [93]	0.956(0.205)	0.956(0.205)	0.964(0.186)	0.970(0.171)
	0.046(0.004)	0.006(0.001)	0.051(0.007)	0.022(0.004)
	0.625(0.009)	2.036(0.036)	2.536(0.040)	2.260(0.009)

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TABLE 17  
 Performance measures of 95% confidence intervals  
 $n = 160, \rho_n \asymp n^{-1/4}$ , graphon: non-smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.959(0.198)	0.961(0.194)	0.961(0.195)	0.962(0.192)
	Length = 0.058(0.003)	0.034(0.002)	0.150(0.007)	0.134(0.008)
	LogTime = -7.164(0.080)	-6.145(0.477)	-6.043(0.343)	-5.933(0.330)
Norm. Approx.	0.958(0.201)	0.958(0.201)	0.959(0.198)	0.960(0.196)
	0.058(0.003)	0.034(0.002)	0.150(0.007)	0.134(0.008)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.830(0.376)	0.854(0.353)	0.840(0.367)	0.854(0.353)
	0.040(0.004)	0.024(0.002)	0.104(0.009)	0.093(0.009)
	-1.160(0.005)	0.489(0.054)	0.158(0.046)	0.339(0.018)
Green and Shalizi [61]	0.938(0.241)	0.936(0.245)	0.936(0.245)	0.946(0.226)
	0.056(0.005)	0.033(0.003)	0.145(0.013)	0.130(0.012)
	0.640(0.011)	2.058(0.067)	2.727(0.036)	2.164(0.022)
Levin and Levina [93]	0.952(0.214)	0.952(0.214)	0.954(0.210)	0.952(0.214)
	0.058(0.005)	0.034(0.003)	0.150(0.013)	0.135(0.012)
	0.640(0.013)	2.059(0.060)	2.727(0.037)	2.345(0.015)

TABLE 18  
 Performance measures of 95% confidence intervals  
 $n = 160, \rho_n \asymp n^{-1/2}$ , graphon: block model

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.969(0.173)	0.964(0.186)	0.965(0.183)	0.962(0.192)
	Length = 0.022(0.002)	0.001(0.000)	0.012(0.002)	0.003(0.000)
	LogTime = -7.301(0.074)	-6.859(0.116)	-6.461(0.416)	-6.281(0.245)
Norm. Approx.	0.966(0.182)	0.960(0.195)	0.961(0.195)	0.954(0.210)
	0.022(0.002)	0.001(0.000)	0.012(0.002)	0.003(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.836(0.371)	0.852(0.355)	0.830(0.376)	0.864(0.343)
	0.015(0.002)	0.001(0.000)	0.009(0.001)	0.002(0.000)
	-1.155(0.006)	0.555(0.056)	0.192(0.114)	0.360(0.038)
Green and Shalizi [61]	0.948(0.222)	0.944(0.230)	0.938(0.241)	0.916(0.278)
	0.021(0.002)	0.001(0.000)	0.011(0.002)	0.002(0.000)
	0.624(0.008)	2.060(0.067)	2.850(0.058)	2.142(0.028)
Levin and Levina [93]	0.970(0.171)	0.968(0.176)	0.970(0.171)	0.968(0.176)
	0.022(0.002)	0.001(0.000)	0.013(0.002)	0.003(0.001)
	0.630(0.009)	2.073(0.057)	2.853(0.052)	2.345(0.026)

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TABLE 19  
*Performance measures of 95% confidence intervals*  
 $n = 160$ ,  $\rho_n \asymp n^{-1/2}$ , *graphon: smooth graphon*

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.967(0.178)	0.959(0.198)	0.964(0.186)	0.955(0.208)
	Length = 0.020(0.002)	0.000(0.000)	0.009(0.001)	0.001(0.000)
	LogTime = -7.279(0.075)	-6.744(0.149)	-6.337(0.377)	-6.378(0.215)
Norm. Approx.	0.966(0.183)	0.951(0.216)	0.958(0.201)	0.949(0.220)
	0.020(0.002)	0.000(0.000)	0.009(0.001)	0.001(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.852(0.355)	0.870(0.337)	0.882(0.323)	0.880(0.325)
	0.014(0.001)	0.000(0.000)	0.006(0.001)	0.001(0.000)
	-1.145(0.006)	0.533(0.082)	0.225(0.116)	0.452(0.012)
Green and Shalizi [61]	0.938(0.241)	0.956(0.205)	0.940(0.238)	0.920(0.272)
	0.019(0.002)	0.000(0.000)	0.008(0.001)	0.001(0.000)
	0.611(0.009)	1.997(0.116)	2.856(0.071)	2.328(0.012)
Levin and Levina [93]	0.962(0.191)	0.976(0.153)	0.966(0.181)	0.972(0.165)
	0.020(0.002)	0.000(0.000)	0.009(0.001)	0.002(0.000)
	0.619(0.013)	1.987(0.074)	2.848(0.055)	2.405(0.008)

TABLE 20  
*Performance measures of 95% confidence intervals*  
 $n = 160$ ,  $\rho_n \asymp n^{-1/2}$ , *graphon: non-smooth graphon*

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.972(0.164)	0.974(0.159)	0.974(0.161)	0.975(0.157)
	Length = 0.028(0.002)	0.003(0.000)	0.029(0.002)	0.011(0.001)
	LogTime = -7.007(0.080)	-6.635(0.252)	-6.126(0.364)	-5.835(0.471)
Norm. Approx.	0.972(0.166)	0.973(0.163)	0.972(0.166)	0.973(0.161)
	0.028(0.002)	0.003(0.000)	0.029(0.002)	0.011(0.001)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.852(0.355)	0.868(0.339)	0.858(0.349)	0.874(0.332)
	0.018(0.002)	0.002(0.000)	0.019(0.002)	0.007(0.001)
	-1.159(0.005)	0.565(0.048)	0.149(0.076)	0.376(0.022)
Green and Shalizi [61]	0.948(0.222)	0.954(0.210)	0.956(0.205)	0.946(0.226)
	0.026(0.002)	0.003(0.000)	0.027(0.003)	0.009(0.001)
	0.636(0.019)	2.079(0.101)	2.505(0.041)	2.241(0.015)
Levin and Levina [93]	0.956(0.205)	0.974(0.159)	0.966(0.181)	0.974(0.159)
	0.028(0.002)	0.003(0.000)	0.029(0.003)	0.011(0.001)
	0.640(0.013)	2.062(0.050)	2.519(0.038)	2.462(0.015)

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TABLE 21  
 Performance measures of 95% confidence intervals  
 $n = 160, \rho_n \asymp n^{-1}$ , graphon: block model

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.989(0.106)	0.000(0.000)	0.975(0.157)	0.813(0.390)
	Length = 0.006(0.000)	0.000(0.000)	0.000(0.000)	0.000(0.000)
	LogTime = -7.284(0.063)	-6.939(0.150)	-6.462(0.321)	-6.293(0.247)
Norm. Approx.	0.988(0.109)	0.738(0.440)	0.976(0.154)	0.948(0.221)
	0.006(0.000)	0.000(0.000)	0.000(0.000)	0.000(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.872(0.334)	0.652(0.477)	0.912(0.284)	0.920(0.272)
	0.004(0.000)	Inf(NaN)	0.000(0.000)	Inf(NaN)
	-1.157(0.017)	0.495(0.039)	0.187(0.099)	0.271(0.018)
Green and Shalizi [61]	0.966(0.181)	0.708(0.455)	0.768(0.423)	0.568(0.496)
	0.005(0.001)	Inf(NaN)	0.000(0.000)	0.000(0.000)
	0.619(0.014)	2.030(0.038)	2.855(0.072)	2.150(0.015)
Levin and Levina [93]	0.980(0.140)	0.728(0.445)	0.990(0.100)	0.992(0.089)
	0.006(0.001)	0.000(0.000)	0.001(0.000)	0.000(0.000)
	0.626(0.009)	2.041(0.030)	2.874(0.060)	2.371(0.017)

TABLE 22  
 Performance measures of 95% confidence intervals  
 $n = 160, \rho_n \asymp n^{-1}$ , graphon: smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.989(0.105)	0.000(0.000)	0.964(0.186)	0.651(0.477)
	Length = 0.005(0.000)	0.000(0.000)	0.000(0.000)	0.000(0.000)
	LogTime = -7.290(0.076)	-6.949(0.103)	-6.458(0.388)	-6.277(0.333)
Norm. Approx.	0.987(0.115)	0.437(0.496)	0.968(0.177)	0.926(0.262)
	0.005(0.000)	0.000(0.000)	0.000(0.000)	0.000(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.870(0.338)	0.290(0.456)	0.910(0.288)	0.900(0.302)
	0.003(0.000)	Inf(NaN)	0.000(0.000)	Inf(NaN)
	-0.193(0.061)	1.592(0.024)	1.132(0.038)	1.146(0.012)
Green and Shalizi [61]	0.970(0.171)	0.380(0.488)	0.760(0.429)	0.560(0.499)
	0.005(0.001)	Inf(NaN)	0.000(0.000)	0.000(0.000)
	1.543(0.035)	2.949(0.030)	3.809(0.034)	3.020(0.005)
Levin and Levina [93]	0.980(0.141)	0.390(0.490)	0.990(0.100)	0.990(0.100)
	0.005(0.001)	0.000(0.000)	0.000(0.000)	0.000(0.000)
	1.541(0.033)	2.968(0.033)	3.803(0.038)	3.069(0.009)

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TABLE 23  
 Performance measures of 95% confidence intervals  
 $n = 160$ ,  $\rho_n \asymp n^{-1}$ , graphon: non-smooth graphon

Method	Edge	Triangle	V-shape	Three star
Our method	Coverage = 0.992(0.090)	0.947(0.223)	0.987(0.112)	0.960(0.197)
	Length = 0.008(0.001)	0.000(0.000)	0.001(0.000)	0.000(0.000)
	LogTime = -7.275(0.072)	-6.576(0.324)	-6.121(0.356)	-6.304(0.268)
Norm. Approx.	0.991(0.097)	0.974(0.159)	0.984(0.127)	0.974(0.158)
	0.008(0.001)	0.000(0.000)	0.001(0.000)	0.000(0.000)
	No time cost	No time cost	No time cost	No time cost
Bhattacharyya and Bickel [17]	0.880(0.325)	0.926(0.262)	0.882(0.323)	0.894(0.308)
	0.005(0.001)	Inf(NaN)	0.001(0.000)	0.000(0.000)
	-1.195(0.007)	0.560(0.055)	0.249(0.058)	0.317(0.019)
Green and Shalizi [61]	0.968(0.176)	0.984(0.126)	0.892(0.311)	0.650(0.477)
	0.007(0.001)	Inf(NaN)	0.001(0.000)	0.000(0.000)
	0.575(0.007)	2.011(0.069)	2.893(0.040)	2.187(0.014)
Levin and Levina [93]	0.988(0.109)	0.994(0.077)	0.988(0.109)	0.986(0.118)
	0.008(0.001)	0.000(0.000)	0.001(0.000)	0.000(0.000)
	0.569(0.007)	1.986(0.070)	2.895(0.035)	2.401(0.014)

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